On the Existence and Constructions of Vectorial Boolean Bent Functions*

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Abstract

Recently, obtaining vectorial Boolean bent functions of the form $Tr_m^n(P(x))$, where $P(x) \in \mathbb{F}_{2^n}[x]$, from Boolean bent functions of the form $Tr_1^n(P(x))$ has attracted several attentions and some open problems about this issue were proposed. This paper first provides three constructions of vectorial Boolean bent functions in the form $Tr_m^n(P(x))$, where two of them imply answers to two open problems proposed by E.Pasalic et al. and A.Muratović-Ribić et al. respectively. And by analyzing known types of Boolean bent functions of the form $Tr_1^n(P(x))$, the existence and constructions of several types of vectorial Boolean bent functions in the form $Tr_m^n(P(x))$ are obtained.

1 Introduction

Bent functions were initially introduced by Rothaus in [33], where it was shown that n-variable Boolean bent functions exist if and only if n is even. The term bent for vectorial Boolean functions, which is an extension of Boolean bent functions, was first considered by Nyberg in [29], where it was shown that bent (n, m)-functions (ie., vectorial Boolean functions from $\mathbb{F}_{2^n}$ to $\mathbb{F}_{2^m}$) exist if and only if n is even and $n \geq 2m$. Vectorial Boolean bent functions are also called perfect nonlinear functions [7], and there are several equivalent descriptions of vectorial Boolean bent functions, such as maximally nonlinear vectorial Boolean functions, vectorial Boolean functions having flat Walsh spectra, vectorial Boolean functions whose non-zero components are Boolean bent functions, vectorial Boolean functions having the minimum differential uniformity, etc. Because of possessing the maximal nonlinearity and the minimum differential uniformity, vectorial Boolean bent functions have the optimum resistance to linear cryptanalysis [16, 17, 22, 23] and differential cryptanalysis [1, 2, 3, 35], Thus, the constructions of vectorial Boolean bent functions are theoretical important and have great practical significance.

The constructions of vectorial Boolean bent functions are divided into two categories: primary constructions and secondary constructions. Primary constructions are also called

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direct constructions, and secondary constructions lead to vectorial Boolean bent functions with using known vectorial Boolean bent functions, which are also called indirect constructions. Among the constructions of vectorial Boolean bent functions, primary constructions hold a key status. Most of primary constructions of vectorial Boolean bent functions stem from the Maiorana-McFarland method [12, 24] or the Partial Spread method [12]. The strict Maiorana-McFarland constructions, the extended Maiorana-McFarland constructions and the general Maiorana-McFarland constructions are the three constructions [7, 30, 31, 34] of vectorial Boolean bent functions in the light of Maiorana-McFarland constructions of Boolean bent functions. By the Partial Spread method, the $\mathcal{PS}_{ap}$ constructions [7] and a Partial Spread construction [9] of vectorial Boolean bent functions were presented. Particularly, by study new connections between vectorial Boolean bent functions and the hyperovals of the projective plane, S. Mesnager [26] introduced a new primary construction of bent $(n, \frac{n}{2})$-functions from o-polynomials, named class $\mathcal{H}$ of vectorial functions. For a comprehensive understanding of primary constructions of vectorial Boolean bent functions, we send the readers to [26]. In addition, except the well known Direct Sum Construction, only two secondary constructions of vectorial Boolean bent functions (Proposition 9.5 and Proposition 9.6 in [7]) can be found in literatures, which are the generalizations of two secondary constructions of Boolean bent functions in [5] and [6] respectively. In addition, Proposition 9.6 in [7] includes Direct Sum Construction as a special case.

Recently, the new primary constructions that obtaining vectorial Boolean bent functions of the form $Tr_m^n(P(x))$ from Boolean bent functions of the form $Tr_1^n(P(x))$ have attracted a lot of attentions, where $P(x) \in \mathbb{F}_{2^n}[x]$.

The Boolean bent functions of the form $Tr_1^n(ax^d)$ are known as monomial bent functions, where d (understood modulo $2^n - 1$) is named a bent exponent that an integer such that $Tr_1^n(ax^d)$ is bent for some a. So far, there are five types of monomial bent functions [25], which are named by respective type of bent exponent, and listed in Table 1. In [32], when $Tr_1^n(ax^d)$ is bent, it was shown that the vectorial monomial function $Tr_m^n(ax^d)$ is bent if x^d is a permutation over $\mathbb{F}_{2^m}$. Following this conclusion, three classes of monomial bent functions, Kasami case, Leander case and Canteaut-Charpin-Kyureghyan case, were analyzed, and as a result, some classes of vectorial monomial bent functions were constructed [32]. However, it remains to be an open problem to judge whether the condition that x^d is a permutation over $\mathbb{F}_{2^m}$ is or not a necessary condition, which also affects whether the constructions of vectorial monomial bent functions in [32] are optimal. In [15], two classes of vectorial monomial bent functions of the form $Tr_{\frac{n}{2}}^n(ax^d)$ are constructed based on monomial bent functions in Gold case and Kasimi case. For a monomial bent function $Tr_1^n(ax^d)$ that belongs to Gold case or Kasimi case, it was shown [15] that $Tr_{\frac{n}{2}}^n(ax^d)$ is bent if $\gcd(d, 2^n - 1) \mid (2^{\frac{n}{2}} + 1)$ and $a \notin \{x^{\gcd(d,2^n-1)} : x \in \mathbb{F}_{2^n}\}$. However, whether the condition that $\gcd(d, 2^n - 1) \mid (2^{\frac{n}{2}} + 1)$ is necessary or not is unclear in [15]. In [28], it was also proved that there is no vectorial monomial bent function of the form $Tr_{\frac{n}{2}}^n(ax^d)$ as in Dillon case.

The bent properties of binomial Boolean functions of the form $Tr_1^n(a_1x^{d_1} + a_2x^{d_2})$ was studied in [14], where d_i, $i = 1, 2$, are Niho exponents i.e. the restriction of x^{d_i} on $\mathbb{F}_{2^k}$ is linear. Soon afterwards, a construction of Boolean bent functions with 2^r Niho Exponents was introduced [19], and the dual of it is computed [8]. In [21], an equivalent form of the construction in [19] were present. In [28], it was shown that the vectorial Boolean function $Tr_{\frac{n}{2}}^n(a_1x^{d_1} + a_2x^{d_2})$ is bent for a_1, a_2, d_1, d_2 as specified in [14]. Based on Boolean bent function of the form $Tr_1^n(\sum_{i=1}^r a_ix^{d_i})$, the construction of vectorial Boolean bent functions of

Table 1: Monomial Bent Functions of the Form $Tr_1^n(ax^d)$

Case	Exponent d	Condition 1	[1]Conditions 2	References
Gold	2^s+1	$s \in \mathbb{N}$	$a \notin \{x^d : x \in \mathbb{F}_{2^n}\}$	[20]
Dillon	$l(2^{\frac{n}{2}}-1)$	$\gcd(l, 2^{\frac{n}{2}}+1)=1$ (or $l=1$)	[2]$K(a)=-1, \mathbb{F}^*_{2^{\frac{n}{2}}}$ (or $K(N^n_{\frac{n}{2}}(a))=-1, \mathbb{F}^*_{2^n}$)	[10, 20] [18]
Kasami	$2^{2s}-2^s+1$	$\gcd(3,n)=1$, $\gcd(s,n)=1$	$a \notin \{x^3 : x \in \mathbb{F}_{2^n}\}$	[13, 20]
Leander	$(2^s+1)^2$	$n=4s$, s odd	$a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$	[11, 20], Theorem 13 this paper
Canteaut-Charpin-Kyureghyan	$2^{2s}+2^s+1$	$n=6s$, $s>1$ integer	$a \in \{x^d : x \in \mathbb{F}^*_{2^n}\}$ $\cdot\{\rho : Tr_s^{3s}(\rho)=0, \rho \in \mathbb{F}^*_{2^{3s}}\}$	[4, 11], Theorem 15 this paper

[1] Necessary and sufficient conditions such that $Tr_1^n(ax^d)$ is bent.

