

‘Good’ Pseudo-Random Binary Sequences from Elliptic Curves

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Abstract. We show that the binary sequences, constructed by L.Goubin et al from elliptic curves, possess ‘good’ pseudo-randomness. Namely, such sequences are of ‘small’ well-distribution measure and ‘small’ correlation measure of ‘small’ order, both of which were introduced by C.Mauduit et al to analyze the pseudo-randomness of binary sequences. The results give a partial answer to a conjecture proposed by L.Goubin et al.

Keywords. pseudorandom sequences, elliptic curves, exponential sums, well-distribution, correlation.

1 Introduction

Mauduit and Sárközy [15] introduced several measures to evaluate the (local) pseudo-randomness of a finite binary sequence:

$$S_N = \{s_1, s_2, \dots, s_N\} \in \{+1, -1\}^N.$$

The most two important measures are the well-distribution measure and the correlation measure of order k .

The *well-distribution measure* of S_N is defined as

$$W(S_N) = \max_{a,b,t} \left| \sum_{j=0}^{t-1} s_{a+jb} \right|,$$

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where the maximum is taken over all a, b, t such that $a, b, t \in \mathbb{N}$ and $1 \leq a \leq a + (t - 1)b \leq N$, while the *correlation measure of order k* (or *order k correlation measure*) of S_N is defined as

$$C_k(S_N) = \max_{M, D} \left| \sum_{n=1}^M s_{n+d_1} s_{n+d_2} \cdots s_{n+d_k} \right|,$$

where the maximum is taken over all $D = (d_1, \dots, d_k)$ with non-negative integers $0 \leq d_1 < \dots < d_k$ and M such that $M + d_k \leq N$.

S_N is considered as a “good” pseudo-random sequence, if both $W(S_N)$ and $C_k(S_N)$ (at least for small k) are “small” in terms of N (in particular, both are $o(N)$ as $N \rightarrow \infty$). It was shown in [2] that for a “truly” random sequence $S_N \in \{+1, -1\}^N$ (i.e., choosing $S_N \in \{+1, -1\}^N$ with probability $1/2^N$), both $W(S_N)$ and $C_k(S_N)$ (for some fixed k) are around $N^{1/2}$ with “near 1” probability.

For the Legendre sequence $S_p = \{s_1, s_2, \dots, s_p\} \in \{+1, -1\}^p$ with

$$s_n = \begin{cases} \left(\frac{n}{p}\right), & \text{if } \gcd(n, p) = 1; \\ 1, & \text{if } p|n, \end{cases}$$

it was shown by Mauduit and Sárközy in [15] that

$$W(S_p) = O(p^{1/2} \log(p)) \quad \text{and} \quad C_k(S_p) = O(kp^{1/2} \log(p)),$$

which indicate that the Legendre sequence forms a “good” pseudo-random sequence. Many other “good” (but slightly inferior) binary sequences were designed in the literature, see for example [2, 3, 7, 8, 16, 18, 19] and references therein.

Recent developments point towards an interest in the elliptic curve analogues of pseudo-random number generators, see [1, 5, 6, 9, 10, 11, 12, 13, 14, 17, 20, 22] and references therein. Such number generators provide strong potential applications in cryptography for generating pseudo-random numbers and session keys.

In [7] Goubin et al applied elliptic curves to construct binary sequences, whose terms were represented by two elements $+1$ and -1 , and analyzed their pseudo-randomness. In fact, the authors of [7] only listed some examples there and proposed a conjecture that such binary sequences have ‘small’ well-distribution measure and ‘small’ correlation measure of order k (for some small value k). We will apply exponential sums on elliptic curves to show that this conjecture is correct.

In the present paper, we use 0 and 1 to represent the terms of a binary sequence as usual. We first introduce some notions and basic facts of elliptic curves over finite fields. Let $p > 3$ be a prime and \mathbb{F}_p the finite field of p elements, which we identify with the set $\{0, 1, \dots, p-1\}$. \mathbb{F}_p^* is the set of non-zero elements of \mathbb{F}_p . Let \mathcal{E} be an elliptic curve over \mathbb{F}_p , given by an affine Weierstrass equation of the standard form

