New multilinear maps from ideal lattices

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Sep 5, 2015

Abstract. Recently, Hu and Jia presented an efficient attack on the GGH map. They show that the MPKE and WE based on GGH with public tools of encoding are not secure. Currently, an open problem is to fix GGH with functionality-preserving. By modifying zero-testing parameter and using modulus switching method, we present a new construction of multilinear map from ideal lattices. Our construction maintains functionality of GGH with public tools of encoding, such as applications of the GGH-based MPKE and WE. The security of our construction depends upon new hardness assumption.

Keywords. Multilinear maps, Ideal lattices, Multipartite Diffie-Hellman key exchange, Witness encryption, Zeroizing attack

1 Introduction

There are at present only three constructions of multilinear maps [GGH13, CLT13, GGH15]. The first candidate construction of multilinear maps is presented by Garg, Gentry, and Halevi (GGH) [GGH13]. Soon after, Coron, Lepoint, and Tibouchi [CLT13] (CLT) described a construction over the integers using same framework of GGH. Recently, Gentry, Gorbunov and Halevi [GGH15] constructed graph-induced multilinear maps from lattices.

However, the zeroizing attacks for CLT and GGH demonstrate that previous constructions require further improvement. On the one hand, Cheon, Han, Lee, Ryu, and Stehle recently broke the CLT construction using zeroizing attack introduced by Garg, Gentry, and Halevi. To fix the CLT construction, Garg, Gentry, Halevi and Zhandry [GGH+14], and Boneh, Wu and Zimmerman [BWZ14] presented two candidate fixes of multilinear maps over the integers. However, Coron, Lepoint, and Tibouchi showed that two candidate fixes of CLT can also be defeated using extensions of the Cheon et al.'s Attack [CHL+14]. By modifying zero-testing parameter, Coron, Lepoint and Tibouchi [CLT15] proposed a new construction of multilinear map over the integers. On the other hand, Hu and Jia [HJ15a] very recently presented an efficient attack on the GGH map, which breaks the GGH-based applications on multipartite key exchange (MPKE) and witness encryption (WE) based on the hardness of 3-exact cover problem. The Cheon and Lee [CL15] proposed an attack for the GGH map by computing a basis of secret ideal lattice.

Gu (Gu map-1) [Gu15] presented a construction of multilinear maps without encodings of zero, which is an variant of the GGH map. Since no encodings of zero are given in the public parameters, MPKE based on Gu map-1 [HJ15c] successfully avoids the attack in [HJ15a]. However, Gu map-1 cannot be used for the instance of witness encryption based on the hardness of 3-exact cover problem [HJ15b]. This is because there is no randomizer in Gu map-1. But the instance of WE based on the hardness of 3-exact cover problem is a strong application of multilinear map. Currently, an open problem is how to fix the GGH map, whilst still maintaining functionality of the original GGH.

Our results.

We first briefly recall the GGH map. The GGH map works in a polynomial ring $R = \mathbb{Z}[x]/\langle x^n + 1\rangle$, where n is a positive integer. A random large integer q, a secret short ring element $\mathbf{g} \in R$, and a secret random element $\mathbf{z} \in R_q = R/qR$ are chosen during construction, where $\mathbf{g}$ generates a principal ideal $I = \langle \mathbf{g} \rangle \subset R$ and $\mathbf{z}$ is invertible in R_q. Elements in R/I are encoded as follows: a level-k encoding of the coset $e_I = \mathbf{e} + I$ is an element of the

form $\left[\mathbf{c}/\mathbf{z}^k\right]_q$, where the $\mathbf{c}\in e_I$ norm is short. Encodings can both be added and multiplied if the numerator norm remains smaller than q. For a level-κ encoding $\mathbf{u}=\left[\mathbf{c}/\mathbf{z}^\kappa\right]_q$, the encoding $\mathbf{u}$ can be determined as zero by computing $\left[\mathbf{u}\cdot\mathbf{p}_{zt}\right]_q$, where $\mathbf{p}_{zt}=\left[\mathbf{h}\mathbf{z}^\kappa/\mathbf{g}\right]_q$ is a zero-testing parameter. If the norm of $\left[\mathbf{u}\cdot\mathbf{p}_{zt}\right]_q$ is small, then $\mathbf{u}$ is the encoding of zero; otherwise, $\mathbf{u}$ is the encoding of non-zero.

Our main contribution is to construct a new multilinear map using ideal lattices. Our construction improves the origin GGH map in three aspects.

(1) We introduce new noise term to avoid the zeroizing attack problem of GGH. Let $\mathrm{par}_0=\left\{q,\mathbf{y}=\left[(1+\mathbf{ag})/\mathbf{z}\right]_q,\mathbf{x}_i=\left[(\mathbf{a}_i\mathbf{g})/\mathbf{z}\right]_q,i\in[\![\tau]\!],\mathbf{p}_{zt}=\left[\mathbf{h}\mathbf{z}^\kappa/\mathbf{g}\right]_q\right\}$ be the public parameters of GGH. Given arbitrary level-k encoding $\mathbf{u}=\left[\mathbf{c}/\mathbf{z}^k\right]_q$, one can compute $\mathbf{v}_{k,i,j}=\left[\mathbf{u}\cdot\mathbf{x}_i^j\cdot\mathbf{y}^{\kappa-k-j}\cdot\mathbf{p}_{zt}\right]_q=\left[\mathbf{hc}(\mathbf{a}_i)^j\mathbf{g}^{j-1}(1+\mathbf{ag})^{\kappa-k-j}\right]_q$, where $1\le k<\kappa,1\le j<\kappa$ and $k+j\le\kappa$ using so-called zeroizing attack method. It is easy to verify that $\mathbf{v}_{k,i,j}$ is not reduced modulo q. As a result, one one can compute a basis of the secret ring element $\mathbf{g}$. Using this method, Hu and Jia [HJ15a] have broken two applications of MPKE and WE based on GGH. To improve GGH and avoid the zeroizing attack, one needs to introduce new noise term for $\mathbf{v}_{k,i,j}$. If one can add a random noise to $\mathbf{v}_{k,i,j}$, then adversary cannot yield a basis of $\mathbf{g}$. We use two ring elements $\mathbf{g}_1,\mathbf{g}_2$ and double-encodings in our construction to achieve this goal.

(2) We modify zero testing parameter and apply modulus switching method to add new noise term. We design new zero-testing parameter $\mathbf{p}_{zt,1}=\left[\mathbf{z}^\kappa(\mathbf{h}_1\cdot\left[\mathbf{g}_1^{-1}\right]_{q_1}+\mathbf{h}_2\cdot\left[\mathbf{g}_2^{-1}\right]_{q_1}\right]_{q_1}$ and adaptively modify encodings $\mathbf{y}_1=\left[(1+\mathbf{ag}_1)/\mathbf{z}_1\right]_{q_1}$, $\mathbf{x}_{1,i}=\left[(\mathbf{a}_i\mathbf{g}_1)/\mathbf{z}_1\right]_{q_1},i\in[\![\tau]\!]$. The problem is how to remove encoding of non-zero element for ideal lattice $\langle\mathbf{g}_2\rangle$. Roughly speaking, one must generate encoding of zero for $\langle\mathbf{g}_2\rangle$. For this purpose, we use another modulo q_2 and generate some encodings $\mathbf{y}_2=\left[(\mathbf{e}+\mathbf{bg}_2)/\mathbf{z}_2\right]_{q_2}$, $\mathbf{x}_{2,i}=\left[(\mathbf{e}_i+\mathbf{b}_i\mathbf{g}_2)/\mathbf{z}_2\right]_{q_2}$, $i\in[\![\tau]\!]$ over modulo q_2 such that $\mathbf{e}=(1+\mathbf{ag}_1)\bmod\mathbf{g}_2$ and $\mathbf{e}_i=(\mathbf{a}_i\mathbf{g}_1)\bmod\mathbf{g}_2$. To obtain encoding of zero for $\langle\mathbf{g}_2\rangle$, we design a new zero testing parameter $\mathbf{p}_{zt,2}=\left[\mathbf{z}_2^\kappa(\mathbf{h}_2+\mathbf{hg}_2)\left\lfloor(q_2/q_1)\cdot\left[\mathbf{g}_2^{-1}\right]_{q_1}\right\rfloor\right]_{q_2}$ over modulo q_2. When generating a level-k encoding $\mathbf{u}_1=\left[\mathbf{c}_1/\mathbf{z}^k\right]_{q_1}$, one also generates another level-k encoding $\mathbf{u}_2=\left[\mathbf{c}_2/\mathbf{z}_2^k\right]_{q_2}$ over modulo q_2 such that $\mathbf{c}_1=\mathbf{c}_2\bmod\mathbf{g}_2$. Hence, given a level-$\kappa$ encoding $(\mathbf{u}_1,\mathbf{u}_2)$, we can determine whether the encoding of $\mathbf{u}_1$ is zero for $\langle\mathbf{g}_1\rangle$ by computing $\left[\left\lfloor(q_2/q_1)\cdot\left[\mathbf{u}_1\cdot\mathbf{p}_{zt,1}\right]_{q_1}\right\rfloor-\left[\mathbf{u}_2\cdot\mathbf{p}_{zt,2}\right]_{q_2}\right]_{q_2}$. For an arbitrary level-k encoding $(\mathbf{u}_1,\mathbf{u}_2)$, one can compute

$$\mathbf{v}_{k,i,j}=\left[\left\lfloor(q_2/q_1)\cdot\left[\mathbf{u}_1\cdot\mathbf{x}_{1,i}^j\cdot\mathbf{y}_1^{\kappa-k-j}\cdot\mathbf{p}_{zt,1}\right]_{q_1}\right\rfloor-\left[\mathbf{u}_2\cdot\mathbf{x}_{2,i}^j\cdot\mathbf{y}_2^{\kappa-k-j}\cdot\mathbf{p}_{zt,2}\right]_{q_2}\right]_{q_2},$$

where $1\le k<\kappa,1\le j<\kappa$ and $k+j\le\kappa$ using zeroizing attack method. It is easy to see that $\mathbf{v}_{k,i,j}$ is not reduced modulo q_2. However, one can no longer obtain a basis of $\mathbf{g}_1$ using $\mathbf{v}_{k,i,j}$.

