

Extended Truncated-differential Distinguishers on Round-reduced AES

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Abstract. Distinguishers on round-reduced AES have attracted considerable attention in the recent years. While the number of rounds covered in key-recovery attacks did not increase, subspace, yoyo, mixture-differential, and multiple-of- n cryptanalysis advanced the understanding of the properties of the cipher.

For substitution-permutation networks, integral attacks are a suitable target for extension since they usually end after a linear layer sums several subcomponents. Based on results by Patarin, Chen et al. already observed that the expected number of collisions for a sum of permutations differs slightly from that for a random primitive. Though, their target remained lightweight primitives.

The present work illustrates how the well-known integral distinguisher on three-round AES resembles a sum of PRPs and can be extended to truncated-differential distinguishers over 4 and 5 rounds. In contrast to previous distinguishers by Grassi et al., our approach allows to prepend a round that starts from a diagonal subspace. We demonstrate how the prepended round can be used for key recovery with a new differential key-recovery attack on six-round AES. Moreover, we show how the prepended round can also be integrated to form a six-round distinguisher. For all distinguishers and the key-recovery attack, our results are supported by implementations with Cid et al.'s established Small-AES version. While the distinguishers do not threaten the security of the AES, they try to shed more light on its properties.

Keywords: Cryptanalysis · block cipher · AES

1 Introduction

During the previous two decades, the Advanced Encryption Standard (AES) [Nat01] has withstood vast amounts of cryptanalysis. Besides the biclique-based accelerated exhaustive search [BKR11]¹, the best-known attacks on AES-128 in the secret-key model cover seven rounds, as had been the state of the art close after its announcement [FKL⁺00]. However, the community's efforts led to attacks with considerably reduced resources. Among the best attacks in terms of the number of rounds [BLNS18, DFJ13, FKL⁺00, MDRM10], the Demirci-Selçuk-based [DS08] meet-in-the-middle attacks by Derbez et al. possess the lowest time and data complexities for more than half a decade [DFJ13].²

¹Instead of exploiting dedicated properties of a given cipher, biclique-based attacks [BKR11] usually represent a general approach to speed up brute force since their outer loop iterates over all possible keys. Some works consider it therefore as accelerated exhaustive search instead of an attack cf. [DDKS13].

²The property exploited by Demirci and Selçuk extended the earlier collision attack by Gilbert and Minier [GM00]

1.1 Distinguishers on Round-reduced AES

Although the number of rounds covered in key-recovery attacks did not increase since then, the recent years were filled with research on distinguishers on round-reduced AES that significantly raised the understanding of the cipher’s components. This direction appears promising – metaphorically, it is comparable to heuristics that sometimes also have to leave a local optimum to improve in the long run.

Negative Results paved a rocky start for the search for new distinguishers. Sun et al. [SLG⁺16b] proved the absence of impossible differentials over more than four rounds for the AES *structure*, which was tantamount with the absence of integrals or zero-correlation attacks over more than four rounds [SLR⁺15].³ Since [SLG⁺16b] ignored the details of the AES S-box and its key schedule, there remained a spark of hope for longer distinguishers of those kinds. The works by Wang and Jin [WJ18, WJ19] extinguished this spark. Under the assumption of independent random round keys, they showed the absence of any (truncated) impossible differentials over more than four-round of the AES taking also the details of the S-box into account.

Key-dependent Distinguishers. Despite the negative results, a series of works has been focusing on novel properties for distinguishers on fewer rounds than the best known attacks. First, a number of key-dependent distinguishers were crafted, e.g., the chosen-ciphertext zero-correlation hull on five rounds by Sun et al. [SLG⁺16a] exploited a known difference in two key bytes to cover five rounds. While their distinguisher required the full codebook when converted to the single-key model, it re-ignited the community’s efforts on analyzing round-reduced AES.

Subsequent works improved on their result and proposed further key-dependent results. In [GRR16], Grassi et al. reconsidered the distinguisher by Sun et al. and derived a key-dependent chosen-plaintext distinguisher with lower complexity. Moreover, differential, linear, and integral distinguishers already implicitly exploited that texts or tuples thereof had a larger probability of lying in (cosets of) certain subspaces, but were often formulated in a complicated manner. Grassi et al. unified them in their notion of subspace trails. In particular for the AES, subspace trails allow to describe diagonal, column, inverse-diagonal, and mixed (results of applying MixColumns to an inverse-diagonal subspace) subspaces quite elegantly.

Grassi proposed further key-dependent distinguishers in [Gra18a]. Starting from the result by Sun et al., Hu et al. [HCGW18] improved the complexity and transformed it into a chosen-plaintext attack, and derived impossible-differential attacks from it. Cui et al. [CCM⁺18] later reduced the complexity. Those considered attacks on the AES with a secret S-box [TKKL15]. Among the attacks on round-reduced AES with a secret S-box, the impossible-differential distinguisher on five rounds by Hu et al. [HCGW18] possesses the lowest time complexity of $2^{80.5}$ and $2^{76.5}$ chosen ciphertexts.

Key-independent Distinguishers. Besides the key-dependent results, several powerful key-independent distinguishers have been proposed recently. In a series of works [GRR16, GRR17, GR18, Gra18b], Grassi et al. outlined novel observations and distinguishers on five-round AES. At their core, [GRR17] observed a strong property dubbed multiple-of- n : starting from a structure of a diagonal space, the number of different ciphertext pairs that belonged to the same coset of a mixed space was always a multiple of n after five rounds. Boura et al. [BCC19] revisited the multiple-of- n property and derived similar distinguishers for further AES-like primitives.

³Sun et al. [SLR⁺15] showed that the existence of an impossible differential on a cipher \mathcal{E} implies the existence of affine layers \mathcal{A}_1 and \mathcal{A}_2 such that there is an integral and a zero-correlation distinguisher on $\mathcal{A}_2 \circ \mathcal{E} \circ \mathcal{A}_1$. Thus, a result on one of those kinds of distinguishers also implies results on the others.

Mixture-differential Distinguishers. In [Gra17, Gra18b], Grassi proposed mixture-differential distinguishers. At their core, they exploited the following property: let (P, P') denote a pair of plaintexts that differ in two bytes (x, y) and (x', y') in the same column, and whose corresponding ciphertexts (C, C') have a difference in a certain (mixed) subspace after four rounds. Then, the difference of the ciphertext pair (C'', C''') that corresponds to the plaintexts (P'', P''') with byte values (x, y') and (x', y) (i.e., which “mixes” the original pair of plaintexts), respectively, will also lie in a coset of the same subspace. This pair of plaintext-ciphertext pairs is called a couple. In [Gra18b], Grassi published the efficient four-round distinguisher and a five-round key-recovery attack from [Gra17]. Soon upon, Bar-On et al. [BODK⁺18] improved the key-recovery to the attack on five-round AES with the lowest data and time complexity known so far.

In [Gra19], Grassi further published a probabilistic mixture-differential distinguisher on five rounds as well as a key-recovery attack on six rounds from [Gra17]. The latter version also contained further threshold as well as an impossible mixture-differential distinguishers. The probabilistic distinguisher exploited a tiny difference between two expectations: the expected number of sets with at least one pair whose both ciphertexts belong to the same coset of a subspace is a little lower for five-round AES than for a random permutation. Grassi’s threshold distinguisher exploited different expectations between sets of couples. Each set was formed by couples whose both plaintexts mixed the byte values in two diagonals, and differed only in those two diagonals. This distinguisher exploited a complex fact: Grassi counted the expected number of sets whose number of couples where both ciphertexts belong to the same coset of a subspace is higher than a threshold. This number is higher for five-round AES than it is for a random permutation. Finally, he considered an impossible mixture-differential distinguisher. Specifically defined sets cannot have at least one pair whose both ciphertexts belong to the same coset of a certain subspace (i.e., for five-round AES, at least one set does not have the corresponding property). Again, this attack considered sets formed by couples whose plaintexts mixed the byte values of the first two diagonals.

To the best of our knowledge, the best previously published distinguishers up to now on round-reduced AES-128 in terms of minimal complexity are the exchange/yoyo-based proposals by Rønjom and Bardeh [BR19b, BR19c]. They extended their earlier yoyo attacks from [RBH17] and reduced the complexity to about 2^{88} XORs and chosen plaintexts. Bardeh [Bar19] recently proposed a variant with adaptively chosen ciphertexts with 2^{83} data and time complexity. Table 1 provides a summary of existing distinguishers on five or more rounds of AES-128; note that many results hold for diverse versions of the AES; however, the focus of this work resides in the 128-bit variant.

1.2 From Integral To Truncated-differential Distinguishers

Integral Distinguishers map multi-sets of inputs that iterate over all values to multi-sets of outputs that are balanced. Traditionally, the properties of bits or bytes are either constant (\mathcal{C}), iterate over all values (\mathcal{A}), are balanced (\mathcal{B}), or unknown (\mathcal{U}). A traditional integral distinguisher usually ends directly before all parts of the state become unknown. For the AES, the three-round integral distinguisher [DKR97] that maps sets of a single active byte to a set of states where each byte is balanced after three rounds is understood well, and so is the extension to a distinguisher on four rounds [DKR97, FKL⁺00] that prepends a round and starts from an active diagonal.

Probabilistic Integrals. Wang et al. [WCC⁺16] proposed so-called statistical integrals. Integral distinguishers map a structure of 2^s inputs that iterate over all values to b pairwise disjoint sets of t output bits each that are uniformly distributed. Statistical integral

Table 1: Existing secret-key distinguishers on five and more rounds of the AES-128, ordered by rounds (descending) then time (ascending). MAs = memory accesses; Encs. = r -round AES encryptions; CP = chosen plaintexts; (A)CC = (adaptive) chosen ciphertexts; MD = mixture differential; TD = truncated differential. ⁽¹⁾: We employ the updated parameters from [GR19]. The original distinguishers had been proposed in [GR18].

Attack Type	Time	Data	Ref.
Five Rounds			
Integral	2^{128}	XORs	2^{128} CC [SLG ⁺ 16a]
Impossible Differential	2^{107}	MAs	$2^{98.2}$ CP [GRR16]
Threshold MD	$2^{98.1}$	MAs	2^{89} CP [Gra17]
Impossible MD	$2^{97.8}$	MAs	2^{82} CP [Gra17]
Truncated differential	$2^{73.3}$	MAs	2^{68} CP Sect. 5
Probabilistic MD	$2^{71.5}$	MAs	2^{52} CP [Gra19, Gra17]
Truncated differential ⁽¹⁾	$2^{52.6}$	MAs	$2^{48.96}$ CP [GR18, GR19]
Variance of TD ⁽¹⁾	$2^{37.6}$	MAs	2^{34} CP [GR18, GR19]
Multiple-of-8	$2^{35.6}$	MAs	2^{32} CP [GRR17]
Yoyo	$2^{26.2}$	XORs	$2^{27.2}$ ACC [BR19a]
Yoyo	$2^{25.8}$	XORs	$2^{26.8}$ ACC [RBH17]
Six Rounds			
Impossible Yoyo	$2^{121.83}$	XORs	$2^{122.83}$ ACC [RBH17]
Truncated differential	$2^{96.52}$	MAs	$2^{89.43}$ CP Sect. 7
Exchange	$2^{88.2}$	Encs.	$2^{88.2}$ CP [BR19b, BR19c]
Exchange	2^{83}	Encs.	2^{83} ACC [Bar19]

attacks exploit that the output values are multivariate-hypergeometrically distributed, in contrast to random permutations whose distribution is close to a multi-normal one. Therefore, the adversary need not reconstruct the full integral of 2^s texts, but can reduce the data requirements to $2^{s-t/2}$. Cui et al. [CSCW17, CCM⁺18] transferred the approach to the AES and AES-like ciphers and exploited multiple structures on input and several integral properties on output. If N_s structures are necessary and consider b pairwise disjoint sets of t -bit output, their approach reduced the data complexity further to $O(\sqrt{N_s/b} \cdot 2^{s-t/2})$. Though, this approach does not directly aim at extending distinguishers, but represents a data-reduction technique instead. The statistical integral distinguisher on five-round AES in [CSCW17, CCM⁺18] is built with slight modification based on the zero-correlation linear hull for five-round decryption of AES by Sun et al. in [SLG⁺16a], but has a great improvement in terms of data and time complexity (from 2^{128} CC to $2^{114.32}$ CC). Moreover, inherited from the distinguisher in [SLG⁺16a], the distinguisher by Cui et al. were in the secret-S-box setting. For Skipjack [WCC⁺16], it nevertheless allowed to extend previous attacks by reducing the number of texts to trace through the cipher.

Many forms of cryptanalysis use biases in the expectation of certain properties to distinguish distributions, e.g. differential or linear cryptanalysis. Nevertheless, biases in other statistical properties such as the standard deviation or variance can also be exploited. In [Pat08, Pat13], Patarin studied the number of pairwise collisions for the sums of multiple permutations in comparison with those from random functions. Follow-up works by Nachev, Marrière or Patarin, and Volte [NMV16, NPV18, NPV14, VNM16a, VNM16b] also employed the standard deviation as distinguishing property. Though, their focus resided mainly on Feistel networks.

Extending Integrals To Truncated-differential Distinguishers. The core observation at the beginning of this work was that an integral distinguisher usually ends with a sum operation. In many Substitution-Permutation Networks (SPNs), the linear layer often consists of a sum of multiple words. At the end of an integral distinguisher, such a sum is equivalent to the sum of words that iterate over all values in the subspace—hence, a sum of permutations. The sum still has a Balanced (i.e., zero-sum) property, which is usually destroyed by the subsequent non-linear layer. As illustrated in [Pat08, Pat13], the number of *collisions* induced by a sum of permutations differs slightly from that of an ideal function. This number of collisions due to the sum of the linear layer will be preserved by the subsequent non-linear S-box operation. Therefore, an integral distinguisher can be extended through the subsequent non-linear operation.

We point out that Chen et al. [CMSZ15] had already considered this approach of Patarin’s analysis for extending integrals of SPNs. They considered Type-II and Nyberg-type Feistel networks and conducted experiments on lightweight ciphers for which they could confirm that this strategy can lead to extended attacks. This work, however, focuses on the AES. We view such extended integral distinguishers as truncated-differential distinguishers, because we consider partial collisions in particular words.

A similar direction was furthermore taken by Heys [Hey14] for extending integrals on BSPN [YTH96], a byte-wise variant of the involutorial SPN of [HT96] with a mixing function of $\circ(0, 1, \dots, 1)$. Besides deterministic integrals similar to those on three- and four-round AES, Heys and Liu also studied a probabilistic extension of the three-round integral. There, they employed an eight-byte version of BSPN with the AES S-box and studied the probability of obtaining a zero sum in the individual bytes after three rounds that naturally allows to be traced through the subsequent round. This approach is interesting as it does not consider the probability that the sum of two texts yields zero, but the sum of an entire δ -set of 256 texts is zero with probability significantly deviating from random. We added similar experiments in Appendix I.

Previous Truncated-differential Distinguishers on the AES. In [GR18], Grassi and Rechberger also considered truncated differentials and exploited smaller statistical differences. The core results are five-round distinguishers that exploit the following property of the AES: A structure of 2^{32} plaintexts from a single diagonal space is mapped to a mixed space through two rounds of AES with probability one; With certain probability, it is mixed again to a single diagonal space after Round 3, from which it is mapped to a mixed space through two further rounds as before. Grassi and Rechberger observed that the probability for the transition in the middle is slightly higher for five-round AES than for a random permutation. Furthermore, they show how to exploit the considerably different variances. Their results are similar to what we aim to in this work. Though, we start from a different angle (from integrals), whereas their work did not depend on Patarin’s result. In contrast, the mean-based distinguishers in [Gra17, Gra19] considered the expectation of couples, whereas we will consider pairs. So, this work starts from Patarin’s results and explores a different path.

Contribution. This work tries to extend the known integral to truncated-differential distinguishers. As a result, it describes a five-round distinguisher from a single byte to a mixed space. Since inputs start from single-byte differences, plaintext structures can form fewer pairs than in e.g., the structures from diagonals as in [GR18]. As a consequence, the data and computational complexity of the five-round distinguisher here are higher than those of the probabilistic distinguishers in [GR18]. However, our proposal allows a straight-forward extension to a six-round key-recovery attack by prepending a round. We present the results of a practical implementation of the five-round distinguisher and the six-round key-recovery attacks with a small-scale variant of the AES. Finally, we

propose a possible extension to a six-round truncated-differential distinguisher and report our results of its implementation for the small variant. The code to all our experiments is freely available online.⁴

Outline. The remainder is structured as follows. Next, Section 2 will revisit the necessary preliminaries, as well as the known results by Patarin on sums of independent permutations and subspaces of the AES by Grassi et al. [GRR16], and briefly recalls the integral distinguishers of round-reduced AES. Section 3 will provide a statistical model. Thereupon, Section 4 will develop from them a four-round truncated distinguisher. Section 5 proceeds similarly for a five-round distinguisher. Section 6 describes a key-recovery attack on six rounds based on our five-round distinguisher. Section 7 derives a six-round distinguisher. We provide a theoretical analysis for all distinguishers for both the AES as well as for Small-AES, and report on our results of experimental verification with the latter. Similarly, we implemented also the key-recovery on six rounds with Small-AES and provide insights to our results. Section 8 discusses and concludes this work.