$$y^2 = x^3 + ax + b, \quad (1)$$

with coefficients $a, b \in \mathbb{F}_p$ and nonzero discriminant, see [4]. It is known that the set $\mathcal{E}(\mathbb{F}_p)$ of \mathbb{F}_p -rational points of \mathcal{E} forms an Abelian group under an appropriate composition rule denoted by \oplus and with the point at infinity \mathcal{O} as the neutral element. We recall that

$$|\#\mathcal{E}(\mathbb{F}_p) - p - 1| \leq 2p^{1/2},$$

where $\#\mathcal{E}(\mathbb{F}_p)$ is the number of \mathbb{F}_p -rational points, including the point at infinity \mathcal{O} . Let $G \in \mathcal{E}(\mathbb{F}_p)$ be a point of order N , that is, N is the size of the cyclic group $\langle G \rangle$ generated by G . A multiple of a point P is taken by $nP = \oplus_{i=1}^n P$. We write $iG = (x_i, y_i) \in \mathbb{F}_p \times \mathbb{F}_p$ on \mathcal{E} for all $1 \leq i \leq N-1$.

Five types of finite binary sequences $S_{N-1} = \{s_1, \dots, s_{N-1}\}$, constructed in [7], are described in the following with a slight modification:

$$\begin{aligned} \text{Construction I :} & \quad s_i := \begin{cases} 1, & y_i > \frac{p}{2}; \\ 0, & \text{otherwise.} \end{cases} \\ \text{Construction II :} & \quad s_i := \begin{cases} 1, & x_i > \frac{p}{2}; \\ 0, & \text{otherwise.} \end{cases} \\ \text{Construction III :} & \quad s_i := \begin{cases} 1, & y_i \text{ is even;} \\ 0, & \text{otherwise.} \end{cases} \\ \text{Construction IV :} & \quad s_i := \begin{cases} 1, & x_i \text{ is even;} \\ 0, & \text{otherwise.} \end{cases} \\ \text{Construction V :} & \quad s_i := \begin{cases} 1, & x_i < y_i; \\ 0, & \text{otherwise.} \end{cases} \end{aligned}$$

In fact, Construction I has been proposed in [12] and the period and the linear complexity of this sequence has also been considered.

In the present paper, we want to prove that these constructions indeed produce ‘good’ pseudo-random sequences. Namely, we show that both the well-distribution measure and the correlation measure of ‘small’ order of the above five sequences are ‘small’. The proof is based on some bounds of character sums over subgroups of the point group of elliptic curves [11].

Throughout this paper, the implied constant in the symbol “ \ll ” may sometimes depend on the integer $\deg(f)$, the degree of a rational function f , and is absolute otherwise.

2 Preparations

Let \mathcal{E} be an elliptic curve over \mathbb{F}_p defined as Eq.(1). Let $f \in \mathbb{F}_p(\mathcal{E})$ be a rational function. We denote by $\deg(f)$ the degree of the pole divisor of f . In particular, $\deg(f) = 2$ if $f = x$ and $\deg(f) = 3$ if $f = y$. The translation map by $W \in \mathcal{E}(\mathbb{F}_p)$ on $\mathcal{E}(\mathbb{F}_p)$ is defined as follows:

$$\tau_W : P \mapsto P \oplus W.$$

It is obvious that $(f \circ \tau_W)(P) = f(\tau_W(P)) = f(P \oplus W)$. We denote by \ominus the inverse operation of \oplus in the rational points group of \mathcal{E} . From Lemma 3.16, Theorem 3.17 and Lemma 3.14 of [4], we have the following statement.

Lemma 1 *Let $f \in \mathbb{F}_p(\mathcal{E})$ be a rational function. If f has a pole at $H \in \mathcal{E}(\overline{\mathbb{F}}_p)$ of multiplicity ρ , then $f \circ \tau_W$ has a pole at $H \ominus W$ of the same multiplicity ρ .*

Let $e_p(z) = \exp(2\pi iz/p)$ be an additive character of \mathbb{F}_p . For any positive m , an additive character of $\mathbb{Z}_m := \{0, 1, \dots, m-1\}$, the residue ring modulo m , is defined as $e_m(z) = \exp(2\pi iz/m)$. We also need the following upper bound which is a special case of Corollary 1 of [11].