(3) Our new construction seemly supports more applications than the GGH map. Owing to adding new noise term, one can no longer yield a basis of $\mathbf{g}_1$. Hence, we conjecture that the membership group problem (SubM) and the decisional linear (DLIN) problem are hard in our construction. However, in the GGH map, one can compute non-reduced ring elements over modulus q_1 and a basis of $\mathbf{g}_1$. As a result, the SubM problem and the DLIN problem are easy in the GGH map.

Our second contribution is to describe the applications of MPKE and WE using our new multilinear map. Since these applications are attacked by [HJ15a], fix for them is urgently required. The construction of MPKE and WE based on our map are same as ones based on GGH.

Organization. Section 2 recalls some background. Section 3 describes our new construction using ideal lattices. Section 4 presents two applications of MPKE and WE based on our construction. Finally, Section 5 draws conclusion.

2 Preliminaries

2.1 Notations

We denote $\mathbb{Z}, \mathbb{Q}, \mathbb{R}$ the ring of integers, the field of rational numbers, and the field of real numbers. We take n as a positive integer and a power of 2. Notation $[\![n]\!]$ denotes the set $\{1, 2, \cdots, n\}$, and $[a]_q$ the absolute minimum residual system $[a]_q = a \bmod q \in (-q/2, q/2]$. Vectors and matrices are denoted in bold, such as $\mathbf{a}, \mathbf{b}, \mathbf{c}$ and $\mathbf{A}, \mathbf{B}, \mathbf{C}$. The j-th entry of $\mathbf{a}$ is denoted as a_j, the element of the i-th row and j-th colomn of $\mathbf{A}$ is denoted as $A_{i,j}$ (or $A[i, j]$). Notation $\|\mathbf{a}\|_\infty$ ($\|\mathbf{a}\|$ for short) denotes the infinity norm of $\mathbf{a}$. The polynomial ring $\mathbb{Z}[x]/\langle x^n + 1\rangle$ is denoted by R, and $\mathbb{Z}_q[x]/\langle x^n + 1\rangle$ by R_q. The elements in R and R_q are denoted in bold as well. Similarly, notation $[\mathbf{a}]_q$ denotes each entry (or each coefficient) $a_i \in (-p/2, p/2]$ of $\mathbf{a}$.

2.2 Lattices and Ideal Lattices

An n-dimension full-rank lattice $L \subset \mathbb{R}^n$ is the set of all integer linear combinations $\sum_{i=1}^n y_i \mathbf{b}_i$ of n linearly independent vectors $\mathbf{b}_i \in \mathbb{R}^n$. If we arrange the vectors $\mathbf{b}_i$ as the columns of matrix $\mathbf{B} \in \mathbb{R}^{n \times n}$, then $L = \{\mathbf{By} : \mathbf{y} \in \mathbb{Z}^n\}$. We say that $\mathbf{B}$ spans L if $\mathbf{B}$ is a basis for L. Given a basis $\mathbf{B}$ of L, we define $P(\mathbf{B}) = \{\mathbf{By} \mid \mathbf{y} \in \mathbb{R}^n, \forall i : -1/2 \le y_i < 1/2\}$ as the parallelization corresponding to $\mathbf{B}$. Let $\det(\mathbf{B})$ denote the determinant of $\mathbf{B}$.

Given $\mathbf{g} \in R$, let $I = \langle \mathbf{g} \rangle$ be the principal ideal lattice in R generated by $\mathbf{g}$, whose $\mathbb{Z}$-basis is $Rot(\mathbf{g}) = (\mathbf{g}, x \cdot \mathbf{g}, ..., x^{n-1} \cdot \mathbf{g})$.

Given $\mathbf{c} \in \mathbb{R}^n, \sigma > 0$, the Gaussian distribution of a lattice L is defined as $\forall \mathbf{x} \in L$, $D_{L,\sigma,\mathbf{c}} = \rho_{\sigma,\mathbf{c}}(\mathbf{x}) / \rho_{\sigma,\mathbf{c}}(L)$, where $\rho_{\sigma,\mathbf{c}}(\mathbf{x}) = \exp(-\pi \|\mathbf{x} - \mathbf{c}\|^2 / \sigma^2)$, $\rho_{\sigma,\mathbf{c}}(L) = \sum_{x \in L} \rho_{\sigma,\mathbf{c}}(\mathbf{x})$. In the following, we will write $D_{\mathbb{Z}^n,\sigma,0}$ as $D_{\mathbb{Z}^n,\sigma}$. We denote a Gaussian sample as $\mathbf{x} \leftarrow D_{L,\sigma}$ (or $\mathbf{d} \leftarrow D_{I,\sigma}$) over the lattice L (or ideal lattice I).

2.3 Multilinear Maps

Definition 2.1 (Multilinear Map [BS03]). For $\kappa+1$ cyclic groups $G_1,...,G_\kappa,G_T$ of the same order q, a κ-multilinear map $e: G_1\times\cdots\times G_\kappa \to G_T$ has the following properties:

(1) Elements $\left\{g_j \in G_j\right\}_{j=1,...,\kappa}$, index $j\in[\![\kappa]\!]$, and integer $a\in\mathbb{Z}_q$ hold that

$$e(g_1,\cdots,a\cdot g_j,\cdots,g_\kappa)=a\cdot e(g_1,\cdots,g_\kappa)$$

(2) Map e is non-degenerate in the following sense: if elements $\left\{g_j \in G_j\right\}_{j=1,...,\kappa}$ are generators of their respective groups, then $e(g_1,\cdots,g_\kappa)$ is a generator of G_T.

Definition 2.2 (κ-Graded Encoding System [GGH13]). A κ-graded encoding system over R is a set system of $S=\left\{S_j^{(\alpha)}\subset R:\alpha\in R, j\in[\![\kappa]\!]\right\}$ with the following properties:

(1) For every index $j\in[\![\kappa]\!]$, the sets $\left\{S_j^{(\alpha)}:\alpha\in R\right\}$ are disjoint.

(2) Binary operations '$+$' and '$-$' exist, such that every α_1,α_2, every index $j\in[\![\kappa]\!]$, and every $u_1\in S_j^{(\alpha_1)}$ and $u_2\in S_j^{(\alpha_2)}$ hold that $u_1+u_2\in S_j^{(\alpha_1+\alpha_2)}$ and $u_1-u_2\in S_j^{(\alpha_1-\alpha_2)}$, where $\alpha_1+\alpha_2$ and $\alpha_1-\alpha_2$ are the addition and subtraction operations in R respectively.

(3) Binary operation '$\times$' exists, such that every α_1,α_2, every index $j_1,j_2\in[\![\kappa]\!]$ with $j_1+j_2\le\kappa$, and every $u_1\in S_{j_1}^{(\alpha_1)}$ and $u_2\in S_{j_2}^{(\alpha_2)}$ hold that $u_1\times u_2\in S_{j_1+j_2}^{(\alpha_1\times\alpha_2)}$, where $\alpha_1\times\alpha_2$ is the multiplication operation in R and j_1+j_2 is the integer addition.

3 Our new construction

Setting the parameters. Let λ be the security parameter, κ the multilinearity level, n the dimension of elements of R. Concrete parameters are set as $\sigma=\sqrt{\lambda n}$, $\sigma'=\lambda n^{1.5}$, $\sigma^*=2^\lambda$, $q_1,q_2\ge 2^{8\kappa\lambda}n^{O(\kappa)}$, $m=2$, $n>\widetilde{O}(\kappa\lambda^2)$, $\tau=O(n^2)$, $\rho=O(n)$.

3.1 Construction

Instance generation: $(\text{par})\leftarrow \text{InstGen}(1^\lambda,1^\kappa)$.

(1) Choose two primes $q_t\ge 2^{8\kappa\lambda}n^{O(\kappa)}, t\in[\![2]\!]$ with $q_1>2^\lambda\cdot q_2$.

(2) Choose $\mathbf{g}_t\leftarrow D_{\mathbb{Z}^n,\sigma}, t\in[\![2]\!]$ in R so that $\left\|\mathbf{g}_t^{-1}\right\|<n^2$.

