# Unbalanced Elementary Symmetric Boolean Functions With The Degree **d** And $wt(d) \ge 3^*$

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#### Abstract

In the paper, for  $d = 2^t k$ ,  $n = 2^t (2k+q) + m$  and special  $k = 2^w (2^0 + 2^1 + \dots + 2^s)$ , we present that a majority of X(d, n) are not balanced. The results include many cases  $wt(d) \ge 3$  and  $n \equiv 0, 1, 2, 3mod4$ . The results are also parts of the conjecture that  $X(2^t, 2^{t+1}l - 1)$  is only nonlinear balanced elementary symmetric Boolean function. Where  $t \ge 2, q \ge 1, s \ge 0, w \ge 0$  and  $m \ge -1$  are integers, and  $X(d, n) = \bigoplus_{\substack{i=1 \le i_1 \le \dots \le i_d \le n}} x_{i_1} \cdots x_{i_d}$ .

**Keywords**: Cryptograph, Boolean functions, balancedness, elementary symmetric.

#### 1 Introduction

Symmetric Boolean function is a subclass of Boolean functions and their outputs only depend on the Hamming weifht of their inputs, namely, for Boolean function f(x), inputs x and y, then f(x) = f(y) when wt(x) = wt(y). They allow reducing memory spaces and gates of hardware implementation and are of great interest to cryptography. Recent years, many significant properties of symmetric Boolean functions have been studied in [1–7], including balancedness, algebraic immunity, resiliency, nonlinearity and so on. In [1–3], some symmetric Boolean functions with maximum algebraic immunity were constructed. The [3] gave the enumeration of symmetric Boolean functions with maximum algebraic is respectively  $2^{n-1} - 2^{n/2-1}$  and  $2^{n-1} - 2^{(n-1)/2}$  when n is respectively even and odd. The [6] gave all balanced symmetric Boolean functions whose degrees are smaller than 7. The [6] and [7] investigated the relationships among the significant properties of symmetric Boolean functions.

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It is well known that balancedness is a primary requirement for Boolean functions in cryptosystem. The balancedness of symmetric Boolean functions should be firstly studied. For fixed algebraic degree, the [6] proved the conjecture that there is not balanced symmetric Boolean function when ngrows. As a subclass of symmetric Boolean functions, elementary symmetric Boolean function is basic unit composing of symmetric Boolean functions. The balancedness of Elementary symmetric Boolean functions have been studied in [8–10]. The [8] prosed a conjecture that  $X(2^t, 2^{t+1}l - 1)$  is only nonlinear balanced elementary symmetric Boolean function. The [9] proved the conjecture holds when wt(d) < 3 and gave some cases that X(d, n) are not balanced when wt(d) = 3. However, the [9] didn't study further when wt(d) > 3. In [10], for  $n = 2^{t+1}l - 1$  with odd l and  $2^{t+1} \nmid d$ , it showed that X(d,n) is balanced if and only if  $d = 2^k, 1 \le k \le t$ . Hence, for  $n = 2^{t+1}l - 1$ with odd l, the only case left is  $2^{t+1} \mid d$ . A majority of conjecture have been proved when  $n \equiv 3mod4$ , however, there are not many results when  $n \equiv 0, 1, 2mod4.$ 

Since X(d, n) is not balanced when  $d > \lceil n/2 \rceil$  [8], we consider special elementary symmetric Boolean functions with forms  $d \leq \lceil n/2 \rceil$ . Combining with [9] and [10], in the paper, we consider special forms  $2^t \mid d, n = 2^t l + m$ and  $d \leq \lceil n/2 \rceil$ . We assume that  $d = k2^t$  and  $n = 2^t(2k+q)+m$ . As the cases wt(d) < 3 were discussed in [9] and the general cases  $wt(d) \geq 3$  are difficult to be discussed, we consider special d whose '1s' are consecutive in the 2-adic description. Namely,  $d = 2^{t+w}(1+2^1+\cdots+2^s)$  and  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq+m$ . For the several kinds of elementary symmetric Boolean functions, we give most cases that X(d, n) are not balanced. The results include many cases  $wt(d) \geq 3$  and  $n \equiv 0, 1, 2, 3mod4$ .