Lemma 2 *Let $f \in \mathbb{F}_p(\mathcal{E})$ be a nonconstant rational function and $G \in \mathcal{E}(\mathbb{F}_p)$ be a rational point of order N . Then the bound*

$$\left| \sum_{\substack{z=0 \\ f(zG) \neq \infty}}^{N-1} e_p(\lambda f(zG)) e_N(\eta z) \right| \leq 2 \deg(f) p^{1/2}$$

holds for all $\lambda \in \mathbb{F}_p^$ and $\eta \in \mathbb{Z}_N$.*

Lemma 3 *Let p be an odd prime number and $\lambda \in \mathbb{Z}$ with $0 \leq |\lambda| \leq \frac{p-1}{2}$. We define*

$$V(\lambda) := \sum_{r=0}^{(p-1)/2} e_p(-\lambda r) - \sum_{r=1}^{(p-1)/2} e_p(\lambda r), \quad (2)$$

$$U(\lambda) := \sum_{r=1}^{(p-1)/2} e_p(2\lambda r) - \sum_{r=0}^{(p-1)/2} e_p(-2\lambda r) \quad (3)$$

and

$$W(\lambda, u) := \sum_{r=0}^u e_p(-\lambda r) - \sum_{r=u+1}^{p-1} e_p(-\lambda r), \quad (4)$$

where $0 \leq u \leq p-1$. Then the following bounds hold:

$$\begin{aligned} \sum_{|\lambda| \leq (p-1)/2} |V(\lambda)| &\leq 2p(1 + \log p) ; \\ \sum_{|\lambda| \leq (p-1)/2} |U(\lambda)| &\leq 2p(1 + \log p) ; \end{aligned}$$

and

$$\sum_{|\lambda| \leq (p-1)/2} |W(\lambda, u)| \leq 2p(1 + \log p).$$

Proof. Since $|V(\lambda)| \leq \left| \sum_{r=0}^{(p-1)/2} e_p(-\lambda r) \right| + \left| \sum_{r=1}^{(p-1)/2} e_p(\lambda r) \right|$, the first desired result follows from Inequality (3.4) of [21]. The other two cases are similar. \square

Lemma 4 *Let N be a positive integer, $1 \leq b \leq N-1$ and $t \in \mathbb{N}$ with $(t-1)b \leq N-1$. Then the following bound holds:*

$$\sum_{\lambda=0}^{N-1} \left| \sum_{x=0}^{t-1} e_N(\lambda bx) \right| \ll N \log N.$$

Proof. Let $d = \gcd(b, N)$, $M = N/d$ and $b_1 = b/d$. Since $(t-1)b \leq N-1$, we have $d(t-1) \leq (t-1)b < N$, and hence $t-1 < M$. We derive

$$\sum_{\lambda=0}^{N-1} \left| \sum_{x=0}^{t-1} e_N(\lambda bx) \right| = d \sum_{\lambda=0}^{M-1} \left| \sum_{x=0}^{t-1} e_N(\lambda bx) \right| = d \sum_{\lambda=0}^{M-1} \left| \sum_{x=0}^{t-1} e_M(\lambda b_1 x) \right| \ll dM \log M.$$

Since $\gcd(M, b_1) = 1$, the last inequality holds by Inequality (3.4) of [21]. \square

Lemma 5 *Let $G \in \mathcal{E}(\mathbb{F}_p)$ be of order N and $f \in \mathbb{F}_p(\mathcal{E})$ a nonconstant rational function. Then for any fixed $a, b, t \in \mathbb{N}$ with $1 \leq a \leq a + (t-1)b \leq N-1$, the following bound holds:*

$$\left| \sum_{x=0}^{t-1} e_p(f((a+bx)G)) \right| \ll p^{1/2} \log N.$$

Proof.

$$\begin{aligned}
\left| \sum_{x=0}^{t-1} e_p(f((a+bx)G)) \right| &= \left| \frac{1}{N} \sum_{n=0}^{N-1} \sum_{x=0}^{t-1} e_p(f(nG)) \sum_{\lambda=0}^{N-1} e_N(\lambda(n - (a+bx))) \right| \\
&= \frac{1}{N} \left| \sum_{\lambda=0}^{N-1} \sum_{x=0}^{t-1} e_N(-\lambda(a+bx)) \sum_{n=0}^{N-1} e_p(f(nG)) e_N(\lambda n) \right| \\
&\leq \frac{1}{N} \sum_{\lambda=0}^{N-1} \left| \sum_{x=0}^{t-1} e_N(-\lambda(a+bx)) \right| \cdot \left| \sum_{n=0}^{N-1} e_p(f(nG)) e_N(\lambda n) \right| \\
&= \frac{1}{N} \sum_{\lambda=0}^{N-1} \left| \sum_{x=0}^{t-1} e_N(-\lambda bx) \right| \cdot \left| \sum_{n=0}^{N-1} e_p(f(nG)) e_N(\lambda n) \right|.
\end{aligned}$$