(3) Choose $\mathbf{a}_1,\mathbf{b}_2,\mathbf{a}_{1,i},\mathbf{b}_{2,i}\leftarrow D_{\mathbb{Z}^n,\sigma'}, i\in[\![\tau]\!]$ in R;

Choose $\mathbf{h},\mathbf{h}_1,\mathbf{h}_2\leftarrow D_{\mathbb{Z}^n,\sqrt{q_2}}$ in R.

(4) Choose random elements $\mathbf{z}_t\leftarrow R_{q_t}, t\in[\![2]\!]$ with $\mathbf{z}_t^{-1}\in R_{q_t}$.

(5) Set $\mathbf{e}_1=(\mathbf{a}_1\mathbf{g}_1+1)\bmod\mathbf{g}_2$, namely $\mathbf{a}_1\mathbf{g}_1+1=\mathbf{b}_1\mathbf{g}_2+\mathbf{e}_1$ so that $\left\|\mathbf{b}_1\right\|<n^2$;

$\mathbf{e}_{1,i}=(\mathbf{a}_{1,i}\mathbf{g}_1)\bmod\mathbf{g}_2, i\in[\![\tau]\!]$, namely $\mathbf{a}_{1,i}\mathbf{g}_1=\mathbf{b}_{1,i}\mathbf{g}_2+\mathbf{e}_{1,i}$ so that $\left\|\mathbf{b}_{1,i}\right\|<n^2$.

(6) Set $\mathbf{y}_1=\left[\frac{\mathbf{a}_1\mathbf{g}_1+1}{\mathbf{z}_1}\right]_{q_1}=\left[\frac{\mathbf{b}_1\mathbf{g}_2+\mathbf{e}_1}{\mathbf{z}_1}\right]_{q_1}$ and $\mathbf{x}_{1,i}=\left[\frac{\mathbf{a}_{1,i}\mathbf{g}_1}{\mathbf{z}_1}\right]_{q_1}=\left[\frac{\mathbf{b}_{1,i}\mathbf{g}_2+\mathbf{e}_{1,i}}{\mathbf{z}_1}\right]_{q_1}$;

$$\mathbf{y}_2 = \left[\frac{\mathbf{b}_2 \mathbf{g}_2 + \mathbf{e}_1}{\mathbf{z}_2} \right]_{q_2} \quad \text{and} \quad \mathbf{x}_{2,i} = \left[\frac{\mathbf{b}_{2,i} \mathbf{g}_2 + \mathbf{e}_{1,i}}{\mathbf{z}_2} \right]_{q_2}.$$

(7) Set $\mathbf{p}_{zt,1} = \left[\mathbf{z}_1^{\kappa} \left(\mathbf{h}_1 \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} + \mathbf{h}_2 \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right) \right]_{q_1}$,

$$\mathbf{p}_{zt,2} = \left[\mathbf{z}_2^{\kappa} (\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2) \mathbf{g}_{2,q_2}^{-1} \right]_{q_2} \quad \text{with} \quad \mathbf{g}_{2,q_2}^{-1} = \left\lfloor (q_2 / q_1) \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1}) \right\rfloor.$$

(8) Output the public parameters $\text{par} = \left\{ \left\{ q_t, \mathbf{y}_t, \left\{ \mathbf{x}_{t,i} \right\}_{i \in [\![\tau]\!]}, \mathbf{p}_{zt,t} \right\}_{t \in [\![2]\!]} \right\}$.

Generating level-k encoding: $(\mathbf{u}_1, \mathbf{u}_2) \leftarrow \text{Enc}(\text{par}, k, \mathbf{d})$.

(1) Sample $\mathbf{r}_i \leftarrow D_{\mathbb{Z}^n, \sigma^*}, i \in [\![\tau]\!]$;

(2) Given $\mathbf{d} \leftarrow D_{\mathbb{Z}^n, \sigma'}$, compute $\mathbf{u}_t = \left[\mathbf{d} \cdot (\mathbf{y}_t)^k + \sum_{i=1}^{\tau} \mathbf{r}_i \cdot (\mathbf{x}_{t,i})^k \right]_{q_t}, t \in [\![2]\!]$;

(3) Output $(\mathbf{u}_1, \mathbf{u}_2)$ as a level-k encoding of $\mathbf{d}$.

Adding encodings: $(\mathbf{u}_1, \mathbf{u}_2) \leftarrow \text{Add}(\text{par}, k, (\mathbf{u}_{1,1}, \mathbf{u}_{2,1}), \cdots, (\mathbf{u}_{1,s}, \mathbf{u}_{2,s}))$.

(1) Given s level-k encodings $(\mathbf{u}_{1,l}, \mathbf{u}_{2,l})$, compute $\mathbf{u}_t = \left[\sum_{l=1}^{s} \mathbf{u}_{t,l} \right]_{q_t}$.

(2) Output a level-k encoding $(\mathbf{u}_1, \mathbf{u}_2)$.

Multiplying encodings: $(\mathbf{u}_1, \mathbf{u}_2) \leftarrow \text{Mul}(\text{par}, 1, (\mathbf{u}_{1,1}, \mathbf{u}_{2,1}), \cdots, (\mathbf{u}_{1,k}, \mathbf{u}_{2,k}))$.

(1) Given k level-1 encodings $(\mathbf{u}_{1,l}, \mathbf{u}_{2,l})$, compute $\mathbf{u}_t = \left[\prod_{l=1}^{k} \mathbf{u}_{t,l} \right]_{q_t}$.

(2) Output a level-k encoding $(\mathbf{u}_1, \mathbf{u}_2)$.

Zero testing: $\text{isZero}(\text{par}, (\mathbf{u}_1, \mathbf{u}_2))$.

Given a level-κ encoding $(\mathbf{u}_1, \mathbf{u}_2)$, to determine whether $\mathbf{u}_1$ is a level-κ encoding of zero for $\mathbf{g}_1$, we compute $\mathbf{v} = \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2}$ and check whether $\|\mathbf{v}\|$ is short:

$$\text{isZero}(\text{par}, (\mathbf{u}_1, \mathbf{u}_2)) = \begin{cases} 1 & \text{if } \|\mathbf{v}\| < q_2^{3/4} \\ 0 & \text{otherwise} \end{cases}.$$

Extraction: $sk \leftarrow \text{Ext}(\text{par}, (\mathbf{u}_1, \mathbf{u}_2))$.

Given a level-κ encoding $(\mathbf{u}_1, \mathbf{u}_2)$, we compute $\mathbf{v} = \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2}$, and collect $\eta = (\log q)/4 - \lambda$ most-significant bits of each of the n coefficients of $\mathbf{v}$:

$$\text{Ext}(\text{par}, (\mathbf{u}_1, \mathbf{u}_2)) = \text{Extract}_s \left(\text{msbs}_{\eta} \left(\left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2} \right) \right).$$

Remark 3.1 In our construction, modulus switching method is to thwart the attack in [CGH+15].

3.2 Correctness

Lemma 3.2 The algorithm $\mathrm{InstGen}(1^{\lambda},1^{\kappa})$ runs in polynomial time.

Lemma 3.3 The encoding $(\mathbf{u}_1,\mathbf{u}_2)\leftarrow \mathrm{Enc}(\mathrm{par},k,\mathbf{d})$ is a level-k encoding.

Proof. We only need to show that $\mathbf{u}_1$ is a level-k encoding of $\mathbf{d}$ for the ideal lattice $\langle \mathbf{g}_1\rangle$, and level-$k$ encodings $\mathbf{u}_1,\mathbf{u}_2$ encode same level-0 encoding for the ideal lattice $\langle \mathbf{g}_2\rangle$.

(1) By $\mathbf{u}_t=\left[\mathbf{d}\cdot(\mathbf{y}_t)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{x}_{t,i})^k\right]_q$, for $\langle \mathbf{g}\rangle$ we have

$$\begin{aligned}\mathbf{u}_1&=\left[\mathbf{d}\cdot(\mathbf{y}_1)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{x}_{1,i})^k\right]_{q_1}\\&=\left[\mathbf{d}\cdot\left(\frac{\mathbf{a}_1\mathbf{g}_1+1}{\mathbf{z}_1}\right)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot\left(\frac{\mathbf{a}_{1,i}\mathbf{g}_1}{\mathbf{z}_1}\right)^k\right]_{q_1}\\&=\left[\frac{\mathbf{d}\cdot(\mathbf{a}_1\mathbf{g}_1+1)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{a}_{1,i}\mathbf{g}_1)^k}{\mathbf{z}_1^k}\right]_{q_1},\\&=\left[\frac{\mathbf{a}\mathbf{g}_1+\mathbf{d}}{\mathbf{z}_1^k}\right]_{q_1}\end{aligned}$$

where $\mathbf{a}=\left(\mathbf{d}\cdot(\mathbf{a}_1\mathbf{g}_1+1)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{a}_{1,i}\mathbf{g}_1)^k-\mathbf{d}\right)/\mathbf{g}_1$.

Thus, $\mathbf{u}_1$ is a level-k encoding of the level-0 encoding $\mathbf{d}$ for $\langle \mathbf{g}_1\rangle$.