2 Preliminaries

General Notations. We denote by \mathbb{F}_2 the finite field of characteristic two. We represent functions and variables by upper case letters and indices by lowercase letters, sets by calligraphic letters. We employ typewriter font for hexadecimal values. Let $X, Y \in \mathbb{F}_2^n$ for some positive integer n in the following. Then, we denote by $X \parallel Y$ the concatenation of X and Y , by $X \oplus Y$ their bitwise XOR. For all $X \in \mathbb{F}_2^n$, we index the bits $X = (X_{n-1} \dots X_1 X_0)$ where X_{n-1} is the most significant and X_0 the least significant bit of X . For integers $x \geq y$, we write $X_{x..y}$ as short form of $(X_x X_{x-1} \dots X_y)$. We denote by $\mathbb{E}[X]$ the expectation of a random variable X , by σ_X its standard deviation, and by $\sigma_X^2 = \text{Var}[X]$ its variance. We denote by μ and σ the mean and the standard deviation of a distribution.

Functions and Permutations. For sets \mathcal{X} and \mathcal{Y} , we define $\text{Func}(\mathcal{X}, \mathcal{Y}) \stackrel{\text{def}}{=} \{F|F : \mathcal{X} \rightarrow \mathcal{Y}\}$ and $\text{Perm}(\mathcal{X})$ for the set of all permutations over \mathcal{X} . We define $\text{Func}_n \stackrel{\text{def}}{=} \text{Func}(\mathbb{F}_2^n, \mathbb{F}_2^n)$ and write $\text{Perm}_n = \text{Perm}(\mathbb{F}_2^n)$. Note that $|\text{Func}_n| = (2^n)^{2^n}$ and $|\text{Perm}_n| = (2^n)!$. We call π an ideal permutation (over \mathbb{F}_2^n) if $\pi \leftarrow \text{Perm}_n$, i.e., if it is sampled uniformly at random from Perm_n . Similarly, we call ρ an ideal function (over \mathbb{F}_2^n) if $\rho \leftarrow \text{Func}_n$. For integers $m \leq n$ and arbitrary $X \in \mathbb{F}_2^n$, we define $\text{trunc}_m(X) \stackrel{\text{def}}{=} \text{msb}_m(X) \stackrel{\text{def}}{=} X_{(n-1)..(n-m)}$ to truncate the input x and return only the most-significant (i.e., leftmost) m bits of x .

2.1 Distinguishers for Sums of Permutations

In the following, we recall briefly the results by [Pat08, Pat13]. Given a function set \mathcal{F} , we define by $\text{Gen}(\mathcal{F})$ a function generator that gives access to multiple pairwise independent instances from \mathcal{F} . Let $\pi_1, \dots, \pi_k \leftarrow \text{Perm}_n$ be independent ideal random permutations, and let $\rho \leftarrow \text{Func}_n$. We define a k -sum of permutations as $\Sigma_k[\pi_1, \dots, \pi_k](x) \stackrel{\text{def}}{=} \bigoplus_{i=1}^k \pi_i(x)$ and write Σ_k as short form.

The goal of a k -sum-distinguisher is to distinguish ρ from Σ_k . The setting where the adversary has access to only a single instance of ρ or Σ_k is usually irrelevant in the context of provable security since the query complexity of such distinguisher is close or exceeds $O(2^n)$, i.e., the codebook. Patarin considered a setting where the distinguisher has access to multiple independent functions in the generator $\text{Gen}(\mathcal{F})$ and can ask more

⁴<https://github.com/medsec/expectation-cryptanalysis-on-round-reduced-aes>.

Table 2: Expected numbers of collisions after q queries for the sums of k permutations and distinguishing complexity for $q \simeq 2^n$ from [Pat08].

k	2	3	4
$\mathbb{E}[\mathbf{N}_k]$	$\frac{g\binom{q}{2}}{2^n} \left(1 + \frac{1}{2^n - 1}\right)$	$\frac{g\binom{q}{2}}{2^n} \left(1 - \frac{1}{(2^n - 1)^2}\right)$	$\frac{g\binom{q}{2}}{2^n} \left(1 + \frac{1}{(2^n - 1)^3}\right)$
Complexity	$O(2^{2n})$	$O(2^{4n})$	$O(2^{6n})$

queries. We denote by g the number of available functions from the generator, and by q the number of queries x_i asked to each of the functions.

Let \mathbf{N} be a random variable for the number of output collisions y_i , i.e., $\mathbf{N} = |\{i, j \in [q] : i \neq j \wedge y_i = y_j\}|$. For a function F , we define the random variable \mathbf{N}_F for the number of collisions of F after q queries. So, let \mathbf{N}_ρ be the number of collisions for ρ . Since all responses y_i are sampled independently uniformly at random, it holds that [Pat08]

$$\mathbb{E}[\mathbf{N}_\rho] = \frac{g\binom{q}{2}}{2^n} \quad \text{and} \quad \sigma(\mathbf{N}_\rho) = O\left(\frac{\sqrt{gq}}{\sqrt{2^n}}\right).$$

Let \mathbf{N}_k be a random variable for the number of collisions of Σ_k (\mathbf{N}_k is short for \mathbf{N}_{Σ_k}). Scenario 2 in [Pat08] gives

$$\mathbb{E}[\mathbf{N}_k] = \frac{g\binom{q}{2}}{2^n} \cdot \left(1 + \frac{(-1)^k}{(2^n - 1)^{k-1}}\right) \quad \text{and} \quad \sigma(\mathbf{N}_k) = O\left(\frac{\sqrt{gq}}{\sqrt{2^n}}\right).$$

Patarin argues with Chebyshev's theorem $\Pr[|X - \mu| \geq c\sigma] \leq 1/c^2$, that the distinguishing advantage between the collision distributions of two functions F and G becomes non-negligible if

$$\sigma(\mathbf{N}_F) \ll |\mathbb{E}[\mathbf{N}_F] - \mathbb{E}[\mathbf{N}_G]| \quad \text{and} \quad \sigma(\mathbf{N}_G) \ll |\mathbb{E}[\mathbf{N}_F] - \mathbb{E}[\mathbf{N}_G]|,$$

hold. For example, the sum of $k = 2$ permutations has

$$\mathbb{E}[\mathbf{N}_2] = \frac{g\binom{q}{2}}{2^n} \left(1 + \frac{1}{2^n - 1}\right).$$

Thus, if

$$\frac{\sqrt{gq}}{\sqrt{2^n}} \ll \frac{g\binom{q}{2}}{2^{2n}},$$

both distributions can be distinguished with non-negligible advantage. So, for $q \simeq 2^n$, the adversary can distinguish both settings in $g \geq 2^n$, i.e., the adversary needs $O(2^{2n})$ queries. For k permutations, Patarin showed that the advantage is non-negligible when

$$\frac{\sqrt{gq}}{\sqrt{2^n}} \ll \frac{gq^2}{2^{kn}}.$$

For $q \simeq 2^n$, this yields that $g \geq 2^{(2k-3)n}$ functions are necessary. Table 2 illustrates the expected number of collisions and distinguishing efforts for $k \in \{2, 3, 4\}$.

2.2 The AES-128 and Subspaces

Brief Definition of The AES-128. We assume that the reader is familiar with the details of the AES and provide only a very brief summary here. Details can be found in, e.g.,

[DR02, Nat01]. The AES-128 is a substitution-permutation network that transforms 128-bit inputs through ten rounds, consisting of SubBytes (SB), ShiftRows (SR), MixColumns (MC), and a round-key addition with a round key K^i . Before the first round, a whitening key K^0 is XORed to the state; the final round omits the MixColumns operation. We write S^i for the state after Round i , and $S^i[j]$ for the j -th byte, for $0 \leq i \leq 10$ and $0 \leq j \leq 15$. Though, we will interchangeably also use the indices for a 4×4 -byte matrix, i.e., 0,0 for Byte 0, and 3,3 for Byte 15. So, the byte ordering is given by either

$$\begin{bmatrix} 0 & 4 & 8 & 12 \\ 1 & 5 & 9 & 13 \\ 2 & 6 & 10 & 14 \\ 3 & 7 & 11 & 15 \end{bmatrix} \quad \text{or} \quad \begin{bmatrix} 0,0 & 0,1 & 0,2 & 0,3 \\ 1,0 & 1,1 & 1,2 & 1,3 \\ 2,0 & 2,1 & 2,2 & 2,3 \\ 3,0 & 3,1 & 3,2 & 3,3 \end{bmatrix}.$$

When using two-dimensional indices, we will assume that all indices are taken modulo four, which will make our life easier in sums. $R[K^i] \stackrel{\text{def}}{=} \text{AK}[K^i] \circ \text{MC} \circ \text{SR} \circ \text{SB}$ denotes one application of the round function and denote by $S^{r,\text{SB}}$, $S^{r,\text{SR}}$, and $S^{r,\text{MC}}$ the states in the r -th round directly after the application of SubBytes, ShiftRows, and MixColumns, respectively. Moreover, we will use \mathbf{M} to denote the MixColumns matrix.

Subspaces of The AES. We adopt the notation of subspaces for the AES from Grassi et al. [GRR16]. Let \mathcal{W} denote a vector space and $\mathcal{V} \subseteq \mathcal{W}$ be a subspace. If a is an element of \mathcal{W} , then a coset $\mathcal{V} \oplus a$ of \mathcal{V} in \mathcal{W} is a subset $\mathcal{V} \oplus a = \{v \oplus a \mid v \in \mathcal{V}\}$. We consider vectors and vector spaces over $\mathbb{F}_{2^8}^{4 \times 4}$, and denote by $\{e_{0,0}, \dots, e_{3,3}\}$ the unit vectors of $\mathbb{F}_{2^8}^{4 \times 4}$, i.e., $e_{i,j}$ has a single 1 in the i -th row and j -th column. For a vector space \mathcal{V} and a function $F : \mathbb{F}_{2^8}^{4 \times 4} \rightarrow \mathbb{F}_{2^8}^{4 \times 4}$, we let $F(\mathcal{V}) \stackrel{\text{def}}{=} \{F(v) \mid v \in \mathcal{V}\}$. For a subset $\mathcal{I} \subseteq \{1, 2, \dots, n\}$ and a subset of vector spaces $\{\mathcal{V}_1, \mathcal{V}_2, \dots, \mathcal{V}_n\}$, we define $\mathcal{V}_{\mathcal{I}} \stackrel{\text{def}}{=} \bigoplus_{i \in \mathcal{I}} \mathcal{V}_i$. We adopt the definitions by Grassi et al. of four families of subspaces for the AES for each $i \in \{0, 1, 2, 3\}$:

- the column spaces $\mathcal{C}_{\{i\}} = \langle e_{0,i}, e_{1,i}, e_{2,i}, e_{3,i} \rangle$,
- the diagonal spaces $\mathcal{D}_{\{i\}} = \text{SR}^{-1}(\mathcal{C}_{\{i\}})$,
- the inverse-diagonal spaces $\mathcal{ID}_{\{i\}} = \text{SR}(\mathcal{C}_{\{i\}})$, and
- the mixed spaces $\mathcal{M}_{\{i\}} = \text{MC}(\mathcal{ID}_{\{i\}})$.

For $\mathcal{I} \subseteq \{0, 1, 2, 3\}$, the spaces $\mathcal{C}_{\mathcal{I}}$, $\mathcal{D}_{\mathcal{I}}$, $\mathcal{ID}_{\mathcal{I}}$, and $\mathcal{M}_{\mathcal{I}}$ are defined as

$$\mathcal{C}_{\mathcal{I}} \stackrel{\text{def}}{=} \bigoplus_{i \in \mathcal{I}} \mathcal{C}_{\{i\}}, \quad \mathcal{D}_{\mathcal{I}} \stackrel{\text{def}}{=} \bigoplus_{i \in \mathcal{I}} \mathcal{D}_{\{i\}}, \quad \mathcal{ID}_{\mathcal{I}} \stackrel{\text{def}}{=} \bigoplus_{i \in \mathcal{I}} \mathcal{ID}_{\{i\}}, \quad \text{and} \quad \mathcal{M}_{\mathcal{I}} \stackrel{\text{def}}{=} \bigoplus_{i \in \mathcal{I}} \mathcal{M}_{\{i\}}.$$

Small-AES. Cid et al. [CMR05] proposed small-scale variants of the AES to help cryptanalysts study attacks whose complexity prohibited tests on the full-fledged cipher. We employ the four-bit variant in the following, which differs from the AES only in the following aspects:

- It operates on a 4×4 -nibble state of 64 bits. States and keys are elements of $\mathbb{F}_{2^4}^{4 \times 4}$.
- The S-box operates on nibbles; it is given in Table 4 for completeness.
- The MixColumns multiplications operate in \mathbb{F}_{2^4} modulo $p(x) = x^4 + x + 1$. The values in the MixColumns matrix are equal to those of \mathbf{M} in the original AES.
- The round constants to derive the round key K^i are \mathbf{x}^{i-1} in $\mathbb{F}_{2^4}/p(x)$.

2.3 S-box Properties

We will study variants of Small-AES with different S-boxes. For this purpose, this section revisits some relevant S-box properties from [GR18, GR19].

Mean. For $\Delta_I, \Delta_O \in \mathbb{F}_2^n$, we denote the number of solutions for a differential $\Delta_I \rightarrow \Delta_O$ through the S-box $S : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ as $\delta_S(\Delta_I, \Delta_O)$, defined as

$$\delta_S(\Delta_I, \Delta_O) \stackrel{\text{def}}{=} |\{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \Delta_I) = \Delta_O\}|.$$

Independently of the S-box details, the expected number of solutions x for arbitrary non-zero input-output differences (Δ_I, Δ_O) is simply the mean over all non-zero columns or rows in the differential-distribution table: $\mathbb{E}[\delta_S(\Delta_I, \Delta_O)] = \frac{2^b}{2^b-1} = \frac{256}{255}$ for the AES S-box.

Differential Uniformity. Let $\Delta_I, \Delta_O \in \mathbb{F}_2^b$. A mapping $S : \mathbb{F}_2^b \rightarrow \mathbb{F}_2^b$ is called differential δ -uniform iff for all non-zero $\Delta_I \in \mathbb{F}_2^b$ and all $\Delta_O \in \mathbb{F}_2^b$, it holds that

$$|X \in \mathbb{F}_2^b | S(X) \oplus S(X \oplus \Delta_I) = \Delta_O| \leq \delta.$$

In this work, we use the tight interpretation of

$$\text{DU} = \max_{\substack{\Delta_I \neq 0 \\ \Delta_O \neq 0}} \delta_S(\Delta_I, \Delta_O).$$

An S-box that has $\text{DU} = 2$ is called almost perfect(ly) non-linear (APN).

Variance. In contrast to the mean, the variance $\text{Var}[\delta_S(\Delta_I, \Delta_O)]$ depends on the S-box internals. For each non-zero input difference Δ_I , the S-box of the AES has one output difference Δ_O , s. t. $\delta_S(\Delta_I, \Delta_O)$ has four solutions, 126 differences Δ_I , for which there exist two solutions, and 128 differences without solutions. For the S-box of the AES, the variance is therefore given by

$$\text{Var}[\delta_{S_{\text{AES}}}(\Delta_I, \Delta_O)] = 2^2 \cdot \frac{126}{255} + 4^2 \cdot \frac{1}{255} - \left(\frac{256}{255}\right)^2 = \frac{67\,064}{65\,025}.$$

Differential Spectrum. The differential spectrum of an S-box is a histogram of the number of values $\delta_S(\Delta_I, \Delta_O)$.

Homogeneity. Let $\Delta_I, \Delta_O \in \mathbb{F}_2^b$ be non-zero. An S-box is called homogeneous if the distribution of $\delta_S(\Delta_I, \Delta_O)$ w.r.t. $\Delta_O \neq 0$ is independent of the value Δ_I . It is called inverse homogeneous if the distribution of $\delta_S(\Delta_I, \Delta_O)$ w.r.t. $\Delta_I \neq 0$ is independent of the value Δ_O . For example, the S-box of the AES is homogenous since each non-zero Δ_I has one output difference Δ_O with four solutions, 126 output differences Δ_O with two solutions and 128 non-zero output differences Δ_O with no solutions.