[2] Kloosterman sums $K(a)=\sum_{x\in\mathbb{F}^*_{2^{\frac{n}{2}}}}(-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+ax)}$.

the form $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ was studied in [28]. For a Boolean bent function $Tr_1^n(\sum_{i=1}^r a_i x^{d_i})$, where $d_i = d_1 + v_i(2^m-1)$, $i = 2, 3, \cdots, r$, and v_i are nonnegative integers, it was shown that [28], the vectorial Boolean function $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent if x^{d_1} is a permutation over $\mathbb{F}_{2^m}$. And one of the open problems left in [28] is to find a similar result to the above conclusion in the case when x^{d_1} is not a permutation over $\mathbb{F}_{2^m}$.

This paper is devoted to the existence and constructions of vectorial Boolean bent functions in the form $Tr_m^n(P(x))$ based on Boolean bent functions of the form $Tr_1^n(P(x))$, where $P(x) \in \mathbb{F}_{2^n}[x]$. We firstly give three constructions of vectorial Boolean bent functions, where two of them provide answers to the two open problems proposed by E.Pasalic et al. [32] and A.Muratović-Ribić et al. [28] respectively. Moreover, by the techniques presented in section 3, we analyze the bent properties of several types of vectorial Boolean functions of the form $Tr_m^n(P(x))$. It is mainly obtained that the existence and constructions of vectorial monomial bent functions of the form $Tr_m^n(ax^d)$ by using the known five types of monomial bent functions in Table 1. In addition, a construction of vectorial Boolean bent functions in the form $Tr_m^n(\sum_{i=1}^{2^s-1} ax^{(i\cdot 2^{\frac{n}{2}-s}+1)(2^{\frac{n}{2}}-1)+1})$ is presented, where $(i\cdot 2^{\frac{n}{2}-s}+1)(2^{\frac{n}{2}}-1)+1$, $i = 1, 2, \cdots, 2^s-1$, are Niho exponents and $\gcd(s, 2^{\frac{n}{2}})=1$.

The rest of this paper is organized as follows. Section 2 provides some preliminaries for the description of the paper. Section 3 presents three constructions of vectorial Boolean bent functions, and gives answers to two open problems proposed by E.Pasalic et al. and A.Muratović-Ribić et al. respectively. Section 4 analyzes the bent properties of several types of vectorial Boolean functions. Section 5 concludes this paper.

2 Preliminaries

Throughout this paper, let $\mathbb{F}_{2^n}$ denote the Galois field $GF(2^n)$, m and n be two positive integers, α be a primitive element of $\mathbb{F}_{2^n}$, $t = 2^{n-m}+2^{n-2m}+\cdots+2^m+1$ if $m \mid n$, and let the Kloosterman sums $K(a)=\sum_{x\in\mathbb{F}^*_{2^{\frac{n}{2}}}}(-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+ax)}$.

For $m \mid n$, the trace function $Tr_m^n : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$, is defined as

$$Tr_m^n(x) = x + x^{2^m} + x^{2^{2m}} + \cdots + x^{2^{(n/m-1)m}}, x \in \mathbb{F}_{2^n}.$$

In particular, Tr_1^n is called the absolute trace function over $\mathbb{F}_{2^n}$. Note that the number of the inputs x satisfying $Tr_m^n(x) = \theta$ is 2^{n-m} for $\forall\ \theta \in \mathbb{F}_{2^m}$, and the trace function has the well known properties that $Tr_1^n(x) = Tr_1^m \circ Tr_m^n(x)$ and $Tr_m^n(x) = Tr_m^n(x^2)$.

For $m \mid n$, the norm function $N_m^n : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$, is defined as

$$N_m^n(x) = x \cdot x^{2^m} \cdot x^{2^{2m}} \cdots x^{2^{(n/m-1)m}}, x \in \mathbb{F}_{2^n}.$$

A Boolean function $f : \mathbb{F}_{2^n} \to \mathbb{F}_2$ can be uniquely represented as the univariate polynomial representation

$$f(x) = \sum_{i=0}^{2^n-1} \sigma_i x^i, \sigma_i \in \mathbb{F}_{2^n},$$

where $f^2(x) \equiv f(x)(\mathrm{mod} x^{2^n} - x)$. The Boolean function $f : \mathbb{F}_{2^n} \to \mathbb{F}_2$ can also be represented as a non-unique way [27]

$$f = Tr_1^n(P(x)), P(x) \in \mathbb{F}_{2^n}[x].$$

The linear Boolean functions $\varphi : \mathbb{F}_{2^n} \to \mathbb{F}_2$ are the functions

$$\varphi(x) = Tr_1^n(ax), a \in \mathbb{F}_{2^n}.$$

The affine Boolean functions on $\mathbb{F}_{2^n}$ are the functions $\varphi(x) + \delta$, where $\delta \in \mathbb{F}_2$.

A mapping $F : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$ is referred to as vectorial Boolean function, which is also known as $(n, m)-$function, multiple output Boolean function or S-box, and can be uniquely represented in

$$F(x) = \sum_{i=0}^{2^n-1} \tau_i x^i, \tau_i \in \mathbb{F}_{2^n},$$

which is called the univariate polynomial representation of vectorial Boolean functions. If $m \mid n$, the vectorial Boolean function $F : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$ can also be represented in a non-unique way [7]

$$F = Tr_m^n(P(x)), P(x) \in \mathbb{F}_{2^n}[x].$$

If $m \mid n$, the linear vectorial Boolean functions $\phi : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$ are the functions

$$\phi(x) = Tr_m^n(ax), a \in \mathbb{F}_{2^n}.$$

If $m \mid n$, the affine vectorial Boolean functions on $\mathbb{F}_{2^n}$ are the functions $\phi(x) + \eta$, where $\eta \in \mathbb{F}_{2^m}$.

The Walsh transform of a Boolean function $f : \mathbb{F}_{2^n} \to \mathbb{F}_2$ is denoted by $W_f(\omega)$ and defined as

$$W_f(\omega) = \Sigma_{x \in \mathbb{F}_{2^n}} (-1)^{f(x) + Tr_1^n(\omega x)}, \forall\ \omega \in \mathbb{F}_{2^n}.$$

The extended Walsh transform of a vectorial Boolean function $F : \mathbb{F}_{2^n} \to \mathbb{F}_{2^m}$ is denoted by $W_F(\omega, \lambda)$ and defined as

$$W_F(\omega, \lambda) = \Sigma_{x \in \mathbb{F}_{2^n}} (-1)^{Tr_1^m(\lambda F(x)) + Tr_1^n(\omega x)}, \forall\ \omega \in \mathbb{F}_{2^n}, \forall\ \lambda \in \mathbb{F}_{2^m}^*.$$

Among the equivalent definitions of the term bent for Boolean functions and vectorial Boolean functions, we recall the following definitions.