(2) Similarly, for $\langle \mathbf{g}_2\rangle$ we have

$$\begin{aligned}\mathbf{u}_t&=\left[\mathbf{d}\cdot(\mathbf{y}_t)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{x}_{t,i})^k\right]_{q_t}\\&=\left[\mathbf{d}\cdot\left(\frac{\mathbf{b}_t\mathbf{g}_2+\mathbf{e}_1}{\mathbf{z}_t}\right)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot\left(\frac{\mathbf{b}_{t,i}\mathbf{g}_2+\mathbf{e}_{1,i}}{\mathbf{z}_t}\right)^k\right]_{q_t},\\&=\left[\frac{\mathbf{c}_t\mathbf{g}_2+\mathbf{e}}{\mathbf{z}_t^k}\right]_{q_t}\end{aligned}$$

where $\mathbf{e}=\mathbf{d}\cdot(\mathbf{e}_1)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{e}_{1,i})^k$, $\mathbf{c}_t=\left(\mathbf{d}\cdot(\mathbf{b}_t\mathbf{g}_2+\mathbf{e}_1)^k+\sum_{i=1}^{\tau}\mathbf{r}_i\cdot(\mathbf{b}_{t,i}\mathbf{g}_2+\mathbf{e}_{1,i})^k-\mathbf{e}\right)/\mathbf{g}_2$.

The level-k encodings $\mathbf{u}_1,\mathbf{u}_2$ encode same level-0 encoding for $\langle \mathbf{g}_2\rangle$ □

Lemma 3.4 The encoding $(\mathbf{u}_1,\mathbf{u}_2)\leftarrow \mathrm{Add}\left(\mathrm{par},k,(\mathbf{u}_{1,1},\mathbf{u}_{2,1}),\cdots,(\mathbf{u}_{1,s},\mathbf{u}_{2,s})\right)$ is a level-k encoding.

Proof. Since for $\langle \mathbf{g}_1\rangle$, a level-$k$ encoding $\mathbf{u}_{1,l}$ has the form $\mathbf{u}_{1,l}=\left[\frac{\mathbf{r}_l\mathbf{g}_1+\mathbf{d}_l}{\mathbf{z}_1^k}\right]_{q_1}$, then the sum is

$$\mathbf{u}_1=\left[\sum_{l=1}^{s}\mathbf{u}_{1,l}\right]_{q_1}=\left[\frac{\sum_{l=1}^{s}(\mathbf{r}_l\mathbf{g}_1+\mathbf{d}_l)}{\mathbf{z}_1^k}\right]_{q_1}=\left[\frac{\mathbf{r}\mathbf{g}_1+\mathbf{d}}{\mathbf{z}_1^k}\right]_{q_1},$$

where $\mathbf{r}=\sum_{l=1}^{s}\mathbf{r}_l$ and $\mathbf{d}=\sum_{l=1}^{s}\mathbf{d}_l$.

Namely, $\mathbf{u}_1$ is a level-k encoding for $\langle \mathbf{g}_1 \rangle$.

Again for $\langle \mathbf{g}_2 \rangle$, a level-$k$ encoding $\mathbf{u}_{t,l}, t \in [\![2]\!]$ has the form $\mathbf{u}_{t,l}=\left[\frac{\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l}{\mathbf{z}_t^k}\right]_{q_t}$.

Thus, we have

$$\mathbf{u}_t=\left[\sum_{l=1}^{s}\mathbf{u}_{t,l}\right]_{q_t}=\left[\frac{\sum_{l=1}^{s}(\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l)}{\mathbf{z}_t^k}\right]_{q_t}=\left[\frac{\mathbf{c}_t\mathbf{g}_2+\mathbf{e}}{\mathbf{z}_t^k}\right]_{q_t},$$

where $\mathbf{c}_t=\sum_{l=1}^{s}\mathbf{c}_{t,l}$ and $\mathbf{e}=\sum_{l=1}^{s}\mathbf{e}_l$.

That is, level-k encodings $\mathbf{u}_1, \mathbf{u}_2$ encode same level-0 encoding for $\langle \mathbf{g}_2 \rangle$. □

Lemma 3.5 The encoding $(\mathbf{u}_1,\mathbf{u}_2) \leftarrow \mathrm{Mul}\left(\mathrm{par}, 1, (\mathbf{u}_{1,1},\mathbf{u}_{2,1}), \cdots, (\mathbf{u}_{1,k},\mathbf{u}_{2,k})\right)$ is a level-k encoding.

Proof. Since $\mathbf{u}_{1,l}=\left[\frac{\mathbf{r}_l\mathbf{g}_1+\mathbf{d}_l}{\mathbf{z}_1}\right]_{q_1}$ for $\langle \mathbf{g}_1 \rangle$, their product is:

$$\mathbf{u}_1=\left[\prod_{l=1}^{k}\mathbf{u}_{1,l}\right]_{q_1}=\left[\prod_{l=1}^{k}\frac{\mathbf{r}_l\mathbf{g}+\mathbf{d}_l}{\mathbf{z}_1}\right]_{q_1}=\left[\frac{\prod_{j=1}^{k}(\mathbf{r}_l\mathbf{g}+\mathbf{d}_l)}{\mathbf{z}_1^k}\right]_{q_1}=\left[\frac{\mathbf{r}\mathbf{g}+\mathbf{d}}{\mathbf{z}_1^k}\right]_{q_1},$$

where $\mathbf{d}=\prod_{l=1}^{k}\mathbf{d}_l$, $\mathbf{r}=(\prod_{l=1}^{k}(\mathbf{r}_l\mathbf{g}_1+\mathbf{d}_l)-\mathbf{d})/\mathbf{g}_1$.

Again for $\langle \mathbf{g}_2 \rangle$, the level-$1$ encoding $\mathbf{u}_{t,l}, t \in [\![2]\!]$ has the form $\mathbf{u}_{t,l}=\left[\frac{\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l}{\mathbf{z}_t}\right]_{q_t}$.

Thus, we have

$$\mathbf{u}_t=\left[\prod_{l=1}^{k}\frac{\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l}{\mathbf{z}_t}\right]_{q_t}=\left[\frac{\prod_{l=1}^{k}(\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l)}{\mathbf{z}_t^k}\right]_{q_t}=\left[\frac{\mathbf{c}_t\mathbf{g}_2+\mathbf{e}}{\mathbf{z}_t^k}\right]_{q_t},$$

where $\mathbf{e}=\prod_{l=1}^{k}\mathbf{e}_l$ and $\mathbf{c}_t=(\prod_{l=1}^{k}(\mathbf{c}_{t,l}\mathbf{g}_2+\mathbf{e}_l)-\mathbf{e})/\mathbf{g}_2$. □

Before proving correctness of the zero-testing algorithm, we first describe two lemmas about modulus switching.

Lemma 3.6. Suppose that $\mathbf{v}_t=\left[\mathbf{w}_t\cdot\left[\mathbf{g}_t^{-1}\right]_{q_1}\right]_{q_1}, t \in [\![2]\!]$ with $\|\mathbf{w}_t\| << q_2$. Then,

$$\left\lfloor\frac{q_2}{q_1}\mathbf{v}_t\right\rceil=\left[\mathbf{w}_t\cdot\left\lfloor\frac{q_2}{q_1}\left[\mathbf{g}_t^{-1}\right]_{q_1}\right\rceil\right]_{q_2}+\boldsymbol{\delta}_t \text{ with } \|\boldsymbol{\delta}_t\|<l_1(\mathbf{w}_t),$$

where $l_1(\mathbf{w}_t)$ is l_1-norm of $\mathbf{w}_t$.

Proof. By $\mathbf{v}_t=\left[\mathbf{w}_t\cdot\left[\mathbf{g}_t^{-1}\right]_{q_1}\right]_{q_1}=\mathbf{w}_t\cdot\left[\mathbf{g}_t^{-1}\right]_{q_1}-q_1\cdot\mathbf{k}_t$, we get

$$\frac{q_2}{q_1}\mathbf{v}_t=\mathbf{w}_t\cdot\left(\frac{q_2}{q_1}\cdot\left[\mathbf{g}_t^{-1}\right]_{q_1}\right)-q_2\cdot\mathbf{k}_t,$$

$$\mathbf{w}_t \cdot \left(\frac{q_2}{q_1} \cdot \left[\mathbf{g}_t^{-1} \right]_{q_1} \right) = \mathbf{w}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor + \mathbf{w}_t \cdot \left(\frac{q_2}{q_1} \cdot \left[\mathbf{g}_t^{-1} \right]_{q_1} - \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right),$$
$$= \mathbf{w}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor + \mathbf{w}_t \cdot \boldsymbol{\varepsilon}_t$$

where $\|\boldsymbol{\varepsilon}_t\| < 1$.