#### 2 Preliminaries

There are some general definitions about Boolean functions. Denote by GF(2) the finite field with two elements 0 and 1, and denote by  $\bigoplus$  the addition over GF(2). We consider function f(x) called *n*-variable Boolean function from  $GF^n(2)$  to GF(2), where  $GF^n(2)$  is the *n*-dimensional vector space over GF(2) and  $x = (x_1, x_2, \ldots, x_n) \in GF^n(2)$ . f(x) can be represented as a polynomial, called its algebraic normal form (ANF):

$$f(x_1, \dots, x_n) = \bigoplus_{u \in GF^n(2)} \lambda_u(\prod_{i=1}^n x_i^{u_i}), \qquad \lambda_u \in GF(2).$$

The number of variables in the highest order product term with nonzero coefficient is called its algebraic degree. The Hamming weight of a binary vector  $x = (x_1, x_2, ..., x_n)$  is the number of its nonzero coordinates, denoted by wt(x). Denote by |A| the size of the group A.  $|\{x \in GF^n(2)|f(x) = 1\}|$  is called the Hamming weight of Boolean function f(x), denoted by wt(f(x)).

f(x) is called balanced if  $wt(f(x)) = 2^{n-1}$ . Hence,  $| \{x \in GF^n(2) | f(x) = 0\} | -wt(f(x)) = 2^n - 2wt(f(x))$  can refer the balancedness of f(x), namely, f(x) is balanced if and only if  $2^n - 2wt(f(x)) = 0$ .

An *n*-variable Boolean function f(x) is called symmetric if its output is invariant under any permutation of its input bits. Equivalently, the output of f(x) only depends on the Hamming weight of its input vector. The form of elementary symmetric Boolean functions is as follows:

$$X(d,n) = \bigoplus_{1 \le i_1 < \dots < i_d \le n} x_{i_1} \cdots x_{i_d}.$$

Let  $Z(d,n) = 2^n - 2wt(X(d,n))$ , then X(d,n) is balanced if and only if Z(d,n) = 0. We write  $\binom{n}{i}$  as  $C_n^i$  for short, and it is easy to get  $X(d,n) = [1 - (-1)^{C_i^d}]/2$  when  $x \in GF^n(2)$  and wt(x) = i. Hence,

$$Z(d,n) = 2^{n} - 2 \cdot \sum_{i=0}^{n} [C_{n}^{i} \cdot \frac{1 - (-1)^{C_{i}^{d}}}{2}]$$
  
$$= \sum_{i=0}^{d-1} C_{n}^{i} + \sum_{i=d}^{n} [C_{n}^{i} \cdot (-1)^{C_{i}^{d}}].$$
(1)

**Definition 1.** [2] Let  $a = (a_1, \dots, a_n) \in GF^n(2)$ ,  $b = (b_1, \dots, b_n) \in GF^n(2)$ , we say  $a \leq b$  if  $a_i \leq b_i$  for all  $1 \leq i \leq n$ ; we say  $a \neq b$  if  $a_i > b_i$  for some *i*.

**Lemma 1.** [11] Let k and t be nonnegative integers,  $k \ge t$ , their 2-adic descriptions are  $a = (k_0, k_1, \dots, k_l)$  and  $t = (t_0, t_1, \dots, t_l)$ , then

$$C_k^t \equiv C_{k_0}^{t_0} C_{k_1}^{t_1} \cdots C_{k_l}^{t_l} \equiv \begin{cases} 1mod2 & , & a \leq b \\ 0mod2 & , & a \neq b \end{cases}$$

We get the following two lemmas from basic knowledge of mathematical.

**Lemma 2.** For fixed real number a > 1 and b, then  $a^x > bx$  when x > N(a, b), where N(a, b) is a real number and is only relative to a and b.

**Lemma 3.** For fixed real number a > 1, then  $x^a - o(x^a) > 0$  when x > N(a), where N(a) is a real number and is only relative to a.  $o(x^a)$  is higher order indefinite small than  $x^a$ , namely,  $x^a/o(x^a) \to \infty$  when  $x \to \infty$ .