Definition 1. *For even n, a Boolean function $f : \mathbb{F}_{2^n} \to \mathbb{F}_2$ is called bent if and only if $W_f(\omega) = \pm 2^{\frac{n}{2}}$ holds for all $\omega \in \mathbb{F}_{2^n}$.*

Definition 2. *For even n, a vectorial Boolean function $F : \mathbb{F}_{2^n} \to \mathbb{F}_m$ is called bent if and only if $W_F(\omega, \lambda) = \pm 2^{\frac{n}{2}}$ holds for all $\omega \in \mathbb{F}_{2^n}$, $\lambda \in \mathbb{F}_{2^m}^*$.*

3 Vectorial Boolean bent functions from Boolean bent function in trace form

In this section, using Boolean bent functions in trace form, three constructions of vectorial Boolean bent functions are given. The first two constructions, Theorem 1 and Theorem 3, imply answers to two open problems proposed by E.Pasalic et al. [32] and A.Muratović-Ribić et al. [28] respectively.

Before the discussions, two useful lemmas are presented below, which will be used in the sequel frequently. The two lemmas are more or less known in basic Algebra and usually used in literatures, however, it is difficult to find a explicit precise reference. For the convenience of the readers, we include the proofs here.

Lemma 1. *Let $G = \{x^d : x \in \mathbb{F}_{2^n}^*\}$. Then $G = \langle \alpha^d \rangle = \langle \alpha^{\gcd(d, 2^n-1)} \rangle$.*

Proof. Note the fact that $G = \{x^d : x \in \mathbb{F}_{2^n}^*\}$ is a cyclic subgroup of $\mathbb{F}_{2^n}^*$. For $\forall\ x \in \mathbb{F}_{2^n}^*$, there is an integer u such that $x = \alpha^u$, where $1 \leq u \leq 2^n - 1$. Then $G \subseteq \langle \alpha^{ud} \rangle \subseteq \langle \alpha^d \rangle$. For $\forall\ \alpha^{ud} \in \langle \alpha^d \rangle$, according to $\alpha^u \in \mathbb{F}_{2^n}^*$, $\alpha^{ud} \in G$, thus, $\langle \alpha^d \rangle \subseteq G$. Therefore, $G = \langle \alpha^d \rangle$.

Due to $|\langle \alpha^{\gcd(d,2^n-1)} \rangle| = \frac{2^n-1}{\gcd(d,2^n-1)} = |\langle \alpha^d \rangle|$ and there is a unique cyclic subgroup of order $\frac{2^n-1}{\gcd(d,2^n-1)}$ in $\mathbb{F}_{2^n}^*$, we have $\langle \alpha^{\gcd(d,2^n-1)} \rangle = \langle \alpha^d \rangle$. □

Lemma 2. *Let $m \mid n$ and $G = \{x^d : x \in \mathbb{F}_{2^n}^*\}$. Then $\mathbb{F}_{2^m}^* \subseteq G$ if and only if $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$.*

Proof. By Lemma 1, it is known that G is a cyclic subgroup of $\mathbb{F}_{2^n}^*$ and the order of G is $\frac{2^n-1}{\gcd(d,2^n-1)}$.

Due to $m \mid n$, we have $(2^m - 1) \mid (2^n - 1)$, thus $\mathbb{F}_{2^m}^*$ is a cyclic subgroup of order $2^m - 1$ in $\mathbb{F}_{2^n}^*$. If $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$, then there is a cyclic subgroup of order $2^m - 1$ in G. According to $G \subseteq \mathbb{F}_{2^n}^*$ and there is a unique cyclic subgroup of order $2^m - 1$ in $\mathbb{F}_{2^n}^*$, we have that G and $\mathbb{F}_{2^n}^*$ have the same cyclic subgroup $\mathbb{F}_{2^m}^*$. Thus, $\mathbb{F}_{2^m}^* \subseteq G$.

The necessity is obvious. Hence the conclusion of Lemma 2 holds. □

The following Theorem 1 gives a sufficient condition for the vectorial Boolean functions of the form $Tr_m^n(ax^d)$ to be bent, which includes E.Pasalic et al's Theorem 1 in [32] as a special case and implies the answer to one open problem in [32] (named Open Problem 1 in this paper).

Theorem 1. *Let n be even and $m \mid n$, and let $f(x) = Tr_1^n(ax^d)$ be a Boolean bent function. If $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$, then the vectorial Boolean function $F(x) = Tr_m^n(ax^d)$ is bent.*

Proof. Let $G = \{x^d : x \in \mathbb{F}_{2^n}^*\}$. If $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$, according to Lemma 2, then $\mathbb{F}_{2^m}^* \subseteq G$. Thus, for $\forall\ \lambda \in \mathbb{F}_{2^m}^*$, $\exists\ \beta \in \mathbb{F}_{2^n}^*$ such that $\lambda = \beta^d$. Therefore, for $\forall\ \omega \in \mathbb{F}_{2^n}, \forall\ \lambda \in \mathbb{F}_{2^m}^*$, we have

$$\begin{aligned}
W_F(\omega,\lambda) &= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^m(\lambda F(x))+Tr_1^n(\omega x)}\\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^m(\lambda Tr_m^n(ax^d))+Tr_1^n(\omega x)}\\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(a\lambda x^d)+Tr_1^n(\omega x)}\\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(a\beta^d x^d)+Tr_1^n(\omega x)}\\
&= \Sigma_{y\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(ay^d)+Tr_1^n(\omega\beta^{-1}y)}\\
&= W_f(\omega\beta^{-1})\\
&= \pm 2^{\frac{n}{2}}
\end{aligned}$$

By definition 2, $F(x)$ is bent and the conclusion holds. □

In [32], it is proved that, if the Boolean function $Tr_1^n(ax^d)$ is bent, then x^d is a permutation over $\mathbb{F}_{2^m}$ is a sufficient condition for the vectorial Boolean function $Tr_m^n(ax^d)$ to be bent, and an open problem is left as below.

Open Problem 1 ([32])**.** *Assume that the Boolean function $Tr_1^n(ax^d)$ is bent, prove or disprove that the condition x^d is a permutation over $\mathbb{F}_{2^m}$ is necessary for the vectorial Boolean function $Tr_m^n(ax^d)$ to be bent.*

In order to answer Open problem 1, we first present Theorem 2 and Remark 1.

Theorem 2. *Let $m \mid n$. If x^d is a permutation over $\mathbb{F}_{2^m}$, then $(2^m-1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$.*

Proof. Due to $m \mid n$, we have $(2^m-1) \mid (2^n-1)$. If x^d is a permutation over $\mathbb{F}_{2^m}$, then $\gcd(d, 2^m-1) = 1$, and thus $\gcd(\gcd(d, 2^n-1), 2^m-1) = 1$. Therefore, $(2^m-1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$. □

Remark 1. Theorem 2 is not vice versa. For example, let $m = 2, n = 6, d = 3$, then $(2^m-1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$. However, $\gcd(d, 2^m-1) = 3 \neq 1$. Thus, x^d is not a permutation over $\mathbb{F}_{2^m}$.