Therefore,

$$\begin{aligned} \left\lfloor \frac{q_2}{q_1} \mathbf{v}_t \right\rfloor &= \left\lfloor \frac{q_2}{q_1} \left[\mathbf{w}_t \cdot \left[\mathbf{g}_t^{-1} \right]_{q_1} \right]_{q_1} \right\rfloor \\ &= \left\lfloor \mathbf{w}_t \cdot \left(\frac{q_2}{q_1} \cdot \left[\mathbf{g}_t^{-1} \right]_{q_1} \right) - q_2 \cdot \mathbf{k}_t \right\rfloor \\ &= \mathbf{w}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor + \left\lfloor \mathbf{w}_t \cdot \boldsymbol{\varepsilon}_t \right\rfloor - q_2 \cdot \mathbf{k}_t \\ &= \left[\mathbf{w}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right]_{q_2} + q_2 \cdot \boldsymbol{\beta}_t + \left\lfloor \mathbf{w}_t \cdot \boldsymbol{\varepsilon}_t \right\rfloor - q_2 \cdot \mathbf{k}_t \end{aligned},$$

By $\left\| \left\lfloor \frac{q_2}{q_1} \mathbf{v}_t \right\rfloor \right\| < q_2 / 2$, we have $\boldsymbol{\beta}_t = \mathbf{k}_t$ with high probability. That is, $\left\lfloor \frac{q_2}{q_1} \mathbf{v}_t \right\rfloor = \left[\mathbf{w}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right]_{q_2} + \boldsymbol{\delta}_t$ with $\boldsymbol{\delta}_t = \left\lfloor \mathbf{w}_t \cdot \boldsymbol{\varepsilon}_t \right\rfloor$. By $\|\boldsymbol{\varepsilon}_t\| < 1$, we obtain $\|\boldsymbol{\delta}_t\| < l_1(\mathbf{w}_t)$.

Lemma 3.7 Suppose that $\boldsymbol{\alpha}_t = \left[\mathbf{g}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right]_{q_2}$. Then $\|\boldsymbol{\alpha}_t\| < l_1(\mathbf{g}_t)$.

Proof. Assume that $\mathbf{v}_t = \left\lfloor \mathbf{w}_t \cdot \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor_{q_1}$ with $\mathbf{w}_t = \mathbf{g}_t$. Then by Lemma 3.6, we have

$$0 = \left\lfloor \frac{q_2}{q_1} \mathbf{v}_t \right\rfloor = \left[\mathbf{g}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right]_{q_2} + \boldsymbol{\delta}_t \text{ and } \|\boldsymbol{\delta}_t\| \le l_1(\mathbf{g}_t).$$

Thus, $\boldsymbol{\alpha}_t = \left[\mathbf{g}_t \cdot \left\lfloor \frac{q_2}{q_1} \left[\mathbf{g}_t^{-1} \right]_{q_1} \right\rfloor \right]_{q_2} = -\boldsymbol{\delta}_t$ and $\|\boldsymbol{\alpha}_t\| < l_1(\mathbf{g}_t)$.

Lemma 3.8 The zero testing $\text{isZero}\left(\text{par}, \left(\mathbf{u}_1, \mathbf{u}_2\right)\right)$ correctly determines whether $\mathbf{u}_1$ is a level-κ encoding of zero for $\langle \mathbf{g}_1 \rangle$.

Proof. Given a level-κ encoding $\left(\mathbf{u}_1, \mathbf{u}_2\right)$, we compute $\mathbf{v} = \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2}$ and check whether $\|\mathbf{v}\|$ is short:

$$\text{isZero}\left(\text{par}, \left(\mathbf{u}_1, \mathbf{u}_2\right)\right) = \begin{cases} 1 & \text{if } \|\mathbf{v}\| < q_2^{3/4} \\ 0 & \text{otherwise} \end{cases}.$$

Without loss of generality, we assume $\mathbf{u}_1 = \left\lceil \frac{\mathbf{r}\mathbf{g}_1 + \mathbf{d}}{\mathbf{z}_1^{\kappa}} \right\rfloor_{q_1}$ and $\mathbf{u}_t = \left\lceil \frac{\mathbf{c}_t\mathbf{g}_2 + \mathbf{e}}{\mathbf{z}_t^{\kappa}} \right\rfloor_{q_t}$. So,

$$\begin{aligned}
\mathbf{v}_1 &= \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor \\
&= \left\lfloor \frac{q_2}{q_1} \cdot \left[\left(\mathbf{h}_1 \left(\mathbf{r}\mathbf{g}_1 + \mathbf{d} \right) \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} + \mathbf{h}_2 \left(\mathbf{c}_1\mathbf{g}_2 + \mathbf{e} \right) \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right) \right]_{q_1} \right\rfloor \\
&= \left\lfloor \frac{q_2}{q_1} \cdot \left[\left(\mathbf{h}_1\mathbf{r}\mathbf{g}_1 \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} + \mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} + \mathbf{h}_2\mathbf{c}_1\mathbf{g}_2 \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} + \mathbf{h}_2\mathbf{e} \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right) \right]_{q_1} \right\rfloor, \\
&= \left\lfloor \frac{q_2}{q_1} \cdot \left[\left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1 + \mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} + \mathbf{h}_2\mathbf{e} \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right) \right]_{q_1} \right\rfloor \\
&= \left\lfloor \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1 \right) \right\rfloor + \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1} \right]_{q_1} \right]_{q_1} \right\rfloor + \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_2\mathbf{e} \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right]_{q_1} \right\rfloor \right\rfloor_{q_2} \\
\mathbf{v}_2 &= \left\lfloor \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right\rfloor_{q_2} \\
&= \left[\left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2 \right) \mathbf{c}_2\mathbf{g}_2 \cdot \mathbf{g}_{2,q_2}^{-1} + \mathbf{h}\mathbf{g}_2 \cdot \mathbf{g}_{2,q_2}^{-1} + \mathbf{h}_2\mathbf{e} \cdot \mathbf{g}_{2,q_2}^{-1} \right]_{q_2}, \\
&= \left[\left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2 \right) \mathbf{c}_2\boldsymbol{\alpha}_2 + \mathbf{h}\boldsymbol{\alpha}_2 + \mathbf{h}_2\mathbf{e} \cdot \mathbf{g}_{2,q_2}^{-1} \right]_{q_2}
\end{aligned}$$

where $\left\| \boldsymbol{\alpha}_2 \right\| < l_1(\mathbf{g}_2)$ by Lemma 3.7.

If $\mathbf{u}_1$ is a level-κ encoding of zero for $\langle \mathbf{g}_1 \rangle$, namely $\mathbf{d} = 0$. Thus, we have

$$\begin{aligned}
\mathbf{v} &= \left\lfloor \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_1 \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2 \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right\rfloor_{q_2} \\
&= \left[\mathbf{v}_1 - \mathbf{v}_2 \right]_{q_2} \\
&= \left\lfloor \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1 \right) \right\rfloor + \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_2\mathbf{e} \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right]_{q_1} \right\rfloor - \left[\left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2 \right) \mathbf{c}_2\boldsymbol{\alpha}_2 + \mathbf{h}\boldsymbol{\alpha}_2 + \mathbf{h}_2\mathbf{e} \cdot \mathbf{g}_{2,q_2}^{-1} \right]_{q_2} \right\rfloor_{q_2} \\
&= \left\lfloor \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_t \right) \right\rfloor + \mathbf{h}_2\mathbf{e} \cdot \left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{g}_2^{-1} \right]_{q_1} \right\rfloor + \boldsymbol{\delta}_2 - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2 \right) \mathbf{c}_2\boldsymbol{\alpha}_2 - \mathbf{h}\boldsymbol{\alpha}_2 - \mathbf{h}_2\mathbf{e} \cdot \mathbf{g}_{2,q_2}^{-1} \right\rfloor_{q_2} \\
&= \left\lfloor \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1 \right) \right\rfloor + \boldsymbol{\delta}_2 - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2 \right) \mathbf{c}_2\boldsymbol{\alpha}_2 - \mathbf{h}\boldsymbol{\alpha}_2 \right\rfloor_{q_2}
\end{aligned}$$

,

where $\left\| \boldsymbol{\delta}_2 \right\| < l_1(\mathbf{h}_2\mathbf{e})$ by Lemma 3.6.