2.4 Known Integral Distinguishers on Round-reduced AES

Integral attacks consider the propagation of sets of inputs through a cryptographic transform such that the sum of their corresponding output parts sums to zero, similarly as cubes [DS09]. The principle had been extended and generalized in the division property by Todo [Tod15]. Integral distinguishers consider parts of the state in a given set, and label the parts at certain states as either

- CONSTANT (**C**) if all texts in a set have a constant value in the considered parts,
- BALANCED (**B**) if the sum of all texts in the considered parts is zero,
- ALL (**A**) (or PERMUTATION) if the texts in a set cover all values of the part exactly once, or
- UNKNOWN (**U**) if no further knowledge of the part is known.

We briefly recall the well-known three-round integral distinguisher for the AES [DKR97].

Three-round Integral Distinguisher. Let $\mathcal{I}, \mathcal{I}', \mathcal{J}, \mathcal{J}' \subseteq \{0, 1, 2, 3\}$. Let X^i denote the i -th element in \mathcal{S} and define $S^{r,i} = R^r(X^i)$ denote the encryption of X^i through r consecutive rounds of AES. Let $r, c \in \{0, 1, 2, 3\}$. We denote by $\mathcal{S} = (\mathcal{D}_{\{r\}} \cap \mathcal{C}_{\{c\}}) \oplus A$ a δ -set, that is a set of 2^8 texts that iterate over all values $x_{c-r \bmod 4, c}$ (the byte at Row $(c-r) \bmod 4$ and Column c is indicated as **A**) and are constant at all other bytes (**C**). $A \in \mathbb{F}_{2^8}^{4 \times 4}$ is any given text. Then, the texts in \mathcal{S} iterate over all 2^8 values in each byte of S^2 after two rounds of AES.

$$\begin{bmatrix} \mathbf{A} & \mathbf{C} & \mathbf{C} & \mathbf{C} \\ \mathbf{C} & \mathbf{C} & \mathbf{C} & \mathbf{C} \\ \mathbf{C} & \mathbf{C} & \mathbf{C} & \mathbf{C} \\ \mathbf{C} & \mathbf{C} & \mathbf{C} & \mathbf{C} \end{bmatrix} \xrightarrow{R^2} \begin{bmatrix} \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \end{bmatrix} \xrightarrow{\text{SR} \circ \text{SB}} \begin{bmatrix} \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \end{bmatrix} \xrightarrow{\text{MC}} \begin{bmatrix} \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \end{bmatrix}.$$

The ALL property is preserved through SubBytes and ShiftRows of the third round, but is no longer guaranteed after the MixColumns operation at the end of Round 3. Since MixColumns is linear, it preserves balanced input sets, i.e., the sum of all 2^8 states $\bigoplus_{i=1}^{2^8} S^{3,i} = 0$. This is indicated by **B**. The subsequent SubBytes operation in Round 4 destroys this Balanced property.

Four-round Integral Distinguisher. It is well-known that the three-round distinguisher can be extended to a four-rounds higher-order distinguisher [DKR97]. Let $\mathcal{D}_{\mathcal{I}}$ for $|\mathcal{I}| = 1$ be a diagonal space. Then, the texts in $\mathcal{D}_{\mathcal{I}}$ iterate over all 2^{32} values in each column of S^3 after almost three rounds of AES.

$$\begin{bmatrix} \mathbf{A} & \mathbf{C} & \mathbf{C} & \mathbf{C} \\ \mathbf{C} & \mathbf{A} & \mathbf{C} & \mathbf{C} \\ \mathbf{C} & \mathbf{C} & \mathbf{A} & \mathbf{C} \\ \mathbf{C} & \mathbf{C} & \mathbf{C} & \mathbf{A} \end{bmatrix} \xrightarrow{R^3} \begin{bmatrix} \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \end{bmatrix} \xrightarrow{\text{SR} \circ \text{SB}} \begin{bmatrix} \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \\ \mathbf{A} & \mathbf{A} & \mathbf{A} & \mathbf{A} \end{bmatrix} \xrightarrow{\text{MC}} \begin{bmatrix} \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \\ \mathbf{B} & \mathbf{B} & \mathbf{B} & \mathbf{B} \end{bmatrix}.$$

The ALL property is no longer guaranteed by the MixColumns operation at the end of Round 3. Since MixColumns is linear, it preserves the balancedness of the outputs, i.e., the sum of all 2^{32} states $\bigoplus_{i=1}^{2^{32}} S^{4,i} = 0$. The SubBytes operation in the subsequent round destroys the Balanced property.

3 Statistical Framework

We follow the framework by [GR19]. The binomial distribution $\mathcal{B}(N, p)$ yields the number of successes in a sequence of N independent Boolean experiments, each of which is successful with probability p . The values μ and variance σ^2 are given by $\mu = N \cdot p$ and $\sigma^2 = N \cdot p \cdot (1 - p)$. The theorem by De Moivre and Laplace allows to approximate the binomial distribution by a normal distribution if the skew is close to zero. Moreover, the difference as well as the sum of two normal distributions also yields a normal distribution. Let $X_0 \sim \mathcal{N}(\mu_0, \sigma_0^2)$ follow a normal distribution with mean μ_0 and variance σ_0^2 . Moreover, let $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$. Then, $X_0 - X_1 \sim \mathcal{N}(\mu_0 - \mu_1, \sigma_0^2 + \sigma_1^2)$.

We consider two normal distributions, where we approximate that their difference is normally distributed with $\mathcal{N}(\mu, \sigma^2)$, with

$$\mu = N \cdot |p_0 - p_1| \quad \text{and} \quad (1)$$

$$\sigma^2 = \sigma_0^2 + \sigma_1^2 = N \cdot p_0 \cdot (1 - p_0) + N \cdot p_1 \cdot (1 - p_1). \quad (2)$$

We will distinguish between at least two settings. In the setting when the distribution for the AES follows a single binomial distribution, we can replace the rightmost term of Equation (2) by $\sigma_1^2 = N \cdot p_1 \cdot (1 - p_1)$.

In the other setting, the distribution $\mathcal{N}(\mu_1, \sigma_1^2)$ will be composed of several binomial distributions. While this affects its mean slightly, it has a strong effect on the variance. In this case, we will determine the quotient σ_1^2/σ_0^2 , and will obtain for the distribution of their differences that

$$\sigma^2 = \sigma_0^2 + \sigma_1^2 = N \cdot \left(p_0 \cdot (1 - p_0) + \frac{\sigma_1^2}{\sigma_0^2} \cdot p_1 \cdot (1 - p_1) \right). \quad (3)$$

Note that Equation (3) generalizes Equation (2), and we can continue to use the latter. Since the probability density of the normal distribution is

$$F(x|\mu, \sigma^2) \stackrel{\text{def}}{=} e^{-\frac{(x-\mu)^2}{2\sigma^2}} \cdot \frac{1}{\sigma\sqrt{2\pi}},$$

it follows that

$$P_S = \int_0^{+\infty} \frac{e^{-\frac{(x-\mu)^2}{2\sigma^2}}}{\sigma\sqrt{2\pi}} dx = \int_{-\frac{\mu}{\sigma}}^{+\infty} \frac{e^{-\frac{x^2}{2}}}{\sqrt{2\pi}} dx = \frac{1}{2} \left(1 + \text{erf} \left(\frac{-\mu}{\sigma\sqrt{2}} \right) \right),$$

where $\text{erf}(x)$ is the error function, i.e., the probability that a normally distributed random variable $X \sim \mathcal{N}(0, 0.5)$ falls into the interval $[-x, x]$. To obtain a success probability of at least P_S , the number of trials N has to satisfy

$$N \geq \frac{2 \left(p_0(1 - p_0) + \frac{\sigma_1^2}{\sigma_0^2} p_1(1 - p_1) \right)}{(p_0 - p_1)^2} \cdot (\text{erfinv}(2 \cdot P_S - 1))^2, \quad (4)$$

where $\text{erfinv}(x)$ is the inverse error function.

4 Four-round Truncated-differential Distinguisher

We extend the deterministic three-round integral distinguisher to a probabilistic four-round truncated-differential distinguisher in the following.

We consider a δ -set of 2^8 plaintexts that iterate over all values in a single plaintext byte and leave all other plaintext bytes constant. For a δ -set, all output bytes after the operation sequence $\widehat{\mathbf{R}}^3 \stackrel{\text{def}}{=} \text{SR} \circ \text{SB} \circ \mathbf{R}[K^2] \circ \mathbf{R}[K^1] \circ \text{AK}[K^0]$ iterate over all values. In general, we will denote by $\widehat{\mathbf{R}}^r$ r rounds, without the key addition and `MixColumns` operation in the final round. So, for each column, the `MixColumns` operation in Round 3 can be viewed as the sum of the results of four permutations where the inputs iterate over all values.

4.1 Setting

Adapting Patarin’s Setting to Three-round AES. We denote by M^i the i -th plaintext message in a δ -set, and let x^i denote the iterated byte. Let v_0^i, v_1^i, v_2^i , and v_3^i be the four bytes of the first column after $\widehat{\mathbf{R}}^3$ that correspond to x^i . For $j \in \{0, 1, 2, 3\}$, we denote

by $\pi'_j : \mathbb{F}_2^8 \rightarrow \mathbb{F}_2^8$ the permutations that map x^i to v_{j-1}^i . In `MixColumns`, the bytes are multiplied by 1, 2, or 3 in \mathbb{F}_{2^8} , whose results are summed together finally. We denote by $P_j : \mathbb{F}_2^8 \rightarrow \mathbb{F}_2^8$ the permutations, consisting of the multiplications from the topmost row of the `MixColumns` matrix and the byte rotation for the XOR operation when our target is Byte $y^i = S_i^{3,MC}[0]$. We write ${}^3_i{}^{MC}$ as the state in Round 3 directly after the application of `MixColumns` for an input x^i . So, $P_0(x) = 2x$, $P_1(x) = 3x$, $P_2(x) = x$, $P_3(x) = x$, in \mathbb{F}_{2^8} .

To be compatible with Patarin's setting, the permutations π'_j would have to be random and independent. Then, if the permutations P_j are independent from the maps π'_j , the resulting compositions $P_j \circ \pi'_j$ are also random permutations. Though, neither the permutations π'_j nor P_j can be fully random; it is an *approximation* to consider their composition as random to apply Patarin's setting. Our setting focuses on $q = 2^n$ queries per permutation evaluation. Of course, the setting can be generalized also to the other columns.

Assumptions. We can adopt the assumptions from [GR19]:

- (1) The number of solutions for non-trivial differentials $\Delta_I \xrightarrow{S} \Delta_O$ through the S-box S

$$S(X) \oplus S(X \oplus \Delta_I) = \Delta_O, \quad (5)$$

are uniformly distributed, for $X, \Delta_I, \Delta_O \in \mathbb{F}_2^b$ for $b \in \{4, 8\}$ ($b = 8$ for AES and $b = 4$ for Small-AES, respectively) and all non-zero differentials Δ_I, Δ_O .

- (2) There is no non-trivial relation between the solutions of several equations (16) that we will study. This means that the probability of one equation must not be a solution of a different one with probability different from that in Equation (17).

Later, we will provide proofs for the mean and variances of our truncated differentials under those assumptions. Moreover, we will also study the behavior of Small-AES experimentally, and will outline deviations from this idealized model in practice.

Byte-collision Probabilities. As a result of our assumption above, we approximate by $\mathbb{E}[N_4]$ the expected number of byte collisions $S_{r,c}^{3,i} = S_{r,c}^{3,j}$. So, according to Table 2 the probability for a byte collision is approximately

$$\Pr_{\text{AES}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}] \simeq \frac{1}{2^8} + \frac{1}{2^8 \cdot (2^8 - 1)^3} \simeq 2^{-8} + 2^{-31.983}. \quad (6)$$

For a random permutation, the probability is roughly

$$\Pr_{\text{rand}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}] = \frac{2^{120} - 1}{2^{128} - 1} \simeq 2^{-8}. \quad (7)$$

Thus, the difference between those two probabilities can be exploited to build a truncated-differential distinguisher on four-round AES, because $S_{r,c}^{3,i} = S_{r,c}^{3,j}$ directly implies a collision between the corresponding bytes in $S^{4,SB}$, which can be computed from the ciphertext by inverting the final `ShiftRows` operation (and the final `MixColumns` operation if present). The distinguisher is depicted on the left side of Figure 1.

Settings. Throughout this work, we will consider two settings:

- (1) We form pairs from δ -sets that have no non-trivial relation to each other.
- (2) We form pairs from δ -sets such that all $\binom{4}{1} \cdot 2^{24}$ δ -sets of a single-diagonal space are contained exactly once.

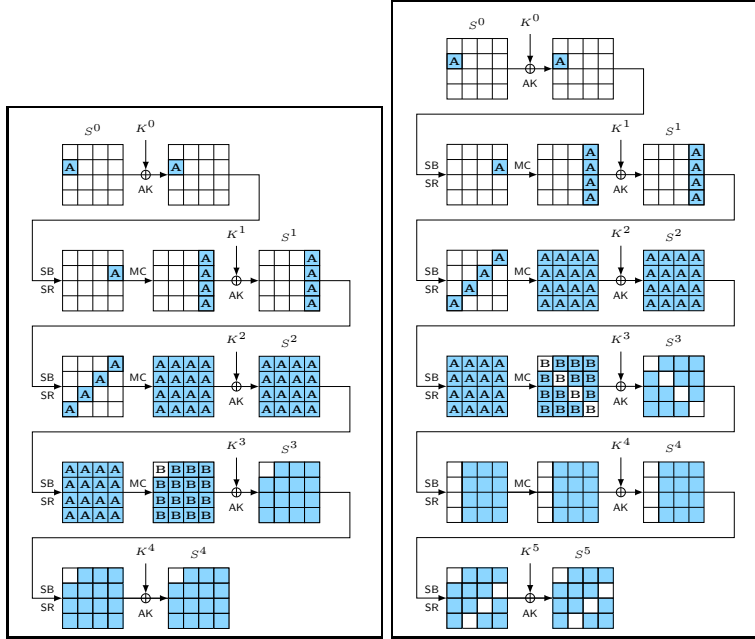


Figure 1: Our distinguishers over four (**left**) and five (**right**) rounds of AES. Darkened cells represent bytes with active (non-zero) difference, white cells represent bytes with zero difference.

The former provides more flexibility to the adversary. We consider the latter only since it serves as a subset of the diagonal space. This will allow us to show our theoretical result. The means of the differentials will not deviate considerably between both settings. However, as Appendix A shows, the latter inherits the multiple-of-eight property from the five-round distinguisher by Grassi et al. [GRR17]. Hence, the probability distributions differ with strong effects, e.g., on the variance.

4.2 Theoretical Analysis

Let Y be a random variable for the number of single-byte collisions after four rounds. Both settings consider the distribution of k collisions from N pairs. In Setting (1), this distribution is described approximately by a binomial distribution, $Y \sim \mathcal{B}(N, p)$, where p is the probability of Equation (6) or Equation (7), respectively.

In Setting (2), the distribution is more complex. Here, it holds that $N = x \cdot 4 \cdot 2^{24} \cdot \binom{2^8}{2}$ stems from x diagonal spaces $\mathcal{D}_{\{c\}}$. Since the number of byte collisions in this setting possesses a multiple-of-eight property, it is reasonable to think that the distribution of eight-tuples would approximately follow a binomial distribution, similar as [GR19]. Let X_3 be a random variable for the number of eight-tuples of collisions. Given N_3 ways of eight-tuples, according to [GR19],

$$\Pr[X_3 = \ell_3] = \binom{N_3}{\ell_3} (p_3)^{\ell_3} (1 - p_3)^{N_3 - \ell_3}$$

where $p_3 = p_{\text{AES}}$ from Equation (6). For the total number of collisions Y , we can derive similarly as Lemma 1 in [GR19] that

$$\Pr[Y = \ell] = \begin{cases} 0 & \text{if } \ell \bmod 8 \neq 0, \\ \binom{N_3}{\ell_3} (p_3)^{\ell_3} (1 - p_3)^{N_3 - \ell_3} & \text{otherwise,} \end{cases}$$

where $\ell = 2^3 \cdot \ell_3$, $N_3 = 2^{31} \cdot 255^4 / 8$ and $p_3 = 2^{-8} + 2^{-31.983}$. In other words, $\Pr[Y] = 2^3 \cdot X_3$, where $X_3 \sim \mathcal{B}(N_3, p_3)$ is a binomial distribution. The factor ℓ_3 represents the number of pairs in the input space that follow the distribution.

We conduct a proof that shows that the theoretical mean of collisions in the first byte matches our claims above, following the footsteps of the expectation proof in [GR19].