Following from Theorem 1, Theorem 2 and Remark 1, the answer to Open Problem 1 can be made, i.e., the condition that x^d is a permutation over $\mathbb{F}_{2^m}$ is not necessary.

Remark 2. By Remark 1 it is known that the sufficient condition of Theorem 1 is weaker than that in E.Pasalic et al's Theorem 1 of [32], i.e. our condition is closer to the necessary condition. Note that, according to Corollary 1 and Corollary 2, the condition $(2^m-1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$ is necessary in Gold case for $m = \frac{n}{2}$ and in Leander case for any $m \mid n$.

The following theorem provides a sufficient condition for the vectorial Boolean functions of the form $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ to be bent, which implies an answer to A.Muratović-Ribić et al.'s Open Problem 1 in [28] (named Open Problem 2 in this paper).

Theorem 3. *Let n be even, $m \mid n$, and $d_i = d_1 + v_i\frac{2^n-1}{\gcd(d_1,2^n-1)}$, where $i = 2, 3, \cdots, r$ and v_i are nonzero integers, and let $f(x) = Tr_1^n(\sum_{i=1}^r a_i x^{d_i})$ be a Boolean bent function. If $(2^m-1) \mid \frac{2^n-1}{\gcd(d_1^2,2^n-1)}$, then the vectorial Boolean function $F(x) = Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent.*

Proof. Let $G_1 = \{x^{d_1} : x \in \mathbb{F}_{2^n}^*\}$ and $G_2 = \{x^{d_1^2} : x \in \mathbb{F}_{2^n}^*\}$. If $(2^m-1) \mid \frac{2^n-1}{\gcd(d_1^2, 2^n-1)}$, according to Lemma 2, then $\mathbb{F}_{2^m}^* \subseteq G_2$. Thus, for $\forall\ \lambda \in \mathbb{F}_{2^m}^*$, $\exists\ \gamma \in \mathbb{F}_{2^n}^*$ such that $\lambda = \gamma^{d_1^2}$ and $\gamma^{d_1} \in G_1$. For nonnegative integers v_i and $i = 2, \cdots, r$, according to the order of G_1 is $\frac{2^n-1}{\gcd(d_1, 2^n-1)}$, if $d_i = d_1 + v_i\frac{2^n-1}{\gcd(d_1, 2^n-1)}$, then $\gamma^{d_1 d_i} = \gamma^{d_1^2}$.

Therefore, for $\forall\ \omega \in \mathbb{F}_{2^n}, \forall\ \lambda \in \mathbb{F}_{2^m}^*$, we have

$$\begin{aligned}
W_F(\omega, \lambda) &= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^m(\lambda F(x))+Tr_1^n(\omega x)} \\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^m(\lambda Tr_m^n(\sum_{i=1}^r a_i x^{d_i}))+Tr_1^n(\omega x)} \\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(\lambda \sum_{i=1}^r a_i x^{d_i})+Tr_1^n(\omega x)} \\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(\gamma^{d_1^2} \sum_{i=1}^r a_i x^{d_i})+Tr_1^n(\omega x)} \\
&= \Sigma_{x\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(\sum_{i=1}^r a_i \gamma^{d_1 d_i} x^{d_i})+Tr_1^n(\omega x)} \\
&= \Sigma_{y\in\mathbb{F}_{2^n}}(-1)^{Tr_1^n(\sum_{i=1}^r a_i y^{d_i})+Tr_1^n(\omega\gamma^{-d_1} y)} \\
&= W_f(\omega\gamma^{-d_1}) \\
&= \pm 2^{\frac{n}{2}}
\end{aligned}$$

By definition 2 we know that $F(x)$ is bent. □

In [28], under the condition that x^{d_1} is a permutation over $\mathbb{F}_{2^m}$, A.Muratović-Ribić et al. derived that if the Boolean function $Tr_1^n(\sum_{i=1}^r a_i x^{d_i})$ is bent then the vectorial Boolean function $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent, where $d_i = d_1 + v_i(2^m-1)$ and v_i are nonnegative integers, $i = 2, \cdots, r$. And an open problem about finding a similar conclusion if it is not to meet the condition that x^{d_1} is a permutation over $\mathbb{F}_{2^m}$ is left.

Open Problem 2 ([28])**.** *Let $n \geq 4$ be an even positive integer, and let $m \leq \frac{n}{2}$ and $m \mid n$. Let x^{d_1} be a permutation of $\mathbb{F}_{2^n}$, and let $f(x) = Tr_1^n(\sum_{i=1}^r a_i x^{d_i})$ be a Boolean bent function, where $m \mid \frac{n}{2}$ and $d_i = d_1 + v_i(2^m-1)$ for $i = 2, \cdots, r$ and some integers $v_i \geq 0$. Then, the function $F(x) = Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is a vectorial bent function.*

Is there a similar result to the above for the functions of the form $f(x) = Tr_1^n(\sum_{i=1}^r x^{d_i})$, if x^{d_1} is not a permutation over $\mathbb{F}_{2^m}$?

Note that, Open Problem 2 asks for a similar result to the authors' conclusion (Theorem 1 in [28]) when x^{d_1} is not a permutation over $\mathbb{F}_{2^m}$, so, there may be several different answers. Here, we give one answer. Before giving our answer to Open problem 2, we present Theorem 4 and Remark 3.

Theorem 4. *Let $m \mid n$. If x^d is a permutation over $\mathbb{F}_{2^m}$, then $(2^m-1) \mid \frac{2^n-1}{\gcd(d^2, 2^n-1)}$.*

Proof. Due to $m \mid n$, we have $(2^m-1) \mid (2^n-1)$. If x^d is a permutation over $\mathbb{F}_{2^m}$, then $\gcd(d, 2^m-1) = 1$, and thus $\gcd(\gcd(d^2, 2^n-1), 2^m-1) = 1$. Therefore, $(2^m-1) \mid \frac{2^n-1}{\gcd(d^2, 2^n-1)}$. □

Remark 3. Theorem 4 is not vice versa. For example, let $m = 2, n = 18, d = 3$, then $(2^m-1) \mid \frac{2^n-1}{\gcd(d^2, 2^n-1)}$. However, $\gcd(d, 2^m-1) = 3 \neq 1$. Thus, x^d is not a permutation over $\mathbb{F}_{2^m}$.

By Theorem 3, Theorem 4 and Remark 3, we get the following theorem which is an answer to Open Problem 2.

Theorem 5. *Let $m \mid n$, $v_i = \frac{t \cdot u_i}{\gcd(d_1, 2^n-1)}$, $d_i = d_1 + v_i(2^m - 1)$, where $i = 2, \cdots, r$ and u_i are nonzero integers, and let $f(x) = Tr_1^n(\sum_{i=1}^r a_i x^{d_i})$ be a Boolean bent function. If $(2^m - 1) \mid \frac{2^n-1}{\gcd(d_1^2, 2^n-1)}$, then the vectorial Boolean function $F(x) = Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent.*

In the above, we have given two sufficient conditions for the vectorial Boolean function $F(x)$ to be bent. However, neither of the sufficient conditions has been proved necessary or not, although they seem to be very close to be. Thus, up to now, the techniques to analyze the bent property of vectorial Boolean functions in the form $Tr_m^n(P(x))$ roundly are lack, where $P(x) \in \mathbb{F}_{2^n}[x]$.