For our choice of parameter, $\left\| \mathbf{r} \right\| < q_2^{1/8}$, $\left\| \mathbf{c}_1 \right\| < q_2^{1/8}$, $\left\| \mathbf{c}_2 \right\| < q_2^{1/8}$, and $\left\| \mathbf{h} \right\| < n^{O(1)} q_2^{1/2}, \left\| \mathbf{h}_1 \right\| < n^{O(1)} q_2^{1/2}, \left\| \mathbf{h}_2 \right\| < n^{O(1)} q_2^{1/2}$. Thus, $\mathbf{v}$ is not reduced modulo q_2. That is,

$$\begin{aligned}\|\mathbf{v}\| &= \left\| \left[\left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1\right) \right\rfloor + \boldsymbol{\delta}_2 - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2\right)\mathbf{c}_2\boldsymbol{\alpha}_2 - \mathbf{h}\boldsymbol{\alpha}_2 \right]_{q_2} \right\| \\ &\le 2^{-\lambda} \cdot \|\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1\| + \|\boldsymbol{\delta}_2\| + \left\|\left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2\right)\mathbf{c}_2\boldsymbol{\alpha}_2\right\| + \|\mathbf{h}\boldsymbol{\alpha}_2\| \\ &\le 6n^{O(1)} \cdot q_2^{1/8} \cdot n^{O(1)} \cdot q_2^{1/2} \\ &< q^{3/4}\end{aligned} .$$

If $\mathbf{u}_1$ is a level-κ encoding of non-zero element for $\langle \mathbf{g} \rangle$. Namely $\mathbf{u}_1 = \left[\frac{\mathbf{r}\mathbf{g} + \mathbf{d}}{\mathbf{z}_1^\kappa} \right]_q$ with $\mathbf{d} \neq 0 \bmod \mathbf{g}$ and $\|\mathbf{d}\| \le \|\mathbf{g}\|$. Thus,

$$\begin{aligned}\mathbf{v} &= \left[\mathbf{v}_1 - \mathbf{v}_2\right]_{q_2} \\ &= \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1}\right]_{q_1} \right]_{q_1} \right\rfloor + \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1\right) \right\rfloor + \boldsymbol{\delta}_2 - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2\right)\mathbf{c}_2\boldsymbol{\alpha}_2 - \mathbf{h}\boldsymbol{\alpha}_2 \right]_{q_2} \\ &\ge \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1}\right]_{q_1} \right]_{q_1} \right\rfloor \right]_{q_2} - \left[\left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r} + \mathbf{h}_2\mathbf{c}_1\right) \right\rfloor + \boldsymbol{\delta}_2 - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2\right)\mathbf{c}_2\boldsymbol{\alpha}_2 - \mathbf{h}\boldsymbol{\alpha}_2 \right]_{q_2}, \\ &\ge q_2^{1-\varepsilon} - q_2^{3/4} \\ &\ge q_2^{1-\varepsilon'}\end{aligned}$$

where $\varepsilon, \varepsilon'$ are arbitrary small positive constants.

By Lemma 4 in [GGH13], $\left\| \left[\mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1}\right]_{q_1} \right]_{q_1} \right\| \approx q_1$. Thus, $\|\mathbf{v}\| \approx q_2$. □

Lemma 3.9 Given two level-κ encodings $\left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right), \left(\mathbf{u}_{1,2}, \mathbf{u}_{2,2}\right)$, suppose that $\mathbf{u}_{1,1}, \mathbf{u}_{1,2}$ encode same plaintext, then

$$\mathrm{Ext}\left(\mathrm{par}, \left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right)\right) = \mathrm{Ext}\left(\mathrm{par}, \left(\mathbf{u}_{1,2}, \mathbf{u}_{2,2}\right)\right).$$

Proof. Let $\mathbf{u}_{1,s} = \left[\frac{\mathbf{r}_s\mathbf{g} + \mathbf{d}}{\mathbf{z}_1^\kappa} \right]_{q_1} = \left[\frac{\mathbf{c}_{1,s}\mathbf{f} + \mathbf{e}_s}{\mathbf{z}_1^\kappa} \right]_{q_1}$, $s \in [\![2]\!]$ so that $\|\mathbf{r}_s\mathbf{g} + \mathbf{d}\| \le q_2^{1/8}$, and $\mathbf{u}_{2,s} = \left[\frac{\mathbf{c}_{2,s}\mathbf{f} + \mathbf{e}_s}{\mathbf{z}_2^\kappa} \right]_{q_2}$, $s \in [\![2]\!]$. Thus, we have

$$\begin{aligned}\mathbf{v}^{(s)} &= \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{u}_{1,s} \cdot \mathbf{p}_{zt,1}\right]_{q_1} \right\rfloor - \left[\mathbf{u}_{2,s} \cdot \mathbf{p}_{zt,2}\right]_{q_2} \right]_{q_2} \\ &= \left[\mathbf{v}_{1,s} - \mathbf{v}_{2,s}\right]_{q_2} \\ &= \left[\left\lfloor \frac{q_2}{q_1} \cdot \left[\mathbf{h}_1\mathbf{d} \cdot \left[\mathbf{g}_1^{-1}\right]_{q_1} \right]_{q_1} \right\rfloor + \left\lfloor \frac{q_2}{q_1} \cdot \left(\mathbf{h}_1\mathbf{r}_s + \mathbf{h}_2\mathbf{c}_{1,s}\right) \right\rfloor + \boldsymbol{\delta}_{2,s} - \left(\mathbf{h}_2 + \mathbf{h}\mathbf{g}_2\right)\mathbf{c}_{2,s}\boldsymbol{\alpha}_{2,s} - \mathbf{h}\boldsymbol{\alpha}_{2,s} \right]_{q_2}\end{aligned} .$$

By Lemma 3.8, we have $\mathbf{v}^{(1)} \approx \mathbf{v}^{(2)}$ when $\mathbf{d} \neq 0 \bmod \mathbf{g}_1$. Thus, the equality holds. □

3.3 Security

Consider the following security experiment:

(1) $\text{par} \leftarrow \text{InstGen}(1^{\lambda}, 1^{\kappa})$

(2) For $l = 0$ to κ:

Sample $\mathbf{d}_l \leftarrow D_{\mathbb{Z}^n,\sigma'}$, $\mathbf{r}_{l,i} \leftarrow D_{\mathbb{Z}^n,\sigma^*}$;

Generate level-1 encoding $\mathbf{u}_{t,l} = \left[\mathbf{d}_l \mathbf{y}_t + \sum_{i=1}^{\tau} \mathbf{r}_{l,i} \mathbf{x}_{t,i}\right]_{q_t}$, $t \in [\![2]\!]$.

(3) Set $\mathbf{u}_t = \left[\prod_{l=1}^{\kappa} \mathbf{u}_{t,l}\right]_{q_t}$, $t \in [\![2]\!]$.

(4) Set $\mathbf{v}_C = \mathbf{v}_D = \text{Ext}\left(\text{par}, \left(\left[\mathbf{d}_0 \mathbf{u}_{1,1}\right]_{q_1}, \left[\mathbf{d}_0 \mathbf{u}_{2,1}\right]_{q_2}\right)\right)$.

(5) Sample $\mathbf{r}_0 \leftarrow D_{\mathbb{Z}^n,\sigma'}$ and set $\mathbf{v}_R = \text{Ext}\left(\text{par}, \left(\left[\mathbf{r}_0 \mathbf{u}_{1,1}\right]_{q_1}, \left[\mathbf{r}_0 \mathbf{u}_{2,1}\right]_{q_2}\right)\right)$.

Definition 3.10 (ext-GCDH/ext-GDDH). According to the security experiment, the ext-GCDH and ext-GDDH are defined as follows:

Level-κ extraction CDH (ext-GCDH): Given $\left\{\text{par}, \left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right), \cdots, \left(\mathbf{u}_{1,\kappa}, \mathbf{u}_{2,\kappa}\right)\right\}$, output a level-$\kappa$ extraction encoding $\mathbf{w} \in R_{q_2}$ such that $\left\|\left[\mathbf{v}_C - \mathbf{w}\right]_{q_2}\right\|_{\infty} \le q_2^{3/4}$.

Level-κ extraction DDH (ext-GDDH): Given $\left\{\text{par}, \left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right), \cdots, \left(\mathbf{u}_{1,\kappa}, \mathbf{u}_{2,\kappa}\right), \mathbf{v}\right\}$, distinguish between $D_{ext-GDDH} = \left\{\text{par}, \left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right), \cdots, \left(\mathbf{u}_{1,\kappa}, \mathbf{u}_{2,\kappa}\right), \mathbf{v}_D\right\}$ and $D_{ext-RAND} = \left\{\text{par}, \left(\mathbf{u}_{1,1}, \mathbf{u}_{2,1}\right), \cdots, \left(\mathbf{u}_{1,\kappa}, \mathbf{u}_{2,\kappa}\right), \mathbf{v}_R\right\}$.

3.4 Cryptanalysis

In this section, we give easily computable some quantities in our construction, and analyze possible attacks using these quantities.

3.4.1 Easily computable quantities

The encodings $\mathbf{y}_1, \mathbf{x}_{1,i}$ in the public parameters are same as that of GGH for the ideal lattice $\langle \mathbf{g}_1 \rangle$. However, the zero testing parameter $\mathbf{p}_{zt,1}$, which is different from one of GGH, is added a term $\mathbf{z}_1^{\kappa} \mathbf{h}_2 \left[\mathbf{g}_2^{-1}\right]_{q_1}$. As a result, the encodings $\mathbf{x}_{1,i}$ of zero for $\langle \mathbf{g}_1 \rangle$ are not any more encoding of zero for $\langle \mathbf{g}_2 \rangle$. Although the non-zero plaintexts encoded by $\mathbf{y}_1, \mathbf{x}_{1,i}$ are one-to-one corresponding to ones encoded by $\mathbf{y}_2, \mathbf{x}_{2,i}$ for $\langle \mathbf{g}_2 \rangle$, they cannot be subtracted to obtain encoding of zero for $\langle \mathbf{g}_2 \rangle$. This is because: (1) $\mathbf{z}_1$ using as level number of encoding is not equal to $\mathbf{z}_2$; (2) the modulo q_1 in the first type encodings is different from q_2 in the second type encodings. Thus, to remove the non-zero level-0 encodings in $\mathbf{y}_1, \mathbf{x}_{1,i}$ for $\langle \mathbf{g}_2 \rangle$, one must switch from q_1 to q_2 and use zero testing parameter $\mathbf{p}_{zt,2}$. Thus, one can only get easily computable quantities in the following form.