**3** When  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq-1$ 

In the section, we discuss elementary symmetric Boolean functions with form  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq-1$  and  $d = 2^{w+t}(1+2^1+\cdots+2^s)$ . Notice that  $n \equiv 3mod4$ , the section is further work of the [10].

**Theorem 1.** Let q > 0, t > 1, w and s be nonnegative integers,  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq-1$  and  $d = 2^{w+t}(1+2^1+\cdots+2^s)$ . For fixed s and q, then Z(d,n) < 0 when  $w \ge N(s,q)$ , where N(s,q) is a nonnegative integer and is only relative to s and q.

*Proof.* Let  $S = 2^0 + 2^1 + \dots + 2^s$ , then  $d = 2^{w+t}S$  and  $n = 2^{w+t+1}S + 2^tq - 1 = 2d + 2^tq - 1$ . We have  $d \leq i$  when  $d \leq i < d + 2^{w+t}$ . Assume that  $2^w \geq q$ , then  $d \not\leq i$  when  $d + 2^{w+t} \leq i \leq n$ . Since lemma 1, we have  $C_i^d \equiv 1 \mod 2$  when  $d \leq i < d + 2^{w+t}$  and  $C_i^d \equiv 0 \mod 2$  when  $d + 2^{w+t} \leq i \leq n$ . Note that  $C_n^i = C_n^{n-i}$  for all  $0 \leq i \leq n$ . Hence,

$$Z(d,n) = \sum_{i=0}^{2^{w+t}S-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{w+t}-1} C_n^i + \sum_{i=2^{w+t}S+2^{w+t}}^n C_n^i$$

$$= \sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i + \sum_{i=2^{w+t}S+2^tq-2^{w+t}}^{2^{w+t}S-1} C_n^i$$

$$- \sum_{i=2^{w+t}S}^{(n-1)/2} C_n^i - \sum_{i=(n+1)/2}^{n-2^{w+t}S} C_n^i$$

$$- \sum_{i=n-2^{w+t}S+1}^{n-2^{w+t}S+2^tq-2^{w+t}} C_n^i + \sum_{i=n-2^{w+t}S-2^tq+2^{w+t}+1}^n C_n^i$$

$$= 2 \cdot (\sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{t-1}q-1} C_n^i). \quad (2)$$

$$\triangleq 2(A-B)$$

For

$$\frac{C_n^{2^{w+t}S+2^{t-1}q-1}}{C_n^{2^{w+t}S+2^tq-2^{w+t}-1}} = \frac{(2^{w+t}S+2^{w+t}) \times \dots \times (2^{w+t}S+2^{t-1}q+1)}{(2^{w+t}S+2^{t-1}q-1) \times \dots \times (2^{w+t}S+2^tq-2^{w+t})} \\
\geq (\frac{2^{w+t}S+2^{w+t}}{2^{w+t}S+2^{t-1}q-1})^{2^{w+t}-2^{t-1}q}.$$
(3)

And notice that  $2^w \ge q$ ,

$$\frac{2^{w+t}S + 2^{w+t}}{2^{w+t}S + 2^{t-1}q - 1} = \frac{S+1}{S + \frac{2^{t-1}q - 1}{2^{w+t}}} > \frac{S+1}{S+0.5} = \frac{2S+2}{2S+1}.$$
 (4)

Then, on the one hand, we have the following inequations from (3) and (4).

$$B > C_n^{2^{w+t}S+2^{t-1}q-1} > (\frac{2S+2}{2S+1})^{2^{w+t}-2^{t-1}q} C_n^{2^{w+t}S+2^tq-2^{w+t}-1} \triangleq C.$$
(5)

on the other hand, we have the following inequations from  $2^w \ge q$ .

$$A < (2^{w+t}S + 2^{t}q - 2^{w+t})C_{n}^{2^{w+t}S + 2^{t}q - 2^{w+t} - 1} \leq 2S(2^{w+t} - 2^{t-1}q)C_{n}^{2^{w+t}S + 2^{t}q - 2^{w+t} - 1} \triangleq D.$$
(6)

Since lemma 2, then there exists a positive real number  $N_1(S)$  which is only relative to S. When  $N_1(S) \ge 2^{w+t} - 2^{t-1}q$ , namely, we have the following inequation when  $2^w \ge N_1(S)/2^t + q/2$ ,

$$\left(\frac{2S+2}{2S+1}\right)^{2^{w+t}-2^{t-1}q} \ge 2S(2^{w+t}-2^{t-1}q).$$
(7)