Now by Lemmas 2 and 4, we derive the desired result. We note that in the above formulae the poles of f must be ruled out. \square

3 Pseudorandomness of Elliptic Curve Sequences

In this section we will present an upper bound respectively for the well-distribution measure $W(S_N)$ and the correlation measure $C_k(S_N)$ for binary sequences defined in Construction I-V.

Assume that f is a rational function and $f = x$ or $f = y$ in the following context. We remark that $x(iG) = x_i$ and $y(iG) = y_i$ for $iG = (x_i, y_i) \in \mathcal{E}(\mathbb{F}_p)$. For Construction I and II, for any $1 \leq i \leq N-1$, we have

$$\frac{1}{p} \sum_{r=1}^{(p-1)/2} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(f(iG) + r)) = \begin{cases} 1, & p > f(iG) \geq (p+1)/2, \\ 0, & \text{otherwise.} \end{cases} \quad (5)$$

$$\frac{1}{p} \sum_{r=0}^{(p-1)/2} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(f(iG) - r)) = \begin{cases} 0, & p > f(iG) \geq (p+1)/2, \\ 1, & \text{otherwise.} \end{cases} \quad (6)$$

Subtracting (5) from (6) yields

$$\frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda f(iG)) V(\lambda) = \begin{cases} -1, & p > f(iG) \geq (p+1)/2, \\ 1, & \text{otherwise.} \end{cases}$$

where $V(\lambda)$ is defined as (2) in Lemma 3. It is easy to see that

$$(-1)^{s_i} = \frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda f(iG)) V(\lambda), \quad (7)$$

where $f = y$ for Construction I and $f = x$ for Construction II.

While for Construction III and IV, for any $1 \leq i \leq N - 1$, we have

$$\frac{1}{p} \sum_{r=0}^{(p-1)/2} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(f(iG) - 2r)) = \begin{cases} 1, & f(iG) \text{ is even,} \\ 0, & \text{otherwise.} \end{cases} \quad (8)$$

$$\frac{1}{p} \sum_{r=1}^{(p-1)/2} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(f(iG) + 2r)) = \begin{cases} 0, & f(iG) \text{ is even,} \\ 1, & \text{otherwise.} \end{cases} \quad (9)$$

(9)–(8), we get

$$\frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda f(iG))U(\lambda) = \begin{cases} -1, & f(iG) \text{ is even,} \\ 1, & \text{otherwise.} \end{cases}$$

where $U(\lambda)$ is defined as (3) in Lemma 3. Similar to (7), we obtain

$$(-1)^{s_i} = \frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda f(iG))U(\lambda), \quad (10)$$

where $f = y$ for Construction III and $f = x$ for Construction IV.

For Construction V, the following two formulae hold for all x_i with $0 \leq x_i \leq p - 1$:

$$\frac{1}{p} \sum_{r=x_i+1}^{p-1} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(y_i - r)) = \begin{cases} 1, & x_i < y_i, \\ 0, & \text{otherwise.} \end{cases} \quad (11)$$

$$\frac{1}{p} \sum_{r=0}^{x_i} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda(y_i - r)) = \begin{cases} 0, & x_i < y_i, \\ 1, & \text{otherwise.} \end{cases} \quad (12)$$

(12)–(11), we get

$$\frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda y_i)W(\lambda, x_i) = \begin{cases} -1, & x_i < y_i, \\ 1, & \text{otherwise.} \end{cases}$$

where $W(\lambda, x_i)$ is defined as (4) in Lemma 3. Hence we obtain

$$(-1)^{s_i} = \frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} e_p(\lambda y_i)W(\lambda, x_i). \quad (13)$$