Given a level-k encoding $\left(\mathbf{u}_1, \mathbf{u}_2\right)$ with $1 \le k < \kappa$, we can compute using par to get

$$\mathbf{v}=\left[\left\lfloor \frac{q_2}{q_1}\left[\mathbf{u}_1\cdot(\mathbf{x}_{1,i})^j\cdot(\mathbf{y}_1)^{\kappa-k-j}\cdot\mathbf{p}_{zt,1}\right]_{q_1}\right\rfloor-\left[\mathbf{u}_2\cdot(\mathbf{x}_{2,i})^j\cdot(\mathbf{y}_2)^{\kappa-k-j}\cdot\mathbf{p}_{zt,2}\right]_{q_2}\right]_{q_2}.$$

It is easy to see that $\mathbf{v}$ is not reduced modulo q_2. However, modulus switching destroys the structure of the ring element in $\left[\mathbf{u}_1\cdot(\mathbf{x}_{1,i})^j\cdot(\mathbf{y}_1)^{\kappa-k-j}\cdot\mathbf{p}_{zt,1}\right]_{q_1}$. As a result, the attack described by Coron et al. [CGH+15] is not applicable for our construction.

The SubM problem. The subgroup membership problem is seemly hard for our construction. Let $\mathbf{g}_1=\mathbf{g}_{1,1}\mathbf{g}_{1,2}$. Given a level-$1$ encoding $(\mathbf{u}_1,\mathbf{u}_2)$ with $\mathbf{u}_1=\left[\frac{\mathbf{w}}{\mathbf{z}_1}\right]_{q_1}$, determine if $\mathbf{w}\in\langle\mathbf{g}_{1,1}\rangle$. Using $\mathbf{v}$, one cannot decide whether $\mathbf{v}$ belongs to $\langle\mathbf{g}_{1,1}\rangle$ regardless of $\mathbf{w}\in\langle\mathbf{g}_{1,1}\rangle$. This is again the result destroying structure of ring element.

The DLIN problem. The decision linear problem is also seemly hard for our construction. For a matrix of $\mathbf{A}=(\mathbf{a}_{i,j})\in R^{w\times w}$, all encoded at level-$k$, $1\le k<\kappa$ form a matrix $\mathbf{T}$, the DLIN problem is to distinguish between rank w and rank $w-1$ for $\mathbf{A}$. Based on the similar reason above, one cannot compute the rank of $\mathbf{A}$ in our encoding scheme.

3.4.2 Hu-Jia Attack

In this section, we show that the Hu-Jia attack [HJ15a] does not work for our construction.

Hu-Jia Attack Description

Their attack includes three steps. The first step generates an equivalent level-0 encoding for a level-1 encoding; the second step computes an equivalent level-0 encoding for the product of several level-0 encodings; the final step transforms an equivalent product level-0 encoding into the shared secret key of MPKE by the modified encoding/decoding.

By analysis, the first step is the key of the Hu-Jia attack. We describe the concrete details of the first step as follows:

(1) Let $\mathrm{par}_0=\left\{q,\mathbf{y}=\left[(1+\mathbf{ag})/\mathbf{z}\right]_q,\mathbf{x}_i=\left[(\mathbf{a}_i\mathbf{g})/\mathbf{z}\right]_q,i\in[\![2]\!],\mathbf{p}_{zt}=\left[(\mathbf{hz}^\kappa)/\mathbf{g}\right]_q\right\}$ be the public parameters of the GGH map. We generate special decodings $\left\{\mathbf{y}^{(1)},\mathbf{x}^{(i)},i=1,2\right\}$, where

$$\mathbf{y}^{(1)}=\left[\mathbf{p}_{zt}\mathbf{y}^{\kappa-1}\mathbf{x}_1\right]_q=\mathbf{h}(1+\mathbf{ag})^{\kappa-1}\mathbf{a}_1,$$

$$\mathbf{x}^{(i)}=\left[\mathbf{p}_{zt}\mathbf{y}^{\kappa-2}\mathbf{x}_i\mathbf{x}_1\right]_q=\mathbf{h}(1+\mathbf{ag})^{\kappa-2}(\mathbf{a}_i\mathbf{g})\mathbf{a}_1,\ i=1,2.$$

Notice that $\mathbf{y}^{(1)},\mathbf{x}^{(i)}$ are not reduced modulo q.

(2) Given a level-1 encoding $\mathbf{u}$, we have $\mathbf{u}=\left[\mathbf{dy}+\mathbf{r}_1\mathbf{x}_1+\mathbf{r}_2\mathbf{x}_2\right]_q$, where $\mathbf{d}$ is secret level-0 encoding, and $\mathbf{r}_1,\mathbf{r}_2$ random noise elements.

Compute special decoding

$$\mathbf{v}=\left[\mathbf{p}_{zt}\mathbf{u}\mathbf{y}^{\kappa-2}\mathbf{x}_1\right]_q=\mathbf{dy}^{(1)}+\mathbf{r}_1\mathbf{x}^{(1)}+\mathbf{r}_2\mathbf{x}^{(2)}.$$

Since $\mathbf{v}$ is not reduced modulo q, then compute

$$\mathbf{v}\bmod\mathbf{y}^{(1)}=(\mathbf{r}_1\mathbf{x}^{(1)}\bmod\mathbf{y}^{(1)}+\mathbf{r}_2\mathbf{x}^{(2)}\bmod\mathbf{y}^{(1)})\bmod\mathbf{y}^{(1)}.$$

(3) Given $\mathbf{v}\bmod\mathbf{y}^{(1)}$ and $\left\{\mathbf{x}^{(1)}\bmod\mathbf{y}^{(1)},\mathbf{x}^{(2)}\bmod\mathbf{y}^{(1)}\right\}$, we get $\mathbf{v}'=\mathbf{v}\bmod\mathbf{y}^{(1)}\in\langle\mathbf{x}^{(1)},\mathbf{x}^{(2)}\rangle$ such that $(\mathbf{v}-\mathbf{v}')\bmod\mathbf{y}^{(1)}=0$. Let $\mathbf{v}'=\mathbf{r}_1'\mathbf{x}^{(1)}+\mathbf{r}_2'\mathbf{x}^{(2)}$.

(4) Compute $\mathbf{d}^{(0)}=(\mathbf{v}-\mathbf{v}')/\mathbf{y}^{(1)}$ over $\mathbb{k}=\mathbb{R}[x]/\langle x^n+1\rangle$ such that the quotient $\mathbf{d}^{(0)}\in R$. By arranging, we obtain

$$\mathbf{d}^{(0)} = (\mathbf{v} - \mathbf{v}') / \mathbf{y}^{(1)}$$
$$= \mathbf{d} + ((\mathbf{r}_1 - \mathbf{r}_1')\mathbf{a}_1 + (\mathbf{r}_2 - \mathbf{r}_2')\mathbf{a}_2)\mathbf{g} / (1 + \mathbf{ag}) .$$

Again since $\mathbf{g}$ and $1+\mathbf{ag}$ are co-prime, we get $\mathbf{d} - \mathbf{d}^{(0)} \in \langle \mathbf{g} \rangle$. Thus, $\mathbf{d}^{(0)}$ is an equivalent level-0 encoding of $\mathbf{d}$. Although $\left\| \mathbf{d}^{(0)} \right\|$ is not small, Hu and Jia [HJ15a] controlled the size of $\mathbf{d}^{(0)}$ by using $\mathbf{x}^{(i)} \in \langle \mathbf{g} \rangle$.

Non-applicabiltiy of Hu-Jia Attack

(1) Let $\text{par} = \left\{ \left\{ q_t, \mathbf{y}_t, \left\{ \mathbf{x}_{t,i} \right\}_{i \in [\![\tau]\!]}, \mathbf{p}_{zt,t} \right\}_{t \in [\![2]\!]} \right\}$ be the public parameters of our construction. Similarly, we generate special decodings $S = \left\{ \mathbf{y}^{(1)}, \left\{ \mathbf{x}^{(i)} \right\}_{i \in [\![\tau]\!]} \right\}$ as follows:

$$\mathbf{y}^{(1)} = \left[\left\lfloor q_2 / q_1 \cdot \left[\mathbf{x}_{1,1} \cdot (\mathbf{y}_1)^{\kappa-1} \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{x}_{2,1} \cdot (\mathbf{y}_2)^{\kappa-1} \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2},$$
$$\mathbf{x}^{(i)} = \left[\left\lfloor q_2 / q_1 \cdot \left[\mathbf{x}_{1,i}\mathbf{x}_{1,1} \cdot (\mathbf{y}_1)^{\kappa-2} \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{x}_{2,i}\mathbf{x}_{2,1} \cdot (\mathbf{y}_2)^{\kappa-2} \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2}.$$

Notice that $\mathbf{y}^{(1)}, \mathbf{x}^{(i)}$ are not reduced modulo q.