The first two rounds of the well-known integral distinguisher on three-round AES can be represented easily as a subspace trail. Let the δ -set iterate over all values of plaintext byte $S_{r,c}^0$, and let $k \stackrel{\text{def}}{=} c - r \bmod 4$ be the zero diagonal after three rounds. Then, there exists the deterministic subspace trail

$$\mathcal{C}_{\{c\}} \cap \mathcal{D}_{\{k\}} \oplus A \xrightarrow{\mathbb{R}} \mathcal{C}_{\{k\}} \oplus A' \xrightarrow{\mathbb{R}} \mathcal{M}_{\{k\}} \oplus A''.$$

The interesting part consists of the probabilistic transition to

$$\mathcal{M}_{\{k\}} \oplus A'' \xrightarrow[\mathbb{P}]{\mathbb{R}} \mathcal{D}_{\mathcal{J}} \cap \mathcal{C}_{\{c\}} \oplus A'''$$

through Round 3 where $\mathcal{J} \subset \{0, 1, 2, 3\}$, $|\mathcal{J}| = 3$, and $c \in \{0, 1, 2, 3\}$. We denote by $\ell \in \{0, 1, 2, 3\}$ the index of the zero-difference diagonal $\ell \notin \mathcal{J}$. A, A', A'', A''' are arbitrary constants in $\mathbb{F}_{2^8}^{4 \times 4}$.

Theorem 1. Let $\mathcal{X} = \{X^i \in \mathbb{F}_{2^8}^{4 \times 4}\}$, for $0 \leq i < 2^{32}$, denote a set of all texts in a coset $\mathcal{M}_{\{k\}} \oplus A$ for some $A \in \mathbb{F}_{2^8}^{4 \times 4}$, $A \in \mathcal{X}^\perp$, where all columns are active. Let $\mathcal{Y} = \{Y^i \in \mathbb{F}_{2^8}^{4 \times 4} : Y^i = \mathbb{R}(X^i)\}$, $0 \leq i < 2^{32}$, denote the set of corresponding outputs after one round of AES. Let Y be a random variable for the number of different ciphertext pairs (Y^i, Y^j) s. t. $Y^i, Y^j \in \mathcal{D}_{\mathcal{J}} \cap \mathcal{C}_{\{c\}} \oplus A'$, for some $A' \in \mathbb{F}_{2^8}^{4 \times 4}$, $\mathcal{J} \subset \{0, 1, 2, 3\}$, $|\mathcal{J}| = 3$, and $c \in \{0, 1, 2, 3\}$. Then, under the assumption of the uniform distribution of non-trivial solutions of differentials through the S-box in Equation (5),

$$\Pr \left[\mathcal{M}_{\{k\}} \oplus A \xrightarrow{\mathbb{R}} \mathcal{D}_{\mathcal{J}} \cap \mathcal{C}_{\{c\}} \oplus A' \right] \simeq 2^{-8} + 2^{-31.983}.$$

Our proof is described in Appendix D. The analysis is conducted in Setting (2), which consists of two steps: First, determine the average number of solutions which yield a collision in a single byte after three rounds for all pairs from δ -sets from a full single-diagonal plaintext space. Second, divide the number of pairs by the number of all pairs in a single-diagonal space. A proof for the mean in Setting (1) of independent δ -sets can use a similar argumentation. A study of the variance can be found in Appendix F.

Complexity. The complexity of our four-round distinguisher is approximated by Equation (4) (plugging $\Pr_{\text{AES}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}]$ into p_{AES} and $\Pr_{\text{rand}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}]$ into p_{rand}). Consider Setting (1) first. For a success probability of $P_S = 0.95$, we obtain $N > 2^{58.402}$ pairs. In total, a δ -set contains $\binom{2^8}{2} \simeq 2^{15}$ pairs. So, the distinguisher needs approximately $2^{43.402}$ δ -sets with $2^{51.402}$ chosen plaintexts. For $P_S = 0.99$, we obtain $N > 2^{59.402}$ pairs, and therefore a data complexity of approximately $2^{52.402}$ chosen plaintexts. Setting (2) has a stronger multiple-of- n property. In theory, it also possesses a considerably higher variance of approximately a factor of eight. Therefore, the data complexity increases by a factor of about $2\sigma_{\text{AES}}^2/\sigma_{\text{rand}}^2 = 16$. For $P_S = 0.95$, we obtain $2^{62.402}$ pairs and $2^{47.402}$ δ -sets of $2^{55.402}$ chosen plaintexts. For $P_S = 0.99$, we obtain $N > 2^{63.402}$ pairs, and therefore a data complexity of approximately $2^{56.402}$ chosen plaintexts.

Reduced Variant. In order to allow a practical verification, we derive the corresponding probabilities for four rounds of Small-AES, the downscaled variant from [CMR05] which employs four-bit S-boxes. Again, we view the downscaled variant of MixColumns as a sum

Table 3: Means and standard deviations for the number of pairs that collided in a fixed output byte for our four-round distinguisher with Small-AES in Setting (1). Each experiment employed 100 random independent keys and 2^s random δ -sets. Experimental values are rounded. Full-round Speck-64-96 was used as pseudorandom permutation π .

s	Theory				Experiments			
	Small-AES		π		Small-AES		π	
	μ	σ	μ	σ	μ	σ	μ	σ
20	7 866 650	2 715.665	7 863 200	2 715.290	7 870 789	2 918.	7 864 396	2 566.
21	15 733 300	3 840.531	15 728 600	3 840.000	15 742 188	3 809.	15 728 650	3 957.
22	31 466 600	5 431.331	31 457 300	5 430.580	31 484 544	6 007.	31 457 205	5 096.
23	62 933 200	7 681.062	62 914 600	7 680.000	62 967 244	7 030.	62 915 004	7 820.
32	32 221 796 921.	173 802.580	32 212 254 720.	173 778.563				

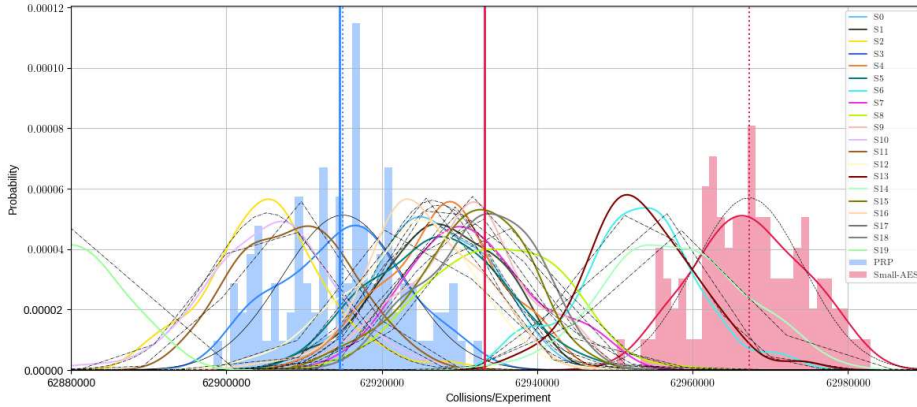


Figure 2: Distribution of means from 100 experiments with random keys for our four-round distinguisher on variants of Small-AES with random S-boxes. Each experiment used 2^{23} independent δ -sets. The bold bars indicate the theoretical predictions for the means of a random permutation and for Small-AES; the dotted bars indicate the experimental means of π and Small-AES.

of four independent permutations on \mathbb{F}_{2^4} . Then, the probability that, after three rounds, a fixed nibble has zero difference is approximately

$$\Pr_{\text{Small-AES}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}] \simeq \frac{1}{2^4} + \frac{1}{2^4 \cdot (2^4 - 1)^3} \simeq 2^{-4} + 2^{-15.721}.$$

For a random permutation, the probability can be approximated by

$$\Pr_{\text{rand}} [S_{r,c}^{3,i} = S_{r,c}^{3,j}] = \frac{2^{60} - 1}{2^{64} - 1} \simeq 2^{-4} - 2^{-64.093}.$$

Applying Equation (4) yields $N > 2^{29.878}$ for $P_S = 0.95$ in Setting 1. Since a δ -set yields $\binom{2^4}{2} = 120$ pairs, we need about 2^{23} δ -sets with 2^{27} chosen plaintexts. For $P_S = 0.99$, we obtain $N > 2^{30.878}$ pairs, and therefore a data complexity of approximately 2^{28} chosen plaintexts.

4.3 Verification

Experimental Verification. We verified the distinguisher experimentally with 100 random keys and 2^s δ -sets, for $s \in \{20, \dots, 23\}$. Note that an experiment means to compute

the number of collisions for 2^s δ -sets and a random independent key. For each experiment, we evaluated the number of collisions of the first nibble of the output. As approximation of a random permutation, we employed full-round Speck-64-96 with 100 random keys. Our results are listed in Table 3. The values μ denote the obtained means of the number of pairs that collide in at least on inactive inverse diagonal, over all experiments, e.g. for 2^{20} δ -sets, one could expect $2^{20} \cdot \binom{2^4}{2} \cdot (2^{-4} + 2^{-15.721}) \approx 7866650$ colliding pairs per experiment.

Transition-based Verification. Rønjom [Røn19] proposed a general approach of transition-distribution matrices for substitution-permutation networks. Starting from a vector that represents the distribution of active-bytes-per-column in plaintexts, Rønjom’s approach allows to derive the probability of active bytes per columns after r rounds of an SPN. Thus, his results can be viewed as a different view on the same distinguishing events that verify our attacks under idealized assumptions of the AES S-box. We employed his implementation for a second form of verification. We obtained

$$\begin{aligned} p_{\text{AES}} &\simeq 0.00390625023558\dots \simeq 2^{-8} + 2^{-31.98306} \text{ and} \\ p_{\text{Small-AES}} &\simeq 0.06251851851852\dots \simeq 2^{-4} + 2^{-15.72067}, \end{aligned}$$

which basically confirmed our biases.

Strange Effects of the Variance. From our theoretical considerations, we initially expected that the number of collisions for each single diagonal yielded always a multiple of eight. Since the variance is defined as $\sigma^2 = N \cdot p \cdot (1 - p)$, one would expect that the variance in Setting (2) is about eight times higher than that in Setting (1) due to the multiple-of-eight property. So, we conducted experiments on the variance in both settings, whose results are provided in Appendix F. Interestingly, our experiments with Small-AES yielded that the numbers of collisions were always multiples of 32. From experiments with alternative S-boxes, it seemed that this is not a property of the S-box. More thoughts and a multiple-of-32 distinguisher are summarized in Appendix B.

Our experiments show that the variance differed significantly between the S-boxes. The variance for the original S-box of Small-AES was about 15.3 times higher in Setting (2) than the variance in Setting (1). For Small-AES with the S-box of PRESENT, the factor was roughly 12.7. Using the S-box of PRIDE, we obtained a huge deviation with a factor of roughly 40. Note that this property affects the mean of our distinguisher only marginally. The reason for the strong deviations in the variances, however, remains an interesting future work.

Alternative S-boxes. The results for Small-AES in Table 3 deviate significantly from the theoretical prediction (62 967 244 vs. 62 933 200) in the sense that Small-AES appears to yield a considerably better distinguisher than anticipated. The naturally arising question is which reasons cause this phenomenon. In the beginning, we aimed to search additional internal properties due to the AES. Inspired by the experiments from [GR18, GR19], we furthermore conducted studies with variants of Small-AES that employed different four-bit S-boxes. First, we considered three real-world ciphers PRESENT [BKL+07], PRINCE [BCG+12], and PRIDE [ADK+14], as well as three artificial constructions called Toy-6, Toy-8, and Toy-10 from [GR19]. For the sake of self-containment, the details of those S-boxes are listed in Table 4a, and their differential spectra in Table 4b. The S-box of Small-AES had been constructed in a similar manner as that of the AES, whereas the real-world examples had been selected for security and lightweight properties. The S-boxes of PRINCE and PRIDE contain one-dimensional linear structures. In this context, the x in Toy- x denotes the differential uniformity of an artificial S-box, i.e., the maximal number of pairs occurring for some differential.

Table 4: Properties of the employed S-boxes.

(a) The employed S-boxes.													(b) Differential spectra of the individual S-boxes.										
x	S(x)												Instance / m	p_m									
	0	1	2	3	4	5	6	7	8	9	a	b		c	d	e	f	0	2	4	6	8	10
Small-AES	6	b	5	4	2	e	7	a	9	d	f	c	3	1	0	8	Small-AES	120/225	90/225	15/225	–	–	–
PRESENT	c	5	6	b	9	0	a	d	3	e	f	8	4	7	1	2	PRESENT	129/225	72/225	24/225	–	–	–
PRINCE	b	f	3	2	1	5	e	9	2	7	a	c	e	5	d	4	PRINCE	120/225	90/225	15/225	–	–	–
PRIDE	0	4	8	f	4	5	6	7	8	9	a	b	b	d	6	3	PRIDE	129/225	72/225	24/225	–	–	–
Toy-6	1	3	6	4	2	5	9	a	0	f	7	e	c	b	d	8	Toy-6	125/225	81/225	18/225	1/225	–	–
Toy-8	1	3	6	4	2	5	a	c	0	f	7	8	e	b	d	9	Toy-8	130/225	74/225	18/225	2/225	1/225	–
Toy-10	6	4	c	5	0	7	2	e	1	f	3	d	8	a	9	b	Toy-10	140/225	60/225	17/225	7/225	–	1/225

For each S-box of them, we employed again 2^{23} independent δ -sets. Surprisingly, our results indicate that the employed S-box can produce tremendous biases, in the positive, but also in the negative side. Our results are listed in Table 11 and illustrated in Figure 2. The precise results are listed in Appendix G. In summary:

- The three real-world S-boxes show lower, but noticeable distance to the random permutation. While PRESENT and PRINCE are below the prediction, the distance to the pseudo-random permutation are more than 1.8 standard deviations.
- Two out of three toy S-boxes, Toy-8 and Toy-10, behave close to random and do not yield a distinguisher for the employed amount of data.
- The Toy-6 S-box strongly deviates from all others by having a significantly lower mean than the random permutation and all others.

To increase our understanding, we also tested randomly generated S-boxes. We sampled 20 S-boxes S_i for $i \in \{0, \dots, 19\}$ in a reproducible manner, and ran our four-round distinguisher in the same setting. The details of the sampling process and the results are also provided in Appendix G. We observed considerable deviations and two drastic outliers (62 626 909.73, 62 876 977.05, and 63 054 909.76); the means of seven experiments were below $\mu - \sigma$; the means of nine clustered around the theoretical prediction of $[\mu - \sigma, \mu + \sigma]$; the means of four experiments exceeded the interval.

It seems that S-boxes that possess a low DU and a low variance produce distinguishers with sufficient distance. However, it is far from trivial to deduce a precise reason why the S-box of Small-AES deviates much from random and from the theoretical prediction. While our experiments can only try to capture an intuition, they provide at least the insight that the considerable deviation from our prediction is less due to our approximation for the probability of truncated differentials, but rather caused by the S-box.