In above discussions, we focus our attention on exponents, input dimension m and output dimension n, but the coefficients of $P(x)$ have not been considered. In the remainder of this section, we focus on the coefficient set of $P(x)$ such that $Tr_1^n(P(x))$ is bent, ie., considering a class of Boolean bent functions as a whole, to which $Tr_1^n(P(x))$ belongs.

When considering a class of Boolean bent function as a whole, a new construction of vectorial Boolean bent functions can be described as in the following theorem.

Theorem 6. *Let $m \mid n$ and*

$$C = \{(c_1, c_2, \cdots, c_r) : Tr_1^n(\sum_{i=1}^r c_i x^{d_i}) \text{ is bent, } (c_1, c_2, \cdots, c_r) \in \mathbb{F}_{2^n}^r\}.$$

The vectorial Boolean function $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent if and only if

$$(a_1, a_2, \cdots, a_r) \cdot \mathbb{F}_{2^m}^* \subseteq C.$$

Proof. For $\forall\ \omega \in \mathbb{F}_{2^n}, \forall\ \lambda \in \mathbb{F}_{2^m}^*$,

$$\begin{aligned} W_{Tr_m^n(\sum_{i=1}^r a_i x^{d_i})}(\omega, \lambda) &= \Sigma_{x \in \mathbb{F}_{2^n}} (-1)^{Tr_1^m(\lambda Tr_m^n(\sum_{i=1}^r a_i x^{d_i})) + Tr_1^n(\omega x)} \\ &= \Sigma_{x \in \mathbb{F}_{2^n}} (-1)^{Tr_1^n(\sum_{i=1}^r a_i \lambda x^{d_i}) + Tr_1^n(\omega x)} \\ &= W_{Tr_1^n(\sum_{i=1}^r a_i \lambda x^{d_i})}(\omega). \end{aligned}$$

Thus, $Tr_m^n(\sum_{i=1}^r a_i x^{d_i})$ is bent if and only if $Tr_1^n(\sum_{i=1}^r a_i \lambda x^{d_i})$ is bent for all $\lambda \in \mathbb{F}_{2^m}^*$. For all $\lambda \in \mathbb{F}_{2^m}^*$, $Tr_1^n(\sum_{i=1}^r a_i \lambda x^{d_i})$ is bent if and only if $(a_1, a_2, \cdots, a_r)\lambda \in C$. The condition that for all $\lambda \in \mathbb{F}_{2^m}^*$, $(a_1, a_2, \cdots, a_r)\lambda \in C$ is equivalent to

$$(a_1, a_2, \cdots, a_r) \cdot \mathbb{F}_{2^m}^* \subseteq C.$$

Consequently, the conclusion of the theorem holds. □

4 On the existence and constructions of vectorial Boolean functions

In this section, by the results and techniques presented in section 3, we analyze the bentness of several types of vectorial Boolean functions in the form $Tr_m^n(P(x))$, mainly about the existence and constructions of vectorial monomial bent functions of the form $Tr_m^n(ax^d)$.

Firstly, in order to obtain vectorial monomial bent functions, we analyze the known five types of monomial bent functions in Table 1 by Theorem 6, ie., considering the coefficient and by Theorem 1, ie., considering the exponent, integer s, input and output dimensions.

Theorem 7. *(Gold Case). Let n be even, $m \mid n$, $a \in \mathbb{F}_{2^n}$, $s \in \mathbb{N}$, and $d = 2^s + 1$.*

Then the vectorial Boolean function $Tr_m^n(ax^d)$ is bent if and only if $a \notin \{\langle \alpha^{\gcd(d,t)} \rangle, 0\}$. Moreover, there are $2^n - \frac{2^n-1}{\gcd(d,t)} - 1$ such vectorial monomial bent functions.

Proof. According to Theorem 2 in [20], the set of the coefficients such that $Tr_1^n(ax^d)$ is bent is

$$C = \{c : c \neq x^d, c, x \in \mathbb{F}_{2^n}\}.$$

Then, by Theorem 6, we obtain $Tr_m^n(ax^d)$ is bent if and only if

$$\begin{aligned}
& a \cdot \mathbb{F}_{2^m}^* \subseteq \{c : c \neq x^d, c, x \in \mathbb{F}_{2^n}\} \\
\Leftrightarrow\quad & a \cdot \mathbb{F}_{2^m}^* \bigcap \{x^d : x \in \mathbb{F}_{2^n}\} = \varnothing \\
\Leftrightarrow\quad & a \notin \{x^d : x \in \mathbb{F}_{2^n}\} \cdot \mathbb{F}_{2^m}^* \\
& \text{(Since Lemma 1)} \\
\Leftrightarrow\quad & a \notin \langle \alpha^{\gcd(t, \gcd(d, 2^n-1))} \rangle \\
& \text{(Since } t \mid (2^n - 1)) \\
\Leftrightarrow\quad & a \notin \langle \alpha^{\gcd(d,t)} \rangle
\end{aligned}$$

Thus, we have $Tr_m^n(ax^d)$ is bent if and only if $a \notin \{\langle \alpha^{\gcd(d,t)} \rangle, 0\}$.

Due to that $|\{\langle \alpha^{\gcd(d,t)} \rangle, 0\}| = \frac{2^n-1}{\gcd(d,t)} + 1$, there are $2^n - \frac{2^n-1}{\gcd(d,t)} - 1$ such vectorial monomial bent functions.

Consequently, the conclusion of the theorem holds. □

Before giving the other necessary and sufficient condition for vectorial Boolean functions of the form $Tr_{\frac{n}{2}}^n(ax^d)$ in Gold case to be bent, we give the following lemma.

Lemma 3. *Let the monomial Boolean function $f(x) = Tr_1^n(ax^d)$ be bent. Then $\gcd(d, 2^{\frac{n}{2}} - 1) = 1$ if and only if $\gcd(d, 2^{\frac{n}{2}} + 1) \neq 1$;*

Proof. ⇒) If $\gcd(d, 2^{\frac{n}{2}} - 1) = 1$, according to conclusion (1) of Lemma 1 in [20], then $W_f(0) = -2^{\frac{n}{2}}$. By conclusion (2) of Lemma 1 in [20], we have $\gcd(d, 2^{\frac{n}{2}} + 1) \neq 1$.

⇐) Since $f(x) = Tr_1^n(ax^d)$ is bent, then $W_f(0) = \pm 2^{\frac{n}{2}}$. If $\gcd(d, 2^{\frac{n}{2}} + 1) \neq 1$, according to conclusion (2) of Lemma 1 in [20], then $W_f(0) \neq 2^{\frac{n}{2}}$. Therefore, $W_f(0) = -2^{\frac{n}{2}}$. Thus, by conclusion (1) of Lemma 1 in [20], we have $\gcd(d, 2^{\frac{n}{2}} - 1) = 1$. □

The following corollary shows that the condition $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$ in Theorem 1 is necessary for $m = \frac{n}{2}$ in Gold case.