(2) Given a level-1 encoding $(\mathbf{u}_1, \mathbf{u}_2)$ with $\mathbf{u}_t = \left[\mathbf{d} \cdot \mathbf{y}_t + \sum_{i=1}^{\tau} \mathbf{r}_i \cdot \mathbf{x}_{t,i} \right]_q, t \in [\![2]\!]$, we compute special decoding

$$\mathbf{v} = \left[\left\lfloor q_2 / q_1 \cdot \left[\mathbf{u}_1\mathbf{x}_{1,1} \cdot (\mathbf{y}_1)^{\kappa-2} \cdot \mathbf{p}_{zt,1} \right]_{q_1} \right\rfloor - \left[\mathbf{u}_2\mathbf{x}_{2,1} \cdot (\mathbf{y}_2)^{\kappa-2} \cdot \mathbf{p}_{zt,2} \right]_{q_2} \right]_{q_2} \neq \mathbf{d} \cdot \mathbf{y}^{(1)} + \sum_{i=1}^{\tau} \mathbf{r}_i \cdot \mathbf{x}^{(i)} .$$

Although $\mathbf{v}$ is not reduced modulo q, one cannot find an equivalent level-0 encoding encoded by $(\mathbf{u}_1, \mathbf{u}_2)$ using the Hu-Jia method. Because $\mathbf{x}^{(i)} \notin \langle \mathbf{g}_1 \rangle$ and the probability that $\mathbf{y}^{(1)}, \mathbf{x}^{(i)}$ are co-prime is almost 1. Namely, $\mathbf{y}^{(1)}, \mathbf{x}^{(i)}$ have not common factor. This is different from the case of the original GGH described by Hu and Jia [HJ15a]. Moreover, one cannot efficiently solve $\mathbf{d}, \mathbf{r}_i$ given $\mathbf{y}^{(1)}, \mathbf{x}^{(i)}, \mathbf{v}$. This is because $\mathbf{v} \neq \mathbf{d} \cdot \mathbf{y}^{(1)} + \sum_{i=1}^{\tau} \mathbf{r}_i \cdot \mathbf{x}^{(i)}$.

Thus, one cannot find an equivalent level-0 encoding encoded by $(\mathbf{u}_1, \mathbf{u}_2)$. Namely, the Hu-Jia attack is prevented in our construction.

3.4.3 Cheon-Lee Attack

The Cheon-Lee attack [CL15] for the GGH map consists of three steps. The first step is find a basis of secret ideal lattice $\langle \mathbf{g}_1 \rangle$. The second step is to find the shortest vector of $\langle \mathbf{g}_1 \rangle$ using HNF. The third step is to apply a lattice reduction algorithm on reduced dimension to solve the GDDH on the GGH map.

However, one cannot yield a basis of $\langle \mathbf{g}_1 \rangle$ using the public parameters in our construction. Thus, The Cheon-Lee attack does not work in our construction.

4 Applications

In the following, we describe two applications using our construction: the MPKE protocol and the instance of witness encryption.

4.1 MPKE Protocol

$\textbf{\textit{Setup}}(1^{\lambda},1^{N})$. Output $(\text{par}) \leftarrow \text{InstGen}(1^{\lambda},1^{\kappa})$ as the public parameters.

$\textbf{\textit{Publish}}(\text{par}, j)$. The j-th party samples $\mathbf{d}_j \leftarrow D_{\mathbb{Z}^n,\sigma'}$, $\mathbf{r}_{j,i} \leftarrow D_{\mathbb{Z}^n,\sigma^*}, i \in [\![\tau]\!]$, publishes the public key $\mathbf{u}_{t,j} = \left[\mathbf{d}_j \cdot \mathbf{y}_t + \sum_{i=1}^{\tau} \mathbf{r}_{j,i} \cdot \mathbf{x}_{t,i}\right]_{q_t}, t \in [\![2]\!]$ and remains $\mathbf{d}_j$ as the secret key.

$\textbf{\textit{KeyGen}}\left(\text{par}, j, \mathbf{d}_j, \left\{\left(\mathbf{u}_{1,k}, \mathbf{u}_{2,k}\right)\right\}_{k\neq j}\right)$. The j-th party computes $\mathbf{c}_{t,j} = \left[\prod_{k\neq j} \mathbf{u}_{t,k}\right]_{q_t}, t \in [\![2]\!]$ and extracts the common secret key $sk = \text{Ext}\left(\text{par}, \left(\left[\mathbf{d}_j \mathbf{c}_{1,j}\right]_{q_1}, \left[\mathbf{d}_j \mathbf{c}_{2,j}\right]_{q_2}\right)\right)$.

Theorem 5.1 Suppose the ext-GCDH/ext-GDDH defined in Section 3.3 is hard, then our construction is one round multipartite Diffie-Hellman key exchange protocol.

4.2 Witness Encryption

4.2.1 Construction

Garg, Gentry, Sahai, and Waters [GGSW13] constructed an instance of witness encryption based on the NP-complete 3-exact cover problem and the GGH map. However, Hu and Jia [HJ15a] have broken the GGH-based WE. In this section, we present a new construction of WE based on our new multilinear map.

3-Exact Cover Problem [GGH13, Gol08] Given a collection Set of subsets $T_1, T_2, ..., T_\pi$ of $[\![K]\!] = \{1, 2, ..., K\}$ such that $K = 3\theta$ and $|T_i| = 3$, find a 3-exact cover of $[\![K]\!]$. For an instance of witness encryption, the public key is a collection Set and the public parameters par in our construction, the secret key is a hidden 3-exact cover of $[\![K]\!]$.

$\textbf{\textit{Encrypt}}(1^{\lambda}, \text{par}, M)$**:**

(1) For $k \in [\![K]\!]$, sample $\mathbf{d}_k \leftarrow D_{Z^n,\sigma'}$, $\mathbf{r}_{k,i} \leftarrow D_{\mathbb{Z}^n,\sigma^*}, i \in [\![\tau]\!]$ and generate level-1 encodings $\mathbf{u}_{t,k} = \left[\mathbf{d}_k \cdot \mathbf{y}_t + \sum_{i=1}^{\tau} \mathbf{r}_{k,i} \cdot \mathbf{x}_{t,i}\right]_{q_t}, t \in [\![2]\!]$.

(2) Compute $\mathbf{u}_t = \left[\prod_{k=1}^{K} \mathbf{u}_{t,k}\right]_{q_t}, t \in [\![2]\!]$ and $sk = \text{Ext}\left(\text{par}, \left(\mathbf{u}_1, \mathbf{u}_2\right)\right)$, and encrypt a message M into ciphertext C.

(3) For each element $T_j = \{j_1, j_2, j_3\} \in Set$, sample $\mathbf{r}_{T_j,i} \leftarrow D_{\mathbb{Z}^n,\sigma^*}, i \in [\![\tau]\!]$, and generate a level-3 encoding $\mathbf{u}_{t,T_j} = \left[\mathbf{u}_{t,j_1}\mathbf{u}_{t,j_2}\mathbf{u}_{t,j_3} + \sum_{i=1}^{\tau} \mathbf{r}_{T_j,i}(\mathbf{x}_{t,i})^3\right]_{q_t}, t \in [\![2]\!]$.

(4) Output the ciphertext C and all level-3 encodings $E = \left\{\left(\mathbf{u}_{1,T_j}, \mathbf{u}_{2,T_j}\right), T_j \in Set\right\}$.

$\textbf{\textit{Decrypt}}(C, E, W)$:

(1) Given C, E and a witness set W, compute $\mathbf{u}_t = \left[\prod_{T_j \in W} \mathbf{u}_{t,T_j}\right]_{q_t}$.

(2) Generate $sk = \text{Ext}\left(\text{par}, \left(\mathbf{u}_1, \mathbf{u}_2\right)\right)$, and decrypt C to the message M.

Similar to [GGSW13], the security of our construction depends on the hardness assumption of the Decision Graded Encoding No-Exact-Cover.

Theorem 5.2 Suppose that the Decision Graded Encoding No-Exact-Cover is hard. Then our

construction is a witness encryption scheme.

4.2.2 Hu-Jia Attacks

Since $\mathbf{u}_{t,T_j} = \left[\mathbf{u}_{j_1}\mathbf{u}_{j_2}\mathbf{u}_{j_3} + \sum_{i=1}^{\tau} \mathbf{r}_{T_j,i}(\mathbf{x}_{t,i})^3 \right]_{q_t}$, $t \in [\![2]\!]$ is a level-3 encoding in our encoding method, one cannot obtain $\mathbf{u}_{t,T_i} = \left[\mathbf{u}_{t,T_j}\mathbf{u}_{t,T_k}(\mathbf{u}_{t,T_l})^{-1} \right]_{q_t}$ when $T_i = T_j \bigcup T_k - T_l$. As a result, the Hu-Jia attacks [HJ15a, HJ15b] are prevented in our new construction.

5 Conclusion

In this paper, we describe a new variant of GGH, which supports the applications for public tools of encoding in the original GGH, such as MPKE and WE. Using modulus switching and modifying zero testing parameters, our construction introduces new noise term to avoid weakness of the GGH map. As a result, our new construction not only prevents all known attacks, but also seemly supports the hardness assumption of the SubM problem and the DLIN problem.

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