4.4 Index Dependencies

The mean of the four-round truncated differential depends strongly on the index of (1) the input cell that is iterated over in the δ -sets as well (2) the index of the inactive output cell after three rounds. Let i_{in} and i_{out} denote the cell indices of the relevant input and output cells. We obtain an equation system through three full rounds. Let $x = S(P[i_{\text{in}}] \oplus K^0[i_{\text{in}}])$. Since we consider all values of x , $K^0[i_{\text{in}}]$ is a constant, and the S-box a permutation, we can simply iterate over all values x and neglect K^0 . For simplicity, let us define $x_i \stackrel{\text{def}}{=} i$. Note that the position of x_i is relevant, and the ShiftRows operations will affect the coefficients of the equation system. We define by $(a_i, b_i, c_i, d_i) \in \mathbb{F}_{2^8}^4$ variables in terms of x_i at S_{SB}^2 ,

i.e., after the `SubBytes` operation in the second round. Note that $b = 4$ for the Small-AES and $b = 8$ for the AES. For $i_{\text{in}} = 0$, we define

$$\begin{pmatrix} a_i \\ b_i \\ c_i \\ d_i \end{pmatrix} \stackrel{\text{def}}{=} \begin{pmatrix} S(2x_i \oplus K^1[0]) \\ S(x_i \oplus K^1[1]) \\ S(x_i \oplus K^1[2]) \\ S(3x_i \oplus K^1[3]) \end{pmatrix}.$$

Note that K^1 also covers the constant values of the plaintext bytes 1, 2, 3 here. Since we will iterate over all values of all constants, we consider all possible values. After the second round, and the subsequent `SubBytes` operation in Round 3, we obtain an equation system that will be rotated by `ShiftRows` to

$$\begin{pmatrix} S(2a_i \oplus K^2[0]) & S(d_i \oplus K^2[4]) & S(c_i \oplus K^2[8]) & S(3b_i \oplus K^2[12]) \\ S(d_i \oplus K^2[5]) & S(3c_i \oplus K^2[9]) & S(2b_i \oplus K^2[13]) & S(a_i \oplus K^2[1]) \\ S(2c_i \oplus K^2[10]) & S(b_i \oplus K^2[14]) & S(a_i \oplus K^2[2]) & S(3d_i \oplus K^2[6]) \\ S(b_i \oplus K^2[15]) & S(3a_i \oplus K^2[3]) & S(2d_i \oplus K^2[7]) & S(c_i \oplus K^2[11]) \end{pmatrix} = (\dots), \quad (8)$$

where the right-hand side will reflect the left-hand side with (a_j, b_j, c_j, d_j) for some $j \neq i$. For each output cell, we will only require a column of the state in Equation (8), i.e., four variables after the `MixColumns` operation at the end of Round 3. Each output cell will contain a dependency of a , b , c , and d exactly once. So, we have to consider only four key variables of K^2 . Since we consider collisions in single cells after four rounds, and can invert `MixColumns` and `ShiftRows` in that row, neither K^4 nor the S-box of Round 4 nor K^3 influence the number of collisions. For example, the equation system for a cell collision at position $i_{\text{out}} = 0$ becomes

$$2S(2S(2x_i \oplus K^1[0]) \oplus K^2[0]) + 3S(S(3x_i \oplus K^1[1]) \oplus K^2[5]) \quad (9)$$

$$+ S(2S(x_i \oplus K^1[2]) \oplus K^2[10]) + S(S(x_i \oplus K^1[3]) \oplus K^2[15]) \quad (10)$$

$$= 2S(2S(2x_j \oplus K^1[0]) \oplus K^2[0]) + 3S(S(3x_j \oplus K^1[1]) \oplus K^2[5]) \quad (11)$$

$$+ S(2S(x_j \oplus K^1[2]) \oplus K^2[10]) + S(S(x_j \oplus K^1[3]) \oplus K^2[15]) \quad (12)$$

for $i \neq j$. For other input or output positions, the equations differ naturally. In all cases, however, we have eight cells as degrees of freedom, i.e., 2^{32} possible values of keys and plaintext. For simplicity, we denote them as $\kappa^1[0..3], \kappa^2[0..3]$ in all cases. For each value of the eight $\kappa^1[0..3], \kappa^2[0..3]$, we compute the equation over all 2^4 values of x_i , form their $\binom{2^4}{2}$ pairs, and count the number of collisions. For an ideal permutation, we would expect

$$\mu_{\text{rand}} = (2^4)^8 \cdot \binom{2^4}{2} \cdot (2^{-4} - 2^{-64.093}) \simeq 32\,212\,254\,719.999 \dots \text{ and}$$

$$\sigma_{\text{rand}} \simeq 173\,778.563,$$

whereas we would expect for Small-AES in theory – that is, without an influence of the indices from the `MixColumns` matrix or the S-box:

$$\mu_{\text{Small-AES}} = (2^4)^8 \cdot \binom{2^4}{2} \cdot (2^{-4} + 2^{-15.792}) \simeq 32\,221\,796\,920.671 \dots \text{ and}$$

$$\sigma_{\text{Small-AES}} \simeq 173\,802.580.$$

Implementation. In practice, the influences of S-box and `MixColumns` matrix cause deviations that are predictable by the 16×16 cell-dependent equations. We implemented

(0)	(4)	(8)	(12)
0 4 8 12 1 5 9 13 2 6 10 14 3 7 11 15	12 0 4 8 13 1 5 9 14 2 6 10 15 3 7 11	8 12 0 4 9 13 1 5 10 14 2 6 11 15 3 7	4 8 12 0 5 9 13 1 6 10 14 2 7 11 15 3
(1)	(5)	(9)	(13)
15 3 7 11 12 0 4 8 13 1 5 9 14 2 6 10	11 15 3 7 8 12 0 4 9 13 1 5 10 14 2 6	7 11 15 3 4 8 12 0 5 9 13 1 6 10 14 2	3 7 11 15 0 4 8 12 1 5 9 13 2 6 10 14
(2)	(6)	(10)	(14)
10 14 2 6 11 15 3 7 8 12 0 4 9 13 1 5	6 10 14 2 7 11 15 3 4 8 12 0 5 9 13 1	2 6 10 14 3 7 11 15 0 4 8 12 1 5 9 13	14 2 6 10 15 3 7 11 12 0 4 8 13 1 5 9
(3)	(7)	(11)	(15)
5 9 13 1 6 10 14 2 7 11 15 3 4 8 12 0	1 5 9 13 2 6 10 14 3 7 11 15 0 4 8 12	13 1 5 9 14 2 6 10 15 3 7 11 12 0 4 8	9 13 1 5 10 14 2 6 11 15 3 7 8 12 0 4

Figure 3: Permutations of numbers of solutions for each input byte, where the bold index in parentheses refers to the position of the δ -sets in the plaintext, i_{in} , and the position in the small squares the index of the output (colliding) cell i_{out} . The blue cells are exemplary for the position $(i_{in}, i_{out}) = (0, 15)$, the best distinguisher for Small-AES.

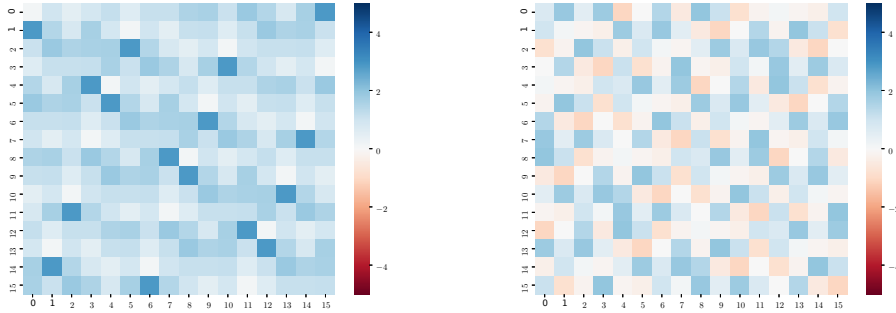


Figure 4: Theoretical index-dependent means from the equation system for our four-round distinguisher with the S-box of Small-AES and that of PRESENT, in $\times |p_{\text{Small-AES}} - p_{\text{rand}}|$ deviations from the mean of the PRP. Each cell represents the result for the active input cell in the plaintexts (y-axis) and the tested inactive cell (x-axis) after four rounds without the final MixColumns operation. Darker color (ignorant of sign) imply better distinguishers.

the system for Small-AES with its original and six alternative S-boxes. The results are provided in Figure 4. For each S-box, we obtained distinct values for each output cell. The numbers of solutions in each row are permutations of each other with strong regularity. More precisely, the numbers of solutions repeat as shown by equal indices in Figure 3. This means, the first row of the left part of Figure 4 corresponds to the top-left of Figure 3. For $i_{in} = 0$, the position $i_{out} = 15$ yields the highest number of collisions, 32 239 747 072. For $i_{in} = 1$, this value repeats at index $i_{out} = 0$; for $i_{in} = 2$, this value repeats at index $i_{out} = 5$; and for $i_{in} = 3$, this value repeats at index $i_{out} = 10$. This corresponds to the darkest cells in left part of Figure 4. The situation is the same for the distinguisher with alternative S-boxes, e.g., compare to that of PRESENT, in the right part of Figure 4.

Practical Verification. We implemented the position-dependent four-round distinguisher also for Small-AES with its original S-box, and with the PRESENT S-box to verify our approach. Figure 5 illustrates the average means from our experiments with 2^{23} δ -sets for Small-AES on the left, in terms of the deviation from the theoretical mean of a

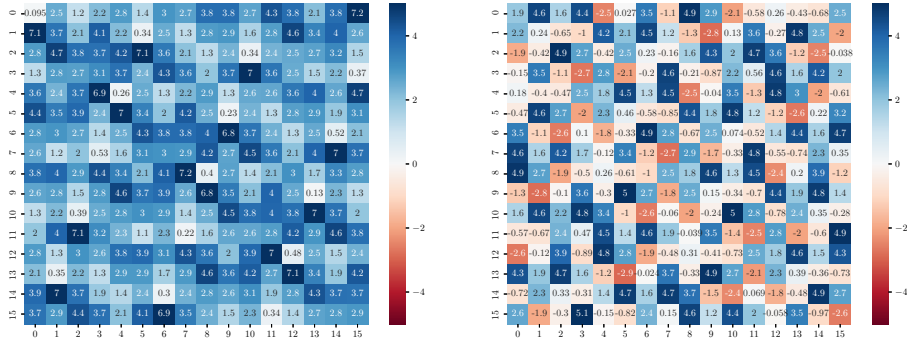


Figure 5: Average means of 50 experiments with random keys and 2^{23} δ -sets on Small-AES with its original S-box (left) and with the PRESENT S-box (right) in $\times\sigma$ deviations from the mean of the PRP. Each cell represents the result for the active input cell in the plaintexts (y-axis) and the tested inactive cell (x-axis) after four rounds without the final MixColumns operation.

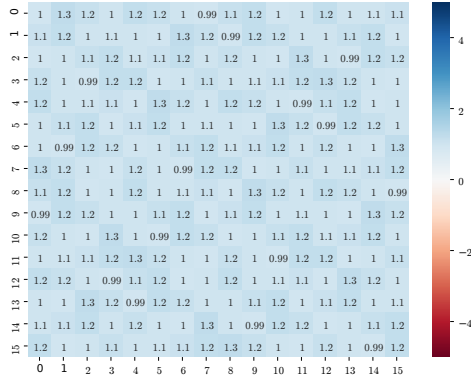


Figure 6: Theoretical index-dependent means from the equation system for our four-round distinguisher for the AES, in $\times|p_{\text{AES}} - p_{\text{rand}}|$ deviations from the mean of the PRP. Each cell represents the result for the active input cell in the plaintexts (y-axis) and the tested inactive cell (x-axis) after four rounds without the final MixColumns operation. 1.0 means that the distinguisher matches our prediction closely.

pseudo-random permutation 62914 600 in multiples of its standard deviation of $\sigma \simeq 7931$. Therefore, we consider the trail from input nibble 1 to output nibble 15 for our distinguisher for Small-AES. The corresponding figure for the PRP possessed about $\pm 0.6\sigma$ as highest deviation. One can observe that the figures are close to our theoretical prediction. One can further observe significant deviations between the combinations. Recall that the range of $[-\sigma.. +\sigma]$ yields no obvious distinguisher. In contrast, the range of $[+2\sigma.. +4\sigma]$ yields roughly the theoretically expected distinguisher. Interestingly, the majority of combinations yields distinguishers for Small-AES. More detailed, for each input index, a single output index possesses a mean close to the pseudorandom value, and two further indices are mostly in the range of $[\sigma.. +2\sigma]$. All further combinations give higher means. In particular, one combination per input index peaked with a derivation of approximately $+7\sigma$.

Application to The AES. The equation for each output cell in (9) consists of eight key-dependent variables. While this is infeasible to be computed naively for the AES, we can

observe that it consists of four terms whose key cells are pair-wise independent. Thus, we could split it into four terms that could be computed individually. For each of the $\binom{2^8}{2}$ combinations of x_i, x_j , and each term $\ell \in \{0, 1, 2, 3\}$, we have $2^8 \cdot 2^8$ solutions that could be stored in a table \mathcal{L}_ℓ . Let $f_0(K^1[0], K^2[0], x_i) \stackrel{\text{def}}{=} 2S(2S(2x_i \oplus K^1[0]) \oplus K^2[0]) = Y$, where $K^1[0]$, $K^2[0]$, and Y are variables in \mathbb{F}_2^8 and x_i is the iterating variable from the δ -set. Then, \mathcal{L}_0 contains the 2^{16} solutions for each possible value of x_i . Thus, \mathcal{L}_0 has 2^8 lists of 2^{16} elements.

Having fixed x_i, x_j , we can derive the differences of $f_\ell(K^1, K^2, x_i) \oplus f_\ell(K^1, K^2, x_j)$ and create a histogram table of all differences, $\Delta\mathcal{L}_\ell$. This means $\Delta\mathcal{L}_\ell[\Delta]$ contains the number of pairs (x_i, x_j) that $f_\ell(K^1[0], K^2[4], x_i) \oplus f_\ell(K^1[0], K^2[4], x_j) = \Delta$. We can repeat this step for all values of ℓ and merge the histogram tables $\Delta\mathcal{L}_0$ and $\Delta\mathcal{L}_1$ by taking the product and create a table $\Delta\mathcal{L}_{0,1}$. The same step can be computed for $\Delta\mathcal{L}_2$ and $\Delta\mathcal{L}_3$ to create a table $\Delta\mathcal{L}_{2,3}$. Finally, we consider the number of collisions by multiplying $\Delta\mathcal{L}_{0,1}[\Delta]$ and $\Delta\mathcal{L}_{2,3}[\Delta]$ element-wise for each $\Delta \in \mathbb{F}_2^8$ and accumulate all results. This yields the number of $(K^1[0, 1, 2, 3], K^2[0, 5, 10, 15], x_i, x_j)$ that collide in a particular output byte. We repeat this step for all input-output combinations. In total, this approach needs

$$16 \cdot 16 \cdot \left(\left(4 \cdot \binom{2^8}{2} \cdot 2^8 \cdot 2^8 \right) + 3 \cdot (2^8 \cdot 2^8) \right)$$

which is feasible with $O(2^{41})$ lookups.

For an ideal permutation, we would expect

$$\begin{aligned} \mu_{\text{rand}} &= (2^8)^8 \cdot \binom{2^8}{2} \cdot 2^{-8} \simeq 2\,351\,959\,869\,397\,967\,831\,040 \quad \text{and} \\ \sigma_{\text{rand}} &\simeq 48\,402\,195\,468.37, \end{aligned}$$

whereas we would expect for AES in theory – that is, without an influence of the indices from the MixColumns matrix or the S-box:

$$\begin{aligned} \mu_{\text{Small-AES}} &= (2^8)^8 \cdot \binom{2^8}{2} \cdot (2^{-8} + 2^{-31.983}) \simeq 2\,351\,960\,011\,247\,373\,747\,803. \quad \text{and} \\ \sigma_{\text{Small-AES}} &\simeq 48\,402\,196\,922.24. \end{aligned}$$

The results are depicted Figure 6 in multiples of $\times |p_{\text{AES}} - p_{\text{rand}}|$ of the deviation from the mean of the PRP. The most important observations are the following that our proposed distinguisher seems to work: all deviations are close to 1.0 – more precisely, in the range $[0.987..1.256]$ of the difference $|p_{\text{AES}} - p_{\text{rand}}|$ – and there exist no significant deviations from it. Moreover, we can observe the regular structure where each row (results of each input cell) is a permutation of the numbers of solutions of the row (of the results for input cell 0).

4.5 S-box Dependencies

The index pattern depends also strongly on the employed S-box. For comparison, the right side of Figure 5 shows the corresponding input-output pattern for Small-AES with the S-box of PRESENT. It differs strongly from that of Small-AES with its original S-box. In general, the obtained means are lower and the peaks lower. All input positions have in common that there exist two to three output cells for each that input-cell position with means of $[4.5\sigma..5.5\sigma]$, i.e., strong distinguishers. Moreover, almost each input position (except two and five) share equally usable distinguishers with means of $[-2\sigma..-3\sigma]$ below the mean of the PRP.

Figure 21 yields the details also for the S-boxes of PRINCE, PRIDE, Toy-6 and Toy-10. The considerably higher number of dark spots in the representations of the latter three illustrate clearly the strong effect of the S-box in both directions, for positive as well as negative biases.

S-box Properties. The question remains which S-box properties cause the deviations. We tested also particularly well-suited S-boxes: the 16 affine-equivalent representatives of optimal S-boxes from [LP07] with respect to differential and linear cryptanalysis. Moreover, we also considered the 10 representatives of platinum S-boxes with minimal CardD1 (number of 1-bit-to-1-bit differentials with non-zero probability) and CardL1 (number of 1-bit-to-1-bit linear approximations) from [ZBRL15].