Corollary 1. *Let n be even, $s \in \mathbb{N}$, $a \in \mathbb{F}_{2^n}$, $d = 2^s + 1$, and let the monomial Boolean function $Tr_1^n(ax^d)$ be bent.*

Then the vectorial Boolean function $Tr_{\frac{n}{2}}^n(ax^d)$ is bent if and only if $\gcd(d, 2^n - 1) \mid (2^{\frac{n}{2}} + 1)$.

Proof. $\Rightarrow$) If $Tr^n_{\frac{n}{2}}(ax^d)$ is bent, according to Theorem 7, then $a \notin \{\langle \alpha^{\gcd(d,2^{\frac{n}{2}}+1)}\rangle, 0\}$, and then $\{\langle \alpha^{\gcd(d,2^{\frac{n}{2}}+1)}\rangle, 0\} \neq \mathbb{F}_{2^n}$, thus $\gcd(d, 2^{\frac{n}{2}}+1) \neq 1$. According to Lemma 3, we have $\gcd(d, 2^{\frac{n}{2}}-1) = 1$. Therefore, $\gcd(d, 2^n-1) = \gcd(d, 2^{\frac{n}{2}}+1)$. Thus, $\gcd(d, 2^n-1) \mid (2^{\frac{n}{2}}+1)$.

$\Leftarrow$) This follows from Theorem 1. □

Example 1. Let $s = 2$ and $n = 4$. Then $d = 2^s + 1 = 5$. $\langle \alpha^{\gcd(d,2^{\frac{n}{2}}+1)}\rangle = \langle \alpha^{\gcd(5,5)}\rangle = \langle \alpha^5\rangle \neq \mathbb{F}^*_{2^4}$. For $\forall\ a \in \mathbb{F}_{2^4}\backslash\{\langle\alpha^5\rangle, 0\} = \{\alpha, \alpha^2, \alpha^3, \alpha^4, \alpha^6, \alpha^7, \alpha^8, \alpha^9, \alpha^{11}, \alpha^{12}, \alpha^{13}, \alpha^{14}\}$, $Tr^4_2(ax^5)$ is a vectorial monomial bent function.

For the Dillion exponent $d = 2^{\frac{n}{2}} - 1$, it was shown that $Tr^n_1(ax^d)$ is bent if and only if $K(a) = -1$ in [12], where $a \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, and $Tr^n_1(ax^d)$ is bent if and only if $K(N^n_{\frac{n}{2}}(a)) = -1$ in [18], where $a \in \mathbb{F}^*_{2^n}$. In [10], it was shown that $Tr^n_1(ax^{l(2^{\frac{n}{2}}-1)})$ is bent if and only if $K(a) = -1$, where $\gcd(l, 2^{\frac{n}{2}} + 1) = 1$ and $a \in \mathbb{F}^*_{2^{\frac{n}{2}}}$. We recall Theorem 5 in [10] and Theorem 3 in [18] as the following theorem. Note that the Kloosterman sums $K(a) = \sum_{x\in\mathbb{F}^*_{2^{\frac{n}{2}}}} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+ax)}$ defined in this paper is the same as in [18] and different from [10].

Theorem 8. *[10, 18] Let n be even, l be an integer and $d = l(2^{\frac{n}{2}} - 1)$.*

If $\gcd(l, 2^{\frac{n}{2}} + 1) = 1$ and $a \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, then monomial Boolean function $Tr^n_1(ax^d)$ is bent if and only if

$$a \in \{\beta : K(\beta) = -1, \beta \in \mathbb{F}^*_{2^{\frac{n}{2}}}\}.$$

If $l = 1$ and $a \in \mathbb{F}^*_{2^n}$, then monomial Boolean function $Tr^n_1(ax^d)$ is bent if and only if

$$a \in \{\beta : K(N^n_{\frac{n}{2}}(\beta)) = -1, \beta \in \mathbb{F}^*_{2^n}\}.$$

Theorem 9. *(Dillion Case) Let n be even, l be an integer, $\gcd(l, 2^{\frac{n}{2}}+1) = 1$ (or $l = 1$), and $d = l(2^{\frac{n}{2}}-1)$, and let $a \in \{\beta : K(\beta) = -1, \beta \in \mathbb{F}^*_{2^{\frac{n}{2}}}\}$ (or $a \in \{\beta : K(N^n_{\frac{n}{2}}(\beta)) = -1, \beta \in \mathbb{F}^*_{2^n}\}$ accordingly).*

If $(2^m - 1) \mid (2^{\frac{n}{2}} + 1)$, then monomial vectorial Boolean function $Tr^n_m(ax^d)$ is bent.

Proof. Note that $\gcd(l, 2^{\frac{n}{2}}+1) = 1$, then $2^{\frac{n}{2}}+1 = \frac{2^n-1}{\gcd(d,2^n-1)}$, and then $(2^m-1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$. According to Theorem 1 and Theorem 8, the conclusion of the theorem holds. □

In [28], for $\gcd(l, 2^{\frac{n}{2}} + 1) = 1$ and $a \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, the nonexistence of vectorial monomial bent function of the form $Tr^n_{\frac{n}{2}}(ax^d)$ in Dillon case was proved. Here, for $l = 1$ and $a \in \mathbb{F}^*_{2^n}$, we prove that there is also no vectorial monomial bent function of the form $Tr^n_{\frac{n}{2}}(ax^d)$.

Lemma 4. *Let n be even. For $\forall\ \lambda \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, let $n_\lambda = N^n_{\frac{n}{2}}(a\lambda)$, where $a \in \mathbb{F}^*_{2^n}$. Then $\{n_\lambda : \lambda \in \mathbb{F}^*_{2^{\frac{n}{2}}}\} = \mathbb{F}^*_{2^{\frac{n}{2}}}$.*

Proof. Note the fact that, for $\forall\ \lambda_1, \lambda_2 \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, if $\lambda_1 \neq \lambda_2$, then $\lambda_1^{2^{\frac{n}{2}}+1} \neq \lambda_2^{2^{\frac{n}{2}}+1}$. Therefore, $\{\lambda^{\frac{n}{2}+1} : \lambda \in \mathbb{F}^*_{2^{\frac{n}{2}}}\} = \mathbb{F}^*_{2^{\frac{n}{2}}}$. Thanks to $n_\lambda = N^n_{\frac{n}{2}}(a\lambda) = a^{2^{\frac{n}{2}}+1}\lambda^{2^{\frac{n}{2}}+1}$ and $a^{2^{\frac{n}{2}}+1} \in \mathbb{F}^*_{2^{\frac{n}{2}}}$, we have $\{n_\lambda : \lambda \in \mathbb{F}^*_{2^{\frac{n}{2}}}\} = \mathbb{F}^*_{2^{\frac{n}{2}}}$. □

Theorem 10. *Let n be even, $a \in \mathbb{F}_{2^n}^*$ and $d = 2^{\frac{n}{2}} - 1$. Then there is no vectorial monomial bent function of the form $Tr_{\frac{n}{2}}^{n}(ax^d)$.*