Theorem 1 Assume that $G \in \mathcal{E}(\mathbb{F}_p)$ is a point of order N and S_{N-1} is one of binary sequences obtained from Construction I – V. Then the well-distribution measure of S_{N-1} holds:

$$W(S_{N-1}) \ll p^{1/2} \log p \log N.$$

Proof. We only prove the statement for S_{N-1} obtained from Construction I and II. Combining with (10), (13) and Lemma 3, one can prove the other three cases in a similar way. For any $a, b, t \in \mathbb{N}$ with $1 \leq a \leq a + (t-1)b \leq N-1$, from (7) we have

$$\begin{aligned} \left| \sum_{j=0}^{t-1} (-1)^{s_{a+jb}} \right| &= \frac{1}{p} \left| \sum_{j=0}^{t-1} \sum_{|\lambda| \leq (p-1)/2} V(\lambda) e_p(\lambda f((a+jb)G)) \right| \\ &= \frac{1}{p} \left| \sum_{|\lambda| \leq (p-1)/2} V(\lambda) \sum_{j=0}^{t-1} e_p(\lambda f((a+jb)G)) \right| \\ &\leq \frac{1}{p} \sum_{|\lambda| \leq (p-1)/2} |V(\lambda)| \cdot \left| \sum_{j=0}^{t-1} e_p(\lambda f((a+jb)G)) \right| \\ &\leq \frac{1}{p} \left(\sum_{|\lambda|=1}^{(p-1)/2} |V(\lambda)| \cdot \left| \sum_{j=0}^{t-1} e_p(\lambda f((a+jb)G)) \right| + t \right). \end{aligned}$$

Now by Lemmas 3 and 5, we obtain the desired result. \square

Theorem 2 Assume that $G \in \mathcal{E}(\mathbb{F}_p)$ is a point of order N and S_{N-1} is one of binary sequences obtained from Construction I – V. Then the correlation measure of order k ($k < p$) holds:

$$C_k(S_{N-1}) \ll k 2^k p^{1/2} (\log p)^k \log N.$$

Proof. Similar to Theorem 1, we only prove the statement for S_{N-1} obtained from Construction I and II. For $D = (d_1, \dots, d_k)$ and M with $0 \leq d_1 < \dots < d_k \leq N-1-M$, we have

$$\begin{aligned} &\left| \sum_{n=1}^M (-1)^{s_{n+d_1} + \dots + s_{n+d_k}} \right| \\ &= \left| \sum_{n=1}^M \prod_{i=1}^k \left(\frac{1}{p} \sum_{|\lambda_i| \leq (p-1)/2} V(\lambda_i) e_p(\lambda_i f((n+d_i)G)) \right) \right| \\ &= \frac{1}{p^k} \left| \sum_{|\lambda_1| \leq (p-1)/2} \dots \sum_{|\lambda_k| \leq (p-1)/2} V(\lambda_1) \dots V(\lambda_k) \sum_{n=1}^M e_p\left(\sum_{i=1}^k \lambda_i f((n+d_i)G)\right) \right| \\ &\leq \frac{1}{p^k} \left(2k \deg(f) p^{1/2} \log N \left(\sum_{|\lambda| \leq (p-1)/2} |V(\lambda)| \right)^k + M \right). \end{aligned}$$

The last inequality holds since the degree of the rational function $\sum_{i=1}^k f \circ \tau_{d_i G}$ is at most $k \deg(f)$ by Lemma 1. The desired result follows from Lemmas 3 and 5. \square

Theorems 1 and 2 indicate that the five types binary sequences are “good” sequences. But it seems that they are slightly inferior to Legendre sequences.

There are a large family of elliptic curves over \mathbb{F}_p with a rational point of large order N . In particular, if $\mathcal{E}(\mathbb{F}_p)$ is a cyclic group, then $N \sim p$. As indicated in [10], from Corollary 6.2 of [23], about 75% of the majority of (isomorphism classes of) elliptic curves have a cyclic point group. By Theorem 2.1 of [23], every cyclic group of order N satisfying $p - 1 - 2p^{1/2} \leq N \leq p - 1 + 2p^{1/2}$ can be realized as the point group of an elliptic curve over \mathbb{F}_p ($p > 5$). An elliptic curve with a rational point of large prime order is necessary for elliptic curve cryptosystems. More information on elliptic curves with cyclic groups can be found in [23, 24].

4 Conclusion

For the binary sequences produced by Construction I-V from a large family of elliptic curves, both the well-distributed measure and the correlation measure of ‘small’ order are ‘small’. The results give a (partially) positive answer to the conjecture proposed by L.Goubin et al in [7] on the randomness of the corresponding sequences. These constructions provide a very attractive alternative to traditional methods.

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