To compare the quality of the S-boxes, we agreed on a metric that captures the average absolute deviation of each output cell to the theoretical expected number of collisions $p_{\text{Small-AES}}$ for our four-round distinguisher. So, we used the Euclidean distance for this purpose. Since the set of numbers of collisions For an S-box $S : \mathbb{F}_2^4 \rightarrow \mathbb{F}_2^4$, let D_S denote our distance, $X_{i_{\text{out}}}^S$ the number of collisions from plaintext δ -sets that iterate over Cell $i_{\text{in}} = 0$ and count the number of pairs that collide in Cell i_{out} after four rounds without the final MixColumns operation. Moreover, let $\mathbb{E}[X] \simeq 32\,221\,796\,921$ denote the expected number of collisions. Note that the multi-set $\{X_{i_{\text{out}}}^S\}_{0 \leq i_{\text{out}} \leq 15}$ is invariant for different values of i_{in} . Then

$$D_S \stackrel{\text{def}}{=} \sqrt{\sum_{i_{\text{out}}=0}^{15} |X_{i_{\text{out}}}^S - \mathbb{E}[X]|^2}$$

We tested whether the distances are correlated with particular S-box properties. Grassi and Rechberger [GR19] already suggested that their five-round distinguisher on the AES would be related to the variance of the S-box. Therefore, the correlation of the D_S measure to the S-box would be interesting, which could be evaluated by the Pearson correlation coefficient ρ . We determined its r -value and p -value; Recall that for two random variables X and Y :

$$\rho_{X,Y} \stackrel{\text{def}}{=} \frac{\text{cov}(X,Y)}{\sigma_X \cdot \sigma_Y},$$

where $\text{cov}(X,Y) \stackrel{\text{def}}{=} \mathbb{E}[(X - \mu_X) \cdot (Y - \mu_Y)]$ is the covariance of X and Y . We obtained $(r, p) \simeq (0.812, 1.637 \cdot 10^{-13})$, i.e., a high correlation of 0.812; the p -value denotes a very small probability that our results were observed but would be uncorrelated. The relationship is depicted in Figure 7. S-boxes with low variances are considerably closer to the expected numbers of collisions. The highest deviations were observed for the random S-boxes S_{17} and S_3 as well as for Toy-10, whose variances were the highest 656, 608, and 608. Though, several S-boxes with low variance of 416/256 (e.g. that of PRIDE) had considerably large values of D_S . As a result, the variance is related, though, does not fully explain the quality of an S-box. One a side note, we further observed that the ten S-boxes (nine optimal S-boxes 3–7, 10–13, and the Small-AES S-box) whose autocorrelation-table entries $a_{i,j}$, for $i, j > 0$ consisted only values ± 8 and contained no ± 16 , (which implied the minimal variance of 344/256) all had low values of D_S . Note that this implied the absence of linear structures.

5 Five-round Truncated-differential Distinguisher

We can extend the four-round distinguisher from Section 4 to five rounds.

5.1 Setting

Consider some diagonal space $\mathcal{D}_{\{c\}}$ for some index $c \in \{0, 1, 2, 3\}$. Then, the expected probability that all four bytes in that diagonal space collide for two texts in a δ -set can

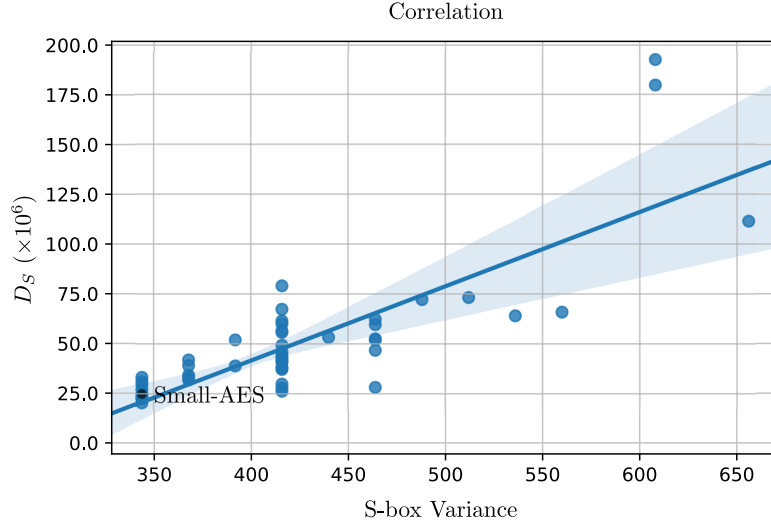


Figure 7: Correlation of D_S , i.e., the squared distance to the expected number of collisions over all output-cell positions and the variance of our tested S-boxes.

be approximated by:

$$\Pr_{\text{AES}} [S^3 \in \mathcal{D}_{\{c\}}] \simeq \left(2^{-8} + \frac{1}{2^8 \cdot (2^8 - 1)^3} \right)^4 \simeq (2^{-8} + 2^{-31.983})^4,$$

Naively, we can approximate the same probability for a random permutation by

$$\Pr_{\text{rand}} [S^3 \in \mathcal{D}_{\{c\}}] = \frac{2^{96} - 1}{2^{128} - 1} \simeq 2^{-32} - 2^{-128}.$$

This approximation is naive since the texts are not independent. In [GR18, Appendix C], Grassi and Rechberger provide arguments for the number of collisions for 2^{32} input texts that can be combined to pairs with each other to a mixed space $\mathcal{M}_{\mathcal{Z}}$ for $|\mathcal{Z}| = 3$. They consider the dependency between multiple texts. Still, they show that the collision probability can be well approximated by 2^{-32} .

Say, we observe the ciphertexts after five rounds that correspond to a δ -set of 2^8 plaintexts. Again, the final MixColumns operation can be simply inverted to obtain the differences before it. Then, the probability to have at least one all-zero anti-diagonal in the difference is given for five-round AES by

$$p_{\text{AES}} \simeq 1 - \left(1 - \Pr_{\text{AES}} [S^3 \in \mathcal{D}_{\{c\}}] \right)^4 \simeq 2^{-30} + 2^{-51.985},$$

whereas for a random permutation, it is approximately

$$p_{\text{rand}} \simeq 1 - \left(1 - \Pr_{\text{rand}} [S^3 \in \mathcal{D}_{\{c\}}] \right)^4 \simeq 2^{-30} - 2^{-61.415}.$$

The distinguisher is depicted on the right side of Figure 1.

5.2 Theoretical Analysis

As for our four-round distinguisher, we show the theoretical probability of inverse-diagonal collisions after almost five rounds in Setting (2). Our proof follows the mean-based proof

Table 5: S-box properties, in ascending order by D_S . DU = differential uniformity, lin. = linearity, NL = nonlinearity, LS = linear structures.

Name	Sbox	D_S	Var.	DU	Lin.	NL	LS
Optimal-06	[0 1 2 13 4 7 15 6 8 12 11 9 10 14 5 3]	20440130	344.00	4	8	4	-
Optimal-04	[0 1 2 13 4 7 15 6 8 12 9 11 10 14 5 3]	22521284	344.00	4	8	4	-
Optimal-03	[0 1 2 13 4 7 15 6 8 12 5 3 10 14 11 9]	22706595	344.00	4	8	4	-
Optimal-11	[0 1 2 13 4 7 15 6 8 14 11 10 5 9 12 3]	23218095	344.00	4	8	4	-
Optimal-07	[0 1 2 13 4 7 15 6 8 12 14 11 10 9 3 5]	24444860	344.00	4	8	4	-
Small-AES	[6 11 5 4 2 14 7 10 9 13 15 12 3 1 0 8]	24893844	344.00	4	8	4	-
Platinum-05	[0 13 4 11 7 14 9 2 6 10 3 5 8 1 15 12]	26272446	416.00	4	8	4	✓
Optimal-12	[0 1 2 13 4 7 15 6 8 14 11 10 9 3 12 5]	27188316	344.00	4	8	4	-
Platinum-02	[0 12 9 7 6 1 15 2 3 11 4 14 13 8 10 5]	27800383	416.00	4	8	4	✓
Random-06	[12 4 0 7 9 6 13 3 5 14 8 10 2 15 1 11]	28205615	344.00	6	12	2	✓
	⋮						
Random-02	[8 4 1 11 0 5 12 13 2 6 15 9 14 7 3 10]	73193659	512.00	6	12	2	✓
PRIDE	[0 4 8 15 15 14 9 2 7 10 12 11 13 6 3]	79057073	416.00	4	8	4	✓
Random-17	[11 14 13 8 2 7 12 6 5 0 4 15 10 9 3 1]	111443494	656.00	8	16	0	✓
Random-03	[7 1 5 10 4 11 12 3 13 8 9 15 14 2 0 6]	179623034	608.00	8	16	0	✓
Toy-10	[6 4 12 5 0 7 2 14 1 15 3 13 8 10 9 11]	192390502	608.00	10	12	2	✓

by Grassi and Rechberger [GR19]. A proof of the mean in Setting (1) δ -sets is similar, and the mean should not change. Moreover, we consider a proof of the variance. Appendix E considers the variance in both settings for the AES and tries to provide arguments for Small-AES.

The core of our truncated-differential distinguisher is formed by the probabilistic transition to

$$\mathcal{M}_{\{k\}} \oplus A \xrightarrow[p]{R} \mathcal{D}_{\mathcal{J}} \oplus B$$

through Round 3 where $\mathcal{J} \subset \{0, 1, 2, 3\}$ and $|\mathcal{J}| = 3$. Let $\ell \in \{0, 1, 2, 3\}$ be the index of the zero-difference diagonal $\ell \notin \mathcal{J}$. Here, $A, B \in \mathbb{F}_{2^8}^{4 \times 4}$ are arbitrary constants. Theorem 2 aims to approximate the probability p for the transition of the third round, under the condition that all columns after two rounds are active.

Theorem 2. Let $\mathcal{X} = \{X^i \in \mathbb{F}_{2^8}^{4 \times 4}\}$, for $0 \leq i < 2^{32}$, denote a set of all texts in a coset $\mathcal{M}_{\{k\}} \oplus A$ for some $A \in \mathbb{F}_{2^8}^{4 \times 4}$, $A \in \mathcal{X}^\perp$, where all columns are active. Let $\mathcal{Y} = \{Y^i \in \mathbb{F}_{2^8}^{4 \times 4} : Y^i = R(X^i)\}$, $0 \leq i < 2^{32}$, denote the set of corresponding outputs after one round of AES. Then, under the assumption of the uniform distribution of non-trivial solutions of S-box differential transitions,

$$\Pr \left[\mathcal{M}_{\{k\}} \oplus A \xrightarrow[p]{R} \mathcal{D}_{\mathcal{J}} \oplus B \right] \simeq 2^{-30} + 2^{-51.985}.$$

The analysis is described in Appendix C, and is again conducted in Setting (2). A proof for the mean in Setting (1) of independent δ -sets can use the same argumentation.

Let Y be a random variable for the number of single-column collisions after four rounds. Both settings consider the distribution of k collisions from N pairs. In Setting (1), this distribution is described approximately by a binomial distribution, $Y \sim \mathcal{B}(N, p)$, where p as above.

Again, in Setting (2), the distribution is more complex. Here, it holds that $N = x \cdot 4 \cdot 2^{24} \cdot \binom{2^8}{2}$ stems from x diagonal spaces $\mathcal{D}_{\{c\}}$. Since this setting maintains a multiple-of-eight

property, it is reasonable to think that the distribution would approximately follow

$$\Pr[Y = \ell] = \begin{cases} 0 & \text{if } \ell \bmod 8 \neq 0, \\ \binom{N_3}{\ell_3} (p_3)^{\ell_3} (1 - p_3)^{N_3 - \ell_3} & \text{otherwise,} \end{cases}$$

where $\ell = 2^3 \cdot \ell_3$, $N_3 = 2^{28} \cdot 255^4$ and $p_3 = 2^{-30} + 2^{-51.985}$. The factor ℓ_3 represents the number of pairs in the input space that follow the distribution.

Steps. The steps of our distinguisher are as follows:

1. Initialize a collision counter.
2. For $i = 1..2^8$, collect a structure \mathcal{S} of 2^{32} texts that iterate over all values in any four bytes and leave the remaining bytes constants. Query the plaintexts of a structure and ask for their ciphertexts. Invert the final **ShiftRows** operations to get the states $S^{5,SB}$ and store them in some list \mathcal{Q} .
3. Form $4 \cdot 2^{24}$ δ -sets from a structure, i.e., the texts in each δ -set iterate over the 2^8 values in one byte and are constant in the 15 remaining bytes. For each δ -set:
 - 3.1 Initialize four lists \mathcal{L}_i , for $i = 0, 1, 2, 3$ of 2^{32} elements.
 - 3.2 For each column i of $S^{5,SB}$, interpret the column as 32-bit integer and append the text to \mathcal{L}_i the index corresponding to the column value for each list.
4. For each of the lists \mathcal{L}_i :
 - 4.1 Look for collisions, e.g., multiple values at the same index.
 - 4.2 For each collision, look those pairs up in other columns if they have already been counted. If not, increment the counter.
5. If the counter is higher than the threshold θ , output real; otherwise, output random.

We can use the approximation for single differential characteristics from [SS17] for an approximation for threshold. The success probability corresponds to $1 - \alpha$, where α is the probability of the Type-I error, that the distinguisher decides that it were a random permutation, but is the AES. A possible choice for a threshold is, according to [SS17],

$$\theta = N \cdot |p_{\text{AES}} - p_{\text{rand}}| - \sqrt{3 \cdot N \cdot p_{\text{AES}} \cdot \ln\left(\frac{1}{\alpha}\right)} \quad (13)$$

Setting $P_S = 0.95$ implies $\alpha = 0.05$.

Complexity. Consider Setting (1) first. For a success probability of approximately $P_S = 0.95$, we obtain $N \geq 2^{76.406}$; for $P_S = 0.99$, we obtain $N \geq 2^{77.406}$ pairs. By fixing any set of four distinct bytes in the plaintexts, we can form 2^{32} texts by iterating over all values of those four bytes and leaving the remaining twelve bytes constant for a structure. So, one structure provides

$$4 \cdot 2^{24} \cdot \binom{2^8}{2} \simeq 2^{40.994} \simeq 2^{41} \text{ pairs.}$$

2^{36} structures with 2^{68} chosen plaintexts with approximately 2^{77} pairs suffice for a success probability of more than 0.95 on average. The memory complexity is given by storing 2^{32} states in \mathcal{Q} and four lists \mathcal{L}_i of $4 \cdot 2^{32}$ columns at a time. The time complexity consists of

- 2^{68} encryptions and $1/5$ partial decryptions or $2^{68.3}$ encryption equivalents.
- $2^{68} + 2^{68} \cdot 4 \simeq 2^{70.4}$ memory accesses for inserting the texts. Depending on the data structure, insertions or lookups need sorting, which requires $N \log_2(N)$ operations on average with e.g., quicksort. Given four lists of 2^8 elements, sorting needs approximately $2^{36} \cdot 2^{24} \cdot 4 \cdot 2^8 \cdot 8 \simeq 2^{73}$ memory accesses.
- At most $2^{36} \cdot 2^{24} \cdot 4 \cdot \binom{2^8}{2} \cdot 2^{-32} \cdot 3 \simeq 2^{46.6}$ additional memory accesses when collisions occur to look up if other columns collide in the other lists.
- Approximately $2^{36} \cdot 2^{24} \cdot 4 \cdot \binom{2^8}{2} \cdot 2^{-32} \simeq 2^{46}$ memory accesses to increment the counter.

So, the computational complexity can be approximated by $2^{70.4} + 2^{73} + 2^{46.6} + 2^{46} \simeq 2^{73.3}$ memory accesses and $2^{68.3}$ encryptions. From Equation (13), we obtain as an approximation for θ for $P_S = 0.95$

$$\theta \simeq 2^{76.406} \cdot 2^{-51.985} - \sqrt{3 \cdot 2^{76.406} \cdot 2^{-51.985} \cdot \ln(1/0.05)} \simeq 2^{24.420}.$$

Setting (2) inherits the multiple-of-eight property. We study its variance in Appendix E. The variance for the AES is approximately a factor of eight higher in this setting than that for a random permutation. Therefore, the data complexity increases theoretically by a factor of about $2 \cdot \sigma_{\text{AES}}^2 / \sigma_{\text{rand}}^2 \approx 16$. For $P_S = 0.95$, we obtain $2^{80.402}$ pairs and therefore a data complexity of approximately 2^{40} structures with 2^{72} chosen plaintexts. For $P_S = 0.99$, we obtain $N > 2^{81.406}$ pairs, and therefore a data complexity of approximately 2^{73} chosen plaintexts. From Equation (13), we obtain as an approximation for θ for $P_S = 0.95$

$$\theta \simeq 2^{80.402} \cdot 2^{-51.985} - \sqrt{3 \cdot 2^{80.402} \cdot 2^{-51.985} \cdot \ln(1/0.05)} \simeq 2^{28.417}.$$

Remark 1. The distinguisher employs structures composed of 2^{32} texts, e.g., from diagonals each, which are close to our analyzed setting. A more efficient approach used in an earlier version of this paper could be to employ structures of 2^{64} texts that iterate over eight bytes, and form pairs from the $8 \cdot 2^{56} \cdot \binom{2^8}{2}$ structures. This approach would approximately halve the necessary data and computational complexities. Its disadvantage is an approximating assumption that the so-created structures are independent. We decided in favor of the more conservative approach here.