Proof. Assume that such a vectorial monomial bent function $Tr_{\frac{n}{2}}^{n}(ax^d)$ exists. Let $n_\lambda = N_{\frac{n}{2}}^{n}(a\lambda)$. According to Theorem 3 in [18] and Theorem 6, we have $Tr_{\frac{n}{2}}^{n}(ax^d)$ is bent if and only if $K(n_\lambda) = \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+n_\lambda x)} = -1$, for all $\lambda \in \mathbb{F}_{2^{\frac{n}{2}}}^*$. Thus,

$$\begin{aligned}
& \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+n_\lambda x)} = -1, \text{ for all } \lambda \in \mathbb{F}_{2^{\frac{n}{2}}}^* \\
& \text{(according to Lemma 4)} \\
\Leftrightarrow\ & \sum_{x,y \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+y)} = -1 \\
\Leftrightarrow\ & \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1})} \sum_{y \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(y)} = -1 \\
\Leftrightarrow\ & (\sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x)})^2 = -1
\end{aligned}$$

This is impossible. This means that the assumption of the existence of vectorial monomial bent function $Tr_{\frac{n}{2}}^{n}(ax^d)$ is wrong, and hence the conclusion of the theorem holds. □

On the other hand, we give a new proof of Theorem 10 in [28].

Theorem 11. *[28] Let $n \geq 4$ be even, $a \in \mathbb{F}_{2^{\frac{n}{2}}}^*$, $d = l(2^{\frac{n}{2}} - 1)$ and $\gcd(l, \frac{n}{2}+1) = 1$. There is no vectorial monomial bent function of the form $Tr_{\frac{n}{2}}^{n}(ax^d)$.*

Proof. Assume that such a vectorial monomial bent function $Tr_{\frac{n}{2}}^{n}(ax^d)$ exists. According to Theorem 5 in [10] and Theorem 6, we have $Tr_{\frac{n}{2}}^{n}(ax^d)$ is bent if and only if $K(a\lambda) = \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+a\lambda x)} = -1$, for all $\lambda \in \mathbb{F}_{2^{\frac{n}{2}}}^*$. Thus,

$$\begin{aligned}
& \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+a\lambda x)} = -1, \text{ for all } \lambda \in \mathbb{F}_{2^{\frac{n}{2}}}^* \\
& \text{(Since } a \in \mathbb{F}_{2^{\frac{n}{2}}}^*\text{)} \\
\Leftrightarrow\ & \sum_{x,y \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1}+y)} = -1 \\
\Leftrightarrow\ & \sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x^{-1})} \sum_{y \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(y)} = -1 \\
\Leftrightarrow\ & (\sum_{x \in \mathbb{F}_{2^{\frac{n}{2}}}^*} (-1)^{Tr_1^{\frac{n}{2}}(x)})^2 = -1
\end{aligned}$$

This is impossible. This means that the assumption of the existence of vectorial monomial bent function $Tr^n_{\frac{n}{2}}(ax^d)$ is wrong, and hence the conclusion of the theorem holds. □

Theorem 12. *(Kasami Case). Let n be even, $m \mid n$, $a \in \mathbb{F}^*_{2^n}$, $\gcd(3,n) = 1$, $\gcd(s,n) = 1$, and $d = 2^{2s} - 2^s + 1$.*

Then the vectorial Boolean function $Tr^n_m(ax^d)$ is bent if and only if $a \notin \{\langle \alpha^{\gcd(3,t)} \rangle, 0\}$. Moreover, there are $2^n - \frac{2^n-1}{\gcd(3,t)} - 1$ such vectorial monomial bent functions.

Proof. According to Theorem 6 in [20] and Theorem 6, the proof is similar to Theorem 7. □

Example 2. Let $s = 3$ and $n = 10$. Then $\gcd(s,n) = \gcd(3,n) = 1, d = 2^{2s} - 2^s + 1 = 57$ and $\langle \alpha^{\gcd(3,2^{\frac{n}{2}}+1)} \rangle = \langle \alpha^3 \rangle \neq \mathbb{F}^*_{2^{10}}$. For $\forall$ $a \in \mathbb{F}_{2^n} \backslash \{\langle \alpha^3 \rangle, 0\}$, $Tr^{10}_5(ax^{57})$ is a vectorial monomial bent function.

In [20], it was proved there are monomial bent functions of the form $Tr^n_1(ax^d)$ with $d = (2^s + 1)^2$. Subsequently, the result was extended in [11]. It was shown that [11] the monomial Boolean function $Tr^{4s}_1(ax^{(2^s+1)^2})$ is bent if and only if there are $\rho \in \varepsilon\mathbb{F}^*_{2^s}$ and $\beta \in \mathbb{F}^*_{2^n}$ such that $a = \rho\beta^d$, where s is positive odd, $n = 4s, \varepsilon \in \mathbb{F}_4 \backslash \mathbb{F}_2$. Note the fact that the condition there are $\rho \in \varepsilon\mathbb{F}^*_{2^s}$ and $\beta \in \mathbb{F}^*_{2^n}$ such that $a = \rho\beta^d$ is equivalent to $a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \mathbb{F}^*_{2^s} \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$. According to Lemma 1, it holds that $\{x^d : x \in \mathbb{F}^*_{2^n}\} = \langle \alpha^{\gcd(d,2^n-1)} \rangle = \langle \alpha^{2^s+1} \rangle$. Thanks to $\mathbb{F}^*_{2^s} = \langle \alpha^{(2^s+1)(2^{2s}+1)} \rangle$, we have $\mathbb{F}^*_{2^s} \subset \{x^d : x \in \mathbb{F}^*_{2^n}\}$. Therefore, $\mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \mathbb{F}^*_{2^s} \cdot \{x^d | x \in \mathbb{F}^*_{2^n}\} = \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$. Thus, Theorem 4.8 in [11] can be described equivalently and more succinctly as the following theorem.

Theorem 13. *Let s be positive odd, $n = 4s$ and $d = (2^s + 1)^2$. The monomial Boolean function $Tr^n_1(ax^d)$ is bent if and only if $a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$.*

Lemma 5. *Let s be positive odd, $n = 4s$ and $d = (2^s+1)^2$. Then $\{x^d : x \in \mathbb{F}^*_{2^n}\} \cap \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\} = \varnothing$.*

Proof. Because s is odd, it holds that $\gcd(3, (2^s - 1)(2^{2s} + 1)) = 1$. For $\forall$ $\varepsilon \in \mathbb{F}_4 \backslash \mathbb{F}_2$, the order of ε is 3, thus $\varepsilon^{(2^s-1)(2^{2s}+1)} \neq 1$. According to Lemma 1, the order of $\{x^d : x \in \mathbb{F}^*_{2^n}\}$ is $(2^s - 1)(2^{2s} + 1)$. Then we have $\{x^d : x \in \mathbb{F}^*_{2^n}\} \cap \mathbb{F}_4 \backslash \mathbb{F}_2 = \varnothing$. Thus, $\{x^d : x \in \mathbb{F}^*_{2^n}\} \cap \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\} = \varnothing$. □

Theorem 14. *(Leander Case). Let s be positive odd, $n = 4s$, $m \mid n$ and $d = (2^s + 1)^2$.*

(1) Let $a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$. Then the vectorial Boolean function $Tr^n_m(ax^d)$ is bent if and only if $m \mid s$.