Therefore, from (2), (5), (6) and (7), if  $2^w \ge q$  and  $2^w \ge N_1(S)/2^t + q/2$  hold at the same time, then

$$Z(d,n) = 2(A - B) < 2(D - C) < 0.$$
(8)

Since  $N_1(S)/2^t + q/2 \le N_1(S)/2 + q/2$ , if we let

$$N(s,q) = \lceil max\{\log_2 q, \log_2 (N_1(S)/2 + q/2)\} \rceil$$
(9)

then Z(d,n) < 0 when  $w \ge N(s,q)$ .

**Remark 1.** In fact, according to theorem 1 and computer exhausting,  $Z(d, n) \neq 0$  when s and q are enough small. Notice that  $N(s,q) = \lceil \log_2 q \rceil$  when  $q \geq N_1(S)$ ,  $N(s,q) = \lceil \log_2 (N_1(S)/2 + q/2) \rceil < \lceil \log_2 N_1(S) \rceil$  when  $q < N_1(S)$ . We can let  $N(s,q) = \lceil \log_2 N_1(S) \rceil$  when  $q < N_1(S)$ . The following table 1 presents the relationships clear. From the table 1, we notice that the relationship between s and N(s,q) is almost linearity when  $q < N_1(S)$ .

Table 1: The relationships among s, q and N(s,q)

s	q <	N(s,q)	s	q <	N(s,q)	s	q <	N(s,q)
0	11	4	7	7772	13	14	1665654	21
1	42	6	8	17068	15	15	3520258	22
2	115	7	9	37156	16	16	7417616	23
3	286	9	10	80318	17	17	15588014	24
4	676	10	11	172590	18	18	32679052	25
5	1554	11	12	368998	19	19	68359552	27
6	3500	12	13	785478	20	20	142713644	28

**Theorem 2.** The conditions are the same as theorem 1. For fixed w, q and t, if  $2^w \ge q$ , then Z(d,n) > 0 when  $s \ge N(w,q,t)$ , where N(w,q,t) is a nonnegative integer and is only relative to w, q and t.

*Proof.* We still let  $S = 2^0 + 2^1 + \cdots + 2^s$ , From the proof of the theorem 1, we have

$$Z(d,n) = 2 \cdot \left(\sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{t-1}q-1} C_n^i\right)$$
  
> 
$$2 \cdot \left(\sum_{i=2^{w+t}S+2^{t}q-2^{w+t}-1}^{2^{w+t}S+2^{t-1}q-2^{w+t}-1} C_n^i - 2^{t-1}qC_n^{2^{w+t}S+2^{t-1}q-1}\right) (10)$$
  
$$\triangleq E.$$

If we assume that  $\prod_{j=0}^{i-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q - j) = 1$  when i = 0, for any  $0 \le i \le 2^{t-1}q$  and  $k = i + 2^{w+t}S + 2^{t-1}q - 2^{w+t} - 1$ , then

$$\frac{C_n^k \cdot (2^{w+t}S + 2^{t-1}q - 1)!(2^{w+t}S + 2^{w+t} + 2^{t-1}q)!}{(2^{w+t+1}S + 2^tq - 1)!} = \prod_{j=0}^{2^{w+t}-i-1} (2^{w+t}S + 2^{t-1}q - 1 - j) \cdot \prod_{j=0}^{i-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q - j). (11)$$

Hence, we have the following equations from (10) and (11).

$$\frac{E \cdot (2^{w+t}S + 2^{t-1}q - 1)!(2^{w+t}S + 2^{w+t} + 2^{t-1}q)!}{2 \cdot (2^{w+t+1}S + 2^{t}q - 1)!} = \sum_{i=0}^{2^{t-1}q} [\prod_{j=0}^{2^{w+t}-i-1} (2^{w+t}S + 2^{t-1}q - 1 - j) \cdot \prod_{j=0}^{i-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q - j)] - 2^{t-1}q \prod_{i=0}^{2^{w+t}-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q - i) = (2^{t-1}q + 1) \cdot [S^{2^{w+t}} + o(S^{2^{w+t}})] - 2^{t-1}q[S^{2^{w+t}} + o(S^{2^{w+t}})] = S^{2^{w+t}} - o(S^{2^{w+t}}).$$
(12)