Reduced Variant. In order to allow a practical verification, we derive the corresponding probabilities for Small-AES. Again, we approximate the downscaled variant of Mix-Columns as a sum of four independent permutations on \mathbb{F}_{2^4} . The expected probability that all four bytes in a diagonal space $\mathcal{D}_{\{c\}}$ collide for two texts in a δ -set can be approximated by $\mathbb{E}[N_4]^4$:

$$\Pr_{\text{Small-AES}} [S^3 \in \mathcal{D}_{\{c\}}] = \left(2^{-4} + \frac{1}{2^4(2^4 - 1)^3} \right)^4 \simeq (2^{-4} + 2^{-15.721})^4,$$

and that for a random permutation by

$$\Pr_{\text{rand}} [S^3 \in \mathcal{D}_{\{c\}}] = \frac{2^{48} - 1}{2^{64} - 1} \simeq 2^{-16} - 2^{-64}.$$

Then, the probability to have at least one all-zero anti-diagonal in the difference is given for five-round Small-AES by

$$p_{\text{Small-AES}} \simeq 1 - \left(1 - \Pr_{\text{Small-AES}} [S^3 \in \mathcal{D}_{\{c\}}] \right)^4 \simeq 2^{-14} + 2^{-23.748},$$

whereas for a random permutation, it is approximately

$$p_{\text{rand}} \simeq 1 - \left(1 - \Pr_{\text{rand}} [S^3 \in \mathcal{D}_{\{c\}}]\right)^4 \simeq 2^{-14} - 2^{-29.415}.$$

In Setting (1), Equation (4) yields, for a success probability of approximately $P_S = 0.95$, that the required number of pairs would have to be $N \geq 2^{35.934}$, and $N \geq 2^{36.934}$ pairs for $P_S = 0.99$. 2^{16} structures of $(2^4)^4 = 2^{16}$ texts each that iterates over four cells of a diagonal for each structure can form $2^{16} \cdot 4 \cdot 2^{12}$ sets of $\binom{2^4}{2}$ pairs each, which corresponds to

$$2^{16} \cdot 4 \cdot 2^{12} \cdot \binom{2^4}{2} \simeq 2^{36.907}$$

pairs. So, 2^{16} structures of 2^{16} texts each may suffice to have a success probability of $P_S \geq 0.95$.

In Setting (2), the variance should be eight times higher, and therefore, $N \geq 2\sigma_{\text{Small-AES}}^2 / \sigma_{\text{rand}}^2 = 16$ should be 16 times higher. Equation (4) yields, for a success probability of approximately $P_S = 0.95$, that the required number of pairs would have to be $N \geq 2^{39.934}$, and $N \geq 2^{40.934}$ pairs for $P_S = 0.99$. So, 2^{20} structures of 2^{16} texts each may suffice to have a success probability of $P_S \geq 0.95$.

Testing on A Single Diagonal. A variant of the distinguisher can also consider only a fixed inverse diagonal after almost five rounds. In this case, the probabilities would become

$$p_{\text{AES}} \simeq \Pr_{\text{AES}} [S^3 \in \mathcal{D}_{\{c\}}] \simeq 2^{-32} + 2^{-53.983},$$

whereas for a random permutation, it is approximately

$$p_{\text{rand}} \simeq \Pr_{\text{rand}} [S^3 \in \mathcal{D}_{\{c\}}] \simeq 2^{-32} - 2^{-128}.$$

Then, the resulting data complexities for $P_S = 0.95$ would become $N \geq 2^{78.402}$ and $N \geq 2^{79.402}$ for $P_S = 0.99$. So, the data complexity would become 2^{70} and 2^{71} chosen plaintexts, respectively.

For Small-AES, the corresponding complexities would become $N \geq 2^{37.878}$ and $N \geq 2^{38.878}$ pairs, i.e., 2^{33} and 2^{34} chosen plaintexts would be required. Therefore, our distinguisher(s) test(s) on any inverse diagonal.

5.3 Verification

Experimental Verification. Again, we tried to verify our distinguishers experimentally with Small-AES. Figure 8 depicts our results with 100 random keys and 2^{30} random δ -sets of 2^4 texts each. All values μ denote the means for the number of pairs in δ -sets that collided in at least one inverse diagonal after five rounds, over all δ sets, e.g., $\mu = \binom{2^4}{2} \cdot 2^{30} \cdot p_{\text{Small-AES}}$ for Small-AES for 2^{30} δ -sets. In addition, we tested the Small-AES with the S-box of PRESENT. Like for our four-round distinguisher, we can observe the following:

- Small-AES deviates significantly from the theory, in the sense that the distinguisher is stronger. Here, its experimental mean is approximately one standard deviation higher than the theoretical prediction.

Instance	μ	σ
π		
Theory	7 864 140	2 804.22
Experiment	7 864 379.	2 492.46
Small-AES		
Theory	7 873 286	2 805.85
Experiments	7 875 860.	2 844.95
PRESENT S-box	7 868 881.	2 785.78

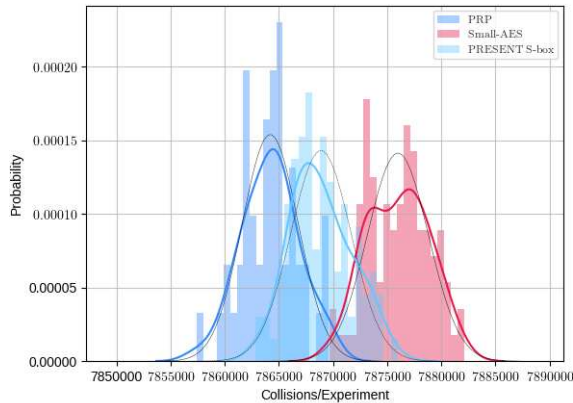


Figure 8: Means and standard deviations for the number of pairs that collided in at least one inverse diagonal for our five-round truncated-differential with Small-AES. Each experiment employed 100 random independent keys and 2^{30} random δ -sets. Speck-64-96 has been used as PRP π . We compare to Small-AES with the PRESENT S-box.

- Small-AES with the S-box of PRESENT has again a strong deviation, here by approximately 1.6 standard-deviation units below the theoretical prediction. As for our four-round distinguisher, it still yields a distinguisher, its mean is approximately 1.7 standard-deviation units above the mean of the PRP, i.e., roughly in the middle between both theoretical predictions.

We conclude that the true distinguisher, at least for Small-AES, depends largely on the internals of the S-box.

Transition-based Verification. Again, we employ the transition-based approach by Rønjom [Røn19] for a second form of verification, using 100-bit precision. We obtained the probabilities

$$\begin{aligned}
 p_{\text{AES}} &\simeq 0.2246672325602152 \dots \cdot 10^{-15} \simeq 2^{-30} + 2^{-51.98306} \\
 p_{\text{Small-AES}} &\simeq 0.7244634283789227 \dots \cdot 10^{-7} \simeq 2^{-14} + 2^{-23.71851}.
 \end{aligned}$$

The results basically confirmed the biases of our five-round distinguishers. The slight deviations can be due to rounding.

A Hidden Property? Again, we also determined the variance in both settings experimentally. The results are given in Appendix E. One would expect that the variance in Setting (2) is about eight times higher than that in Setting (1) due to the multiple-of-eight property. Though, we found that it is roughly 11.4 times larger as outlined in Table 10, i.e., roughly 12 times larger than the variance of Setting (1). The reason for this gap is unknown.

The results give rise to the possibility that there exists some further multiple-of-four property in the distribution that may be hidden by the multiple-of-eight property. Note that this property affects the mean of our distinguisher only marginally. The variance, however, differs significantly, as we could observe experimentally.

S-box Dependencies. Again, we emphasize that our assumptions approximate the structures of AES and Small-AES. For consistency, we iterated again over the cell at index 1

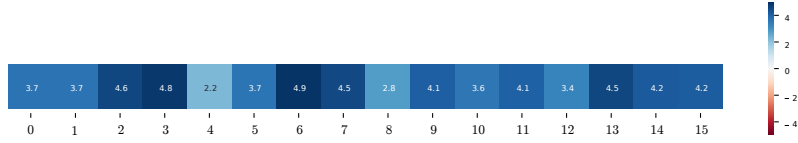


Figure 9: Average means of five experiments with random keys and 2^{30} δ -sets for our five-round distinguisher with Small-AES in $\times\sigma$ deviations from the mean of the PRP. Each cell represents the result for the active input cell in the plaintexts.

in the plaintexts in Setting (1) in our distinguisher on five rounds. Since our distinguisher starts again from a single cell, we expect weaker deviations due to the additional round compared to our four-round distinguisher and due to the fact that it considers four cells after three full rounds at a time. Figure 9 presents our preliminary results. One can observe that all choices for the active input cell yield distinguishers that are at least 2.2σ above the mean of the PRP, with peaks at 4.9σ .

6 Six-round Key-recovery Attack

We can extend the five-round distinguisher from Section 5 to a key-recovery attack on six-round AES that recovers 32 key bits. We apply the attack twice in shifted form to recover 64 key bits and the remaining key bits exhaustively. Table 6 compares it with the state-of-the-art of attacks on six- and seven-round AES-128. Note that various more attacks exist on the AES versions with larger keys.

Precomputation. We construct four tables \mathcal{H}_i for $0 \leq i \leq 3$. For \mathcal{H}_i , we store pairs $(X, \Delta X) \in \mathbb{F}_{2^8}^4 \times \mathbb{F}_{2^8}^4$ such that $R(X) \oplus R(X \oplus \Delta X)$ is non-zero only in Byte i of the output column, where \tilde{R} is the AES round function missing `AddRoundKeys` and restricted to work on one diagonal and output the corresponding column (note that the computation on each diagonal of the state can be done independently of the computation on other diagonals). The table \mathcal{H}_i can be constructed as follows: for each non-zero value y of the active byte ΔY , compute $\Delta O = (\Delta O_0, \Delta O_1, \Delta O_2, \Delta O_3) = MC^{-1}(y \parallel 0 \parallel 0 \parallel 0) \ggg i$. For each ΔO_j ($0 \leq j \leq 3$), check the DDT of the AES S-box which is integrated with the information of correct pairs for each compatible difference, one can get 256 pairs (X_j, X'_j) fulfilling the output difference ΔO_j . For ΔO , obtain all possible combination between (X_0, X'_0) , (X_1, X'_1) , (X_2, X'_2) , (X_3, X'_3) . Each combination corresponds to a valid $(X, \Delta X)$. Thus, one can get 256^4 valid pairs $(X, \Delta X)$ for a non-zero value of the active byte ΔY . In total, each table \mathcal{H}_i contains $255 \cdot 256^4$ entries. Four tables then need $4 \cdot 255 \cdot 256^4 \cdot 8$ bytes, which corresponds to approximately $4 \cdot 255 \cdot 256^4 \cdot 8/16 \simeq 2^{40.1}$ state equivalents.

Steps. The steps are as follows:

1. Initialize a zeroed list \mathcal{K} for the 2^{32} key candidates for $K^0[0, 5, 10, 15]$.
2. Precompute the tables \mathcal{H}_i , for $0 \leq i \leq 3$.
3. For $i = 1..2^s$, collect a structure \mathcal{S} of 2^{32} texts in a coset of $\mathcal{D}_{\{0\}}$.
 - 3.1 So, the texts in each structure iterate over the 2^{32} values in Bytes $P[0, 5, 10, 15]$ and are constant in the 12 remaining bytes.
 - 3.2 Initialize four lists \mathcal{L}_i , for $i = 0, 1, 2, 3$. Query the plaintexts of a structure to an encryption oracle to obtain the corresponding ciphertexts C . Invert the final `ShiftRows` operation to get the states $S^{6, SB}$ and store the tuples of

$(P[0, 5, 10, 15], S^{6, \text{SB}})$ into the four lists \mathcal{L}_i , where the texts are indexed by the i -th anti-diagonal of $S^{6, \text{SB}}$.

3.3 For $j = 0..3$, consider the lists \mathcal{L}_j :

- i. Look up collisions on the j -th anti-diagonal of $(S^{6, \text{SB}}, S'^{6, \text{SB}})$. For each collision, look up in the lists with lower indices j if the same pair collided already in a different anti-diagonal to prevent double counting.
- ii. For each collision, consider their corresponding plaintexts P and P' and derive their difference on the diagonal $\Delta P[0, 5, 10, 15] = P[0, 5, 10, 15] \oplus P'[0, 5, 10, 15]$. From $\mathcal{H}_0, \mathcal{H}_1, \mathcal{H}_2$, and \mathcal{H}_3 , look up all possible pairs $(X, \Delta X)$ such that $\Delta P[0, 5, 10, 15] = \Delta X$. Then use the corresponding $P[0, 5, 10, 15]$ and X to derive the candidate, i.e., $K^0[0, 5, 10, 15] = P[0, 5, 10, 15] \oplus X$, increase its counter in \mathcal{K} .

4. Sort the list of key candidates and output the sorted list.
5. Apply the attack another time from a shifted diagonal, e.g., $\mathcal{D}_{\{1\}}$, and with another 2^s structures to recover another 32 bits of key material.
6. Test the keys in descending order of their counters to recover the remaining 64 key bits.

Complexity. Each structure yields $4 \cdot 2^{24}$ δ -sets, which corresponds to approximately 2^{41} pairs. We choose 2^s structures, which corresponds to 2^{s+32} chosen plaintexts

The time complexity is given by the following:

- The precomputation costs $4 \cdot 255 \cdot 256^4 \cdot 8/16 \cdot 1/6 \simeq 2^{38.5}$ encryption equivalents and $2^{40.1}$ memory accesses for storing the values.
- Encrypting the data takes 2^{s+32} encryptions and $1/6$ partial decryptions, or at most $2^{s+32.3}$ encryption equivalents.
- We need $4 \cdot 2^{s+32} = 2^{s+34}$ memory accesses for storing the 2^{s+32} texts into four lists. Sorting the four lists requires $2^s \cdot 4 \cdot 2^{32} \cdot 32 \simeq 2^{s+39}$ memory accesses.
- Each structure yields about 2^{63} pairs, and each pair has a probability of approximately 2^{-30} to collide on a column. So, we expect approximately $2^{s+63} \cdot 2^{-30} \simeq 2^{s+33}$ collisions. For each collision that occurs in a table \mathcal{L}_i , we need at most three memory accesses for the three other tables in $\mathcal{L}_j, j \neq i$, to prevent double counting. Thus, we need $2^{s+33} \cdot 3 \simeq 2^{s+34.6}$ MAs at that point.
- For each collision, we expect on average $4 \cdot 255$ suggestions (because the match $\Delta P[0, 5, 10, 15] = \Delta X$ imposes 32 bits conditions) from the hash tables \mathcal{H}_i . So, we need about $2^{s+33} \cdot 4 \cdot 255 \simeq 2^{s+43.0}$ memory accesses plus the same amount in \mathcal{K} , plus the same amount of XORs.

Since the attack is performed twice with the same plaintext material but on 32 different key bits, e.g., $K[3, 4, 9, 14]$, we need the same number of encryptions (alternatively, they could be stored) and lookups, plus 2^{64} encryptions for the key search. Let a be the attack advantage. The total time complexity consists of

$$2^{38.5} + 2 \cdot 2^{s+32.3} + 2^{64+2 \cdot (32-a)} \text{ Encryptions and } \\ 2^{40.1} + 2 \cdot (2^{s+34} + 2^{s+39} + 2^{s+34.6} + 2^{s+43.0}) \text{ MAs.}$$

The attack stores 2^{32} tuples of 4 + 16 bytes at a time for the states plus 2^{32} byte counters candidates plus $2^{40.1}$ states for the tables \mathcal{H}_i . So, the memory complexity is circa $2^{40.1}$ states.

6.1 Relations between Success Probability, Advantage and Data

In Step 6, we apply the key-ranking approach, i.e. candidates of the attacked key bits are tested in the order of their decreasing counts. Thus, before reaching the correct key, all candidates ranked above are also checked with combinations of the remaining key bits to see if they are the correct value. Suppose an m -bit key is attacked and the correct key candidate is ranked ℓ -th among all 2^m candidates, then the key-ranking method reduces the complexity by a factor of $2^{m-\lg \ell}$ over the exhaustive search. Denote by $a = m - \lg \ell$, this a is referred to as the advantage by Selçuk in [Sel08].

Remark 2. We note that the original model by Selçuk considers only a single list. The applicability to an attack that reduces the key in two independent columns may be more complex.

Using the Model in [Sel08]. In [Sel08], Selçuk derived formulas to present the relation between the advantage a , the success probability P_s (the probability that the correct key candidate is ranked among the 2^{m-a} first candidates), and the data complexity N both for linear and differential cryptanalysis. The results were derived on order statistics and are applicable for normally distributed random variables. The derived formula is less accurate for differential than that for linear cryptanalysis, mainly because for classical differential cryptanalysis, using the normal distribution to approximate binomial distribution is not quite accurate. However, in the presented truncated differential attack, the binomial distribution of the random variables – the number of collisions for right key (denoted as X_R , i.e., the number of collisions for the 5-round distinguisher) and that for wrong keys (denoted as X_W , i.e., the number of collisions for random permutations) – are allowed to be approximated with normal distributions. Recall that we have $X_R \sim \mathcal{N}(\mu_R, \sigma_R^2)$ and $X_W \sim \mathcal{N}(\mu_W, \sigma_W^2)$, where

$$\begin{aligned} \mu_R &= N \cdot p_{\text{AES}}, & \sigma_R &= N \cdot p_{\text{AES}} \cdot (1 - p_{\text{AES}}), & \text{where } p_{\text{AES}} &\simeq 2^{-30} + 2^{-51.985} \\ \mu_W &= N \cdot p_{\text{rand}}, & \sigma_W &= N \cdot p_{\text{rand}} \cdot (1 - p_{\text{rand}}), & \text{where } p_{\text{rand}} &\simeq 2^{-30} - 2^{-61.415}. \end{aligned}$$

Plugging this configuration into the following Theorem 3, we can draw the trade-off curves (Figure 10a) between the advantage a , the success probability P_s , and the data complexity (note that there is a difference between the required number of plaintext pairs and the required number of plaintexts due to the using of plaintext structures) for the key-recovery attack on six-round AES.