(2) Let m be odd. Then the vectorial Boolean function $Tr^n_m(ax^d)$ is bent if and only if $a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}$.

Proof. (1) $' \Leftarrow'$ If $m \mid s$, then $(2^m - 1) \mid (2^s - 1)(2^{2s} + 1)$. According to Theorem 1 and Theorem 13, $Tr^n_m(ax^d)$ is bent.

$' \Rightarrow'$ According to $Tr^n_m(ax^d)$ is bent and Theorem 6, we obtain

$$a \cdot \mathbb{F}^*_{2^m} \subseteq C = \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}^*_{2^n}\}.$$

If m is even, then $\mathbb{F}_4 \backslash \mathbb{F}_2 \subseteq \mathbb{F}^*_{2^m}$. Let $a = \varepsilon \cdot \tau$, where $\varepsilon \in \mathbb{F}_4 \backslash \mathbb{F}_2$ and $\tau \in \{x^d : x \in \mathbb{F}^*_{2^n}\}$. Then $\varepsilon^2 \in \mathbb{F}_4 \backslash \mathbb{F}_2 \subseteq \mathbb{F}^*_{2^m}$. Let $\lambda = \varepsilon^2$. Then $a\lambda = \varepsilon^3\tau = \tau \in \{x^d : x \in \mathbb{F}^*_{2^n}\}$. By Lemma 5, we have $a\lambda \notin \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}_{2^n}^*\}$. Thus, m can not be even, ie., m is odd. According to m is odd and $m \mid 4s$, we have $m \mid s$.

(2) $\Leftarrow$) According to m is odd and $m \mid 4s$, we have $m \mid s$, and then $(2^m - 1) \mid (2^s - 1)(2^{2s} + 1)$. According to Theorem 1 and Theorem 13, $Tr_m^n(ax^d)$ is bent.

$\Rightarrow$) The proof is trivial. □

By condition (1) of Theorem 14, the following corollary can be obtained, which implies, in Leander case, Theorem 1 is vice versa and there is no vectorial Boolean function of the form $Tr_{\frac{n}{2}}^n(ax^d)$ having maximal output dimension $\frac{n}{2}$.

Corollary 2. *Let s be positive odd, $n = 4s$, $m \mid n$, and $d = (2^s+1)^2$ and let $a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{x^d : x \in \mathbb{F}_{2^n}^*\}$.*

(1) Then the vectorial Boolean function $Tr_m^n(ax^d)$ is bent if and only if $(2^m - 1) \mid \frac{2^n-1}{\gcd(d,2^n-1)}$.

(2) Then there is no vectorial Boolean function of the form $Tr_m^n(ax^d)$ for even m.

Example 3. Let $s = 3, m = 3$. Then $n = 4s = 12$ and $d = (2^s + 1)^2 = 81$. For $\forall\ a \in \mathbb{F}_4 \backslash \mathbb{F}_2 \cdot \{\beta^{81} : \beta \in \mathbb{F}_{2^{12}}^*\}$, $Tr_3^{12}(ax^{81})$ is a vectorial monomial bent function.

In [4], given up to the equivalence induced by replacing a by $a\beta^d$ for $\beta \in \mathbb{F}_{2^n}^*$, it was shown that $Tr_1^n(ax^d)$ is bent if and only if $Tr_s^{3s}(a) = 0$, where the integer $s > 1$, $n = 6s$, $d = 2^{2s} + 2^s + 1$ and $a \in \mathbb{F}_{2^{3s}}^*$. Considering the equivalence induced by replacing a by $a\beta^d$ for $\beta \in \mathbb{F}_{2^n}^*$, the following theorem follows from Theorem 3 in [4].

Theorem 15. *Let $s > 1$ be an integer, $n = 6s$ and $d = 2^{2s}+2^s+1$. The monomial Boolean function $Tr_1^n(ax^d)$ is bent if and only if $a \in \{\rho : Tr_s^{3s}(\rho) = 0, \rho \in \mathbb{F}_{2^{3s}}^*\} \cdot \{x^d : x \in \mathbb{F}_{2^n}^*\}$.*

Theorem 16. *(Canteaut-Charpin-Kyureghyan Case). Let $s > 1$ be an integer, $n = 6s$ and $d = 2^{2s} + 2^s + 1$, and let $a \in \{\rho : Tr_s^{3s}(\rho) = 0, \rho \in \mathbb{F}_{2^{3s}}^*\} \cdot \{x^d : x \in \mathbb{F}_{2^n}^*\}$.*

If $m \mid 2s$, then the vectorial Boolean function $Tr_m^n(ax^d)$ is bent.

Proof. Thanks to $2^n - 1 = d(2^{2s} - 1)(2^{2s} - 2^s + 1)$, so $\frac{2^n-1}{\gcd(d,2^n-1)} = (2^{2s} - 1)(2^{2s} - 2^s + 1)$. If $m \mid 2s$, then $(2^m - 1) \mid (2^{2s} - 1)(2^{2s} - 2^s + 1)$. According to Theorem 1 and Theorem 15, we have $Tr_m^n(ax^d)$ is bent. □

In the remainder of this section, we give a construction of vectorial Boolean bent functions in the form $Tr_{\frac{n}{2}}^n(\sum_{i=1}^{2^s-1} x^{d_i})$, where d_is are Niho exponents.

Theorem 17. *(Niho Case). Let n be even, s be a positive integer, $s < \frac{n}{2}$ and $\gcd(s, \frac{n}{2}) = 1$, and $a + a^{2^{\frac{n}{2}}} \neq 0$.*

If $m \mid \frac{n}{2}$, then $Tr_m^n(\sum_{i=1}^{2^s-1} ax^{(i\cdot 2^{\frac{n}{2}-s}+1)(2^{\frac{n}{2}}-1)+1})$ is bent.

Proof. If $m \mid \frac{n}{2}$, then $(2^m - 1) \mid (2^{\frac{n}{2}} - 1)$, thus $2^{\frac{n}{2}} \equiv 1(mod\ 2^m - 1)$. For all $\lambda \in \mathbb{F}_{2^m}^*$, we obtain that $\lambda + (a\lambda)^{2^{\frac{n}{2}}} = (a + a^{\frac{n}{2}})\lambda \neq 0$.

And by Theorem 2 in [21], we have

$$(a, a, \cdots, a) \cdot \mathbb{F}_{2^m}^* \subseteq C = \{(c_1, c_2, \cdots, c_{2^s-1}) : Tr_1^n(\sum_{i=1}^{2^s-1} c_i x^{(i\cdot 2^{\frac{n}{2}-s}+1)(2^{\frac{n}{2}}-1)+1}) \text{ is bent}\}$$

According to Theorem 6, $Tr_m^n(\sum_{i=1}^{2^s-1} ax^{(i\cdot 2^{\frac{n}{2}-s}+1)(2^{\frac{n}{2}}-1)+1})$ is bent. □

5 Conclusions

This paper presents three constructions of vectorial Boolean bent functions and provides answers to E.Pasalic et al's one open problem in [32] and A.Muratović-Ribić et al.'s Open Problem 1 in [28]. Moreover, the bent properties of several types vectorial Boolean functions is analyzed.

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