Note that the last two equations are relative to q, since lemma 3, then there exists a positive real number  $N_1(w, q, t)$  which is only relative to t, w and q.  $S^{2^{w+t}} - o(S^{2^{w+t}}) > 0$  when  $S = 1 + 2^1 + 2^2 + \cdots + 2^s > N_1(w, q, t)$ . If we let

$$N(w,q,t) = \lceil \log_2(N_1(w,q,t) + 1) - 1 \rceil$$
(13)

then Z(d, n) > 0 when  $s \ge N(w, q, t)$ .

**Remark 2.** Note that w and t are variable in the theorem 1 and s is variable in the theorem 2, although the d and n have the same forms in the two theorems, the two theorems have different meanings. Z(d,n) < 0 in the theorem 1 and Z(d,n) > 0 in the theorem 2. In fact,  $wt(X(d,n))/2^n \to 1$  when  $w \to \infty$  in the theorem 1 and  $wt(X(d, n))/2^n \to 0$  when  $s \to \infty$  in the theorem 2. The two theorems reveal the relationships among w, t, s, q and Z(d, n). The two theorems gave some unbalanced Elementary Symmetric Boolean Functions from two aspects.

4 When 
$$n = 2^{t+w+1}(1+2^1+\dots+2^s)+2^tq$$

In the section, we discuss elementary symmetric Boolean functions with form  $n = 2^{t+w+1}(1+2^1+\cdots+2^s) + 2^t q$  and  $d = 2^{w+t}(1+2^1+\cdots+2^s)$ , where  $n \equiv 0 \mod 4$ . The following theorem 3 and theorem 4 are similar to the theorem 1 and theorem 2.

**Theorem 3.** Let q > 0, t > 1, w and s be nonnegative integers,  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq$  and  $d = 2^{w+t}(1+2^1+\cdots+2^s)$ . For fixed s and q, then Z(d,n) < 0 when  $w \ge N(s,q)$ , where N(s,q) is a nonnegative integer and is only relative to s and q.

*Proof.* Let  $S = 2^0 + 2^1 + \cdots + 2^s$ , then  $d = 2^{t+w}S$  and  $n = 2^{t+w+1}S + 2^tq$ . Assume that  $2^w \ge q+1$ , similarly to the proof of the theorem 1, then

$$Z(d,n) = \sum_{i=0}^{2^{w+t}S-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{w+t}-1} C_n^i + \sum_{i=2^{w+t}S+2^{w+t}}^n C_n^i$$
  
$$= 2 \cdot \left(\sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{t-1}q-1} C_n^i\right) - C_n^{2^{w+t}S+2^{t-1}q}$$
  
$$< 2 \cdot \left(\sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{t-1}q-1} C_n^i\right).$$
(14)

And note that the inequation (14) is the same as the equation (2), similarly to the proof of the theorem 1, if  $2^w \ge q + 1$  and  $2^w \ge N_1(S)/2 + q/2$  hold at the same time, then

$$Z(d,n) < 2 \cdot [2S(2^{w+t} - 2^{t-1}q) - (\frac{2S+2}{2S+1})^{2^{w+t} - 2^{t-1}q}] \cdot C_n^{2^{w+t}S + 2^tq - 2^{w+t} - 1} < 0.$$
(15)

 $N_1(S)$  is a positive real number and is only relative to S. If we let

$$N(s,q) = \lceil max\{ \log_2(q+1), \log_2(N_1(S)/2 + q/2) \} \rceil$$
(16)

then Z(d, n) < 0 when  $w \ge N(s, q)$ .

**Remark 3.** Similarly to the theorem 1, note that  $N(s,q) = \lceil \log_2(q+1) \rceil$ when  $q \ge N_1(S) - 2$ ,  $N(s,q) = \lceil \log_2(N_1(S)/2 + q/2) \rceil < \lceil \log_2(N_1(S)) \rceil$ when  $q < N_1(S) - 2$ , therefore, N(s,q) is only relative to s or q. And note aslo that t not be limited in the theorem 1.

**Theorem 4.** The conditions are the same as the theorem 3. For fixed w, q and t, if  $2^w \ge q+1$ , then Z(d,n) > 0 when  $s \ge N(w,q,t)$ , where N(w,q,t) is a nongetative integer and is only relative to w, q and t.

*Proof.* We still let  $S = 2^0 + 2^1 + \dots + 2^s$ , then  $d = 2^{t+w}S$  and  $n = 2^{t+w+1}S + 2^t q$ . From the proof of the theorem 3, we have

$$Z(d,n) = 2 \cdot \left(\sum_{i=0}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - \sum_{i=2^{w+t}S}^{2^{w+t}S+2^{t-1}q-1} C_n^i\right) - C_n^{2^{w+t}S+2^{t-1}q}$$
  
>  $2 \cdot \sum_{i=2^{w+t}S-2^{w+t}-1}^{2^{w+t}S+2^tq-2^{w+t}-1} C_n^i - 2(2^{t-1}q+1)C_n^{2^{w+t}S+2^{t-1}q}$  (17)  
 $\triangleq F.$ 

Similarly to the proof of the theorem 2, let  $\prod_{j=0}^{i-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q + 2^{-j}) = 1$  when i = 0, then

$$\frac{F \cdot (2^{w+t}S + 2^{t-1}q)!(2^{w+t}S + 2^{w+t} + 2^{t-1}q + 2)!}{2 \cdot (2^{w+t+1}S + 2^{t}q)!} = \sum_{i=0}^{2^{t-1}q+1} [\prod_{j=0}^{2^{w+t}+1-i} (2^{w+t}S + 2^{t-1}q - j)] \\
\cdot \prod_{j=0}^{i-1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q + 2 - j)] \\
- (2^{t-1}q + 1) \prod_{i=0}^{2^{w+t}+1} (2^{w+t}S + 2^{w+t} + 2^{t-1}q + 2 - i) \\
= (2^{t-1}q + 2) \cdot [S^{2^{w+t}+2} + o(S^{2^{w+t}+2})] \\
- (2^{t-1}q + 1) \cdot [S^{2^{w+t}+2} + o(S^{2^{w+t}+2})] \\
= S^{2^{w+t}+2} - o(S^{2^{w+t}+2}).$$
(18)

Since lemma 3, then there exists a positive real number  $N_1(w, q, t)$  which is only relative to w, q and t.  $S^{2^{w+t}+2} - o(S^{2^{w+t}+2}) > 0$  when  $S = 1 + 2^1 + 2^2 + \cdots + 2^s > N_1(w, q, t)$ . If we let  $N(w, q, t) = \lceil \log_2(N_1(w, q, t) + 1) - 1 \rceil$ , then Z(d, n) > 0 when  $s \ge N(w, q, t)$ .

**Remark 4.** For  $n = 2^{t+w+1}(1+2^1+\cdots+2^s) + 2^tq + m$ ,  $d = 2^{w+t}(1+2^1+\cdots+2^s)$ ,  $m \equiv pmod4$ , similarly to the proof of the case  $n = 2^{t+w+1}(1+2^t+\cdots+2^s)$ 

 $2^1 + \cdots + 2^s) + 2^t q - 1$  if p is odd, similarly to the proof of the case  $n = 2^{t+w+1}(1+2^1+\cdots+2^s) + 2^t q$  if p is even, we can get similar results to the foregoing 4 theorems.

### 5 Conclusion

The paper considers Unbalanced elementary symmetric Boolean functions X(d, n) with special form  $d = 2^{t+w}(1+2^1+\cdots+2^s)$ ,  $n = 2^{t+w+1}(1+2^1+\cdots+2^s)+2^tq+m$ . For fixed s, q, or fixed w, q, t, we present a majority of X(d, n) are not balanced. Our results include many X(d, n) that  $d \equiv 0, 1, 2, 3mod4$ , which is supplement of only case  $d \equiv 3mod4$ . Our results are also parts of the conjecture that  $X(2^t, 2^{t+1}l - 1)$  is only nonlinear balanced elementary symmetric Boolean function. For others special forms X(d, n), we also can give many similar results in the same method.

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