Theorem 3 ([Sel08]). Let the statistic for a correct key candidate follow the normal distribution $\mathcal{N}(\mu_R, \sigma_R^2)$, and the statistic for wrong key candidate follow $\mathcal{N}(\mu_W, \sigma_W^2)$. Let m be the number of attacked key bits, a be the advantage of an attack, and N be the number of available samples. Then, the success probability of the attack P_s can be approximated by:

$$P_s \simeq \Phi_{0,1} \left(\frac{\mu_R - \mu_a}{\sqrt{\sigma_R^2 + \sigma_a^2}} \right),$$

where $\mu_a = \mu_W + \sigma_W \Phi_{0,1}^{-1}(1 - 2^{-a})$, and $\sigma_a^2 \simeq \frac{\sigma_W^2 \cdot 2^{-(m+a)}}{\phi_{0,1}^2(\Phi_{0,1}^{-1}(1 - 2^{-a}))}$.

Using the Model in [SS17]. In [SS17], Samajder and Sarkar provide closed-form *upper bounds* on the data complexity in terms of the success probability and the advantage of an multiple (truncated) differential attack (the upper bound on the data complexity shows how much data is *sufficient* to mount an attack with specified advantage and success probability). By employing the Chernoff bounds, they do not need any approximations (e.g.,

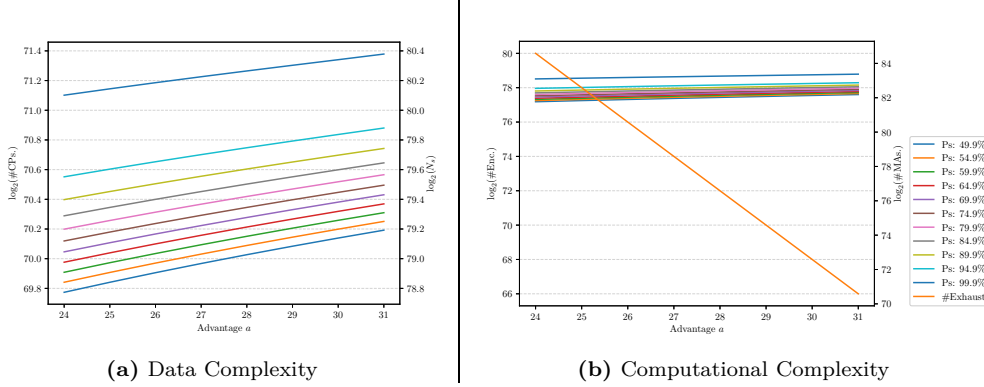


Figure 10: Relations between success probability, advantage, data complexity, and computational complexity of the attack using the model in [Sel08]. Note that in our attack, $\log_2(\#CPs.) = \log_2(N) - 41 + 32$, and we assume $\#Encs. \simeq 6 \cdot 4 \cdot \#MAs.$

the Gaussian approximation of the binomial distribution or the Poisson approximation of the binomial distribution tails). Using their result presented below (Proposition 1) and substituting p_1 with p_{AES} , and q_1 with p_{rand} , we obtain the trade-off curves (Figure 11a) between the advantage a , the success probability P_s , and the upper bound on the data complexity (note that there is a difference between the required number of plaintext pairs N and the required number of chosen plaintexts due to the use of plaintext structures) for the key-recovery attack on six-round AES.

Proposition 1 ([SS17]). Let $0 < \alpha, \beta < 1$ and N be such that

$$N \geq \frac{3 \left(\sqrt{p_1 \ln(1/\alpha)} + \sqrt{q_1 \ln(1/\beta)} \right)^2}{(p_1 - q_1)^2}.$$

Then the probabilities of the Type-1 and Type-2 errors in Hypothesis Test-1 are upper bounded by α and β respectively. Putting $\alpha = 1 - P_s$ and $\beta = 2^{-a}$, it follows that for

$$N \geq \frac{3 \left(\sqrt{p_1 \ln(1 - P_s)} + \sqrt{aq_1 \ln(2)} \right)^2}{(p_1 - q_1)^2}.$$

the success probability will be at least P_s and the advantage will be at least a .

Optimal Computational Complexity. To optimize the computational complexity of the attack, we need to balance the time of the exhaustive search (combining 2^{m-a} candidates ranked at the top of each of the two lists with the 64 remaining key bits. So, we can see if they are the correct value) and that of the attack steps (dominated by the $2 \cdot 2^{s+63-30} \cdot 4 \cdot 255 \simeq 2^{s+44.0}$ memory accesses), where $m = 32$ is the number of attacked key bits, a is the advantage, 2^s is the number of plaintext structure ($s = \log_2(N) - 41$, where N is the required number of pairs and 2^{41} is the number of pairs in each structures). To do that, we assume that one 6-round AES encryption (Encs.) costs equivalently to $6 \times 4 \simeq 2^{4.585}$ times of memory access (MAs.) to a data structure of size no more than 2^{32} . Thus, we need to find a balance between $2^{64} \cdot 2^{2 \cdot (m-a)} = 2^{128-2a}$ and $2 \cdot 2^{\log_2(N)-41+44.0} / (6 \cdot 4) \simeq 2^{\log_2(N)-1.591}$.

Figure 10b shows the trade-off curves under different success probability based on the result obtained using the model in [Sel08], and Figure 11b shows that using the model in [SS17]. With a moderate success probability 63% (note that the success probability

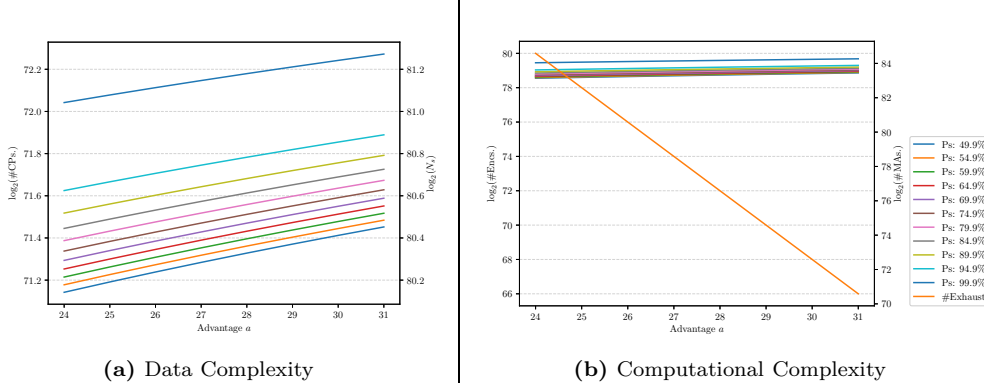


Figure 11: Relations between success probability, advantage, data complexity, and computational complexity of the attack using the model in [SS17]. Note that in our attack, $\log_2(\#CPs.) = \log_2(N) - 41 + 32$, and we assume $\#Encs. \simeq 6 \cdot 4 \cdot \#MAs.$

of the six-round key-recovery attack in [BODK⁺18] is approximately 63%), based on the model by [Sel08], we pick the balance point with $a = 25.5$. At that point, the required number of plaintext pairs is $N = 2^{79.045}$, the data complexity is approximately $2^{70.045}$ chosen plaintexts. The computational complexity is approximately $2^{77.454}$ encryption equivalents. Using the model from Samajder and Sarkar [SS17], we choose the balance point with $a = 25$. At that point, the required number of plaintext pairs is $N = 2^{80.285}$, the data complexity is approximately $2^{71.285}$ chosen plaintexts, and the computational complexity is approximately $2^{78.695}$ six-round encryptions. In Table 6, we chose the more conservative higher complexities from the model by [SS17].

6.2 Experimental Verification

We verified our attack in a variant that omits the second execution on Small-AES. The resulting S-box is inherited 4-differential-uniformity; for each input difference α , there exist six output differences β with $\Pr[S(x) \oplus S(x \oplus \alpha) = \beta] = 2/16$, and one difference β for which $\Pr[S(x) \oplus S(x \oplus \alpha) = \beta] = 4/16$. Similar properties hold for the inverse S-box for Small-AES. So, each hash table \mathcal{H}_i , for $0 \leq i \leq 3$ contains 2^{16} 16-bit entries for each input difference ΔP on average. This corresponds to 2^{20} bytes or 2^{17} 64-bit states for Small-AES.

We chose 2^s structures of 2^{16} texts each, which yielded $2^s \cdot 4 \cdot 2^{12}$ δ -sets. For $s = 16$, this corresponds to approximately $2^{36.907}$ pairs. The time complexity is given by

- $4 \cdot 15 \cdot 15 \cdot 2^{16} \cdot 8/16 \cdot 1/6 \simeq 2^{22.3}$ encryptions for the hash tables.
- 2^{32} encryptions and $1/6$ partial decryptions, or at most $2^{32.22}$ encryption equivalents.
- 2^{34} memory accesses for storing the texts and $4 \cdot 2^{16} \cdot 16 \cdot 2^{16} = 2^{38}$ memory accesses for sorting the lists.
- It yields about $2^{16} \cdot 4 \cdot 2^{12} \cdot \binom{2^4}{2} \cdot 2^{-14} \simeq 2^{22.91}$ collisions. Given that we need at most three accesses to the lists \mathcal{L} per collision, this corresponds to at most $2^{25.49}$ memory accesses.
- Per collision, we expect on average $4 \cdot 15$ suggestions from the hash tables. So, we require $2 \cdot (2^{22.91} \cdot 4 \cdot 15) \simeq 2^{29.82}$ XORs for deriving the keys in \mathcal{K} and the same maximum number of MAs to increment the counters in \mathcal{K} .

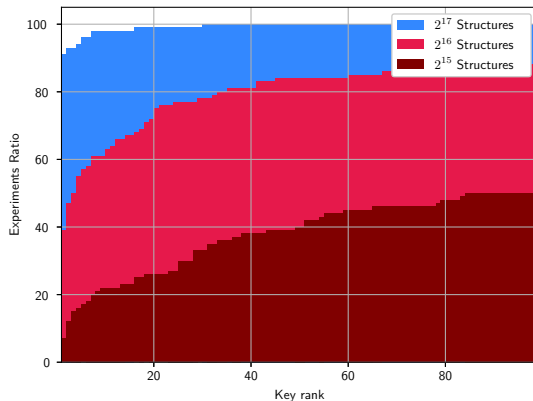


Figure 12: Rank distribution for the correct key among 2^{16} candidates from 100 runs of our six-round attack with Small-AES, with random keys and 2^{15} , 2^{16} , or 2^{17} structures of 2^{16} texts each.

Table 6: Comparison of the best secret-key key-recovery attacks on six-round AES-128 and the best attacks on seven rounds, ordered by rounds (descending) then time (ascending). Enc. = encryptions; MAs = memory accesses; CP = chosen plaintexts.

#Rds.	Attack Type	Time	CP	Ref.
6	Impossible Differential	$2^{122.0}$ Enc.	$2^{91.5}$	[CKK ⁺ 01]
6	MitM	$2^{106.2}$ Enc.	2^8	[DFJ13]
6	Prob. Mixture-differential	$2^{105.0}$ Enc.	$2^{72.8}$	[Gra17, Gra19]
6	Mixture-differential	$2^{81.0}$ Enc.	$2^{27.5}$	[BODK ⁺ 18]
6	Truncated differential	$2^{78.7}$ Enc.	$2^{71.3}$	Sect. 6
6	Integral	$2^{51.7}$ Enc.	2^{35}	[Tod14, TA14]
6	Partial Sum	$2^{42.0}$ Enc.	2^{32}	[Tun12a, Tun12b]
7	Impossible Differential	$2^{106.88}$ Enc.	2^{105}	[BLNS18]
7	MitM	$2^{99.0}$ Enc.	2^{97}	[DFJ13]

We omit the second execution; the attack requires

$$2^{22.3} + 2^{32.22} \simeq 2^{32.23} \text{ encryptions and } 2^{34} + 2^{38} + 2^{24.49} + 2^{29.82} \simeq 2^{38.1} \text{ MAs}$$

to reduce the key space for $K^0[0, 5, 10, 15]$. The attack stores 2^{16} tuples of 4 + 16 nibbles at a time for the states plus 2^{16} four-nibble key candidates, plus 2^{17} states for the hash tables, which corresponds to less than $2^{17.9}$ states.

Figure 12 illustrates the results of 100 experiments that employed independent random keys for 2^{15} , 2^{16} , and 2^{17} structures of 2^{16} plaintexts each. Our experiments aimed at recovering the 16 key bits $K^0[0, 5, 10, 15]$ from the first diagonal. The advantages for the experiments yielded 7.070, 10.550, and 14.243 bits, respectively.

7 Six-round Truncated-differential Distinguisher

We can derive a six-round distinguisher that shares the higher-level concept of our six-round key-recovery attack.

Core Idea. In our five-round distinguisher, we started from a single-byte difference that was transformed through five rounds to a mixed space $\mathcal{M}_{\mathcal{J}}$ for arbitrary $\mathcal{J} \subseteq \{0, 1, 2, 3\}$ and $|\mathcal{J}| = 3$. The idea is to extend the five-round distinguisher at the beginning by starting from a single-diagonal space. Let $X, X' \in (\mathbb{F}_{2^8})^{4 \times 4}$ be distinct and $\Delta X = X \oplus X' \in \mathcal{D}_{\{c\}}$ for some fixed $c \in \{0, 1, 2, 3\}$. Let Y and Y' be the corresponding ciphertexts of X and X' , respectively; either generated by transforming them through six-round AES, or through a random permutation $\pi \leftarrow \text{Perm}((\mathbb{F}_{2^8})^{4 \times 4})$. We define the space $\mathcal{B}_{\{r,c\}} \stackrel{\text{def}}{=} \mathcal{C}_{\{c\}} \cap \mathcal{D}_{\{r\}}$, i.e., the space of the single Byte in Row r and Column c for $r, c \in \{0, 1, 2, 3\}$. With a certain probability p^1 , $\mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,0\}}$ possesses a single-byte difference after the first round. Then, the probability that the difference in their ciphertexts is in $\Delta Y \in \mathcal{M}_{\mathcal{J}}$ is approximately the probability of our five-round distinguisher, let's say p^5 . Let us define the random variable Y to be 1 iff $Y \oplus Y' \in \mathcal{M}_{\mathcal{J}}$ for arbitrary $\mathcal{J} \subseteq \{0, 1, 2, 3\}$ and $|\mathcal{J}| = 3$, and 0 otherwise. In the case that multiple bytes are active in the difference after the first round, we can approximate the probability of Y by that of random permutation. For the random permutation, the probability is always q ignorant of the plaintext pair. So, for a fraction of pairs, we have a higher probability that their ciphertext difference is in $\mathcal{M}_{\mathcal{J}}$. If p^1 is not too small, we can build a distinguisher with smaller mean than before, which can still be used in a distinguisher.

Formulation. In the following, we omit writing the condition $X \oplus X' \in \mathcal{D}_{\{c\}}$ explicitly for brevity. By the law of total probability,

$$\begin{aligned} & \Pr[Y \oplus Y' \in \mathcal{M}_{\mathcal{J}}] \\ &= \Pr[\mathbb{R}^6(X) \oplus \mathbb{R}^6(X') \in \mathcal{M}_{\mathcal{J}}] \\ &= \Pr[\mathbb{R}^6(X) \oplus \mathbb{R}^6(X') \in \mathcal{M}_{\mathcal{J}} | \mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}}] \cdot \Pr[\mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}}] + \\ & \quad \Pr[\mathbb{R}^6(X) \oplus \mathbb{R}^6(X') \in \mathcal{M}_{\mathcal{J}} | \mathbb{R}(X) \oplus \mathbb{R}(X') \notin \mathcal{B}_{\{r,c\}}] \cdot \Pr[\mathbb{R}(X) \oplus \mathbb{R}(X') \notin \mathcal{B}_{\{r,c\}}]. \end{aligned}$$

Under the Markov-cipher assumption, we know from our five-round distinguisher that

$$\begin{aligned} p_{\text{AES}}^5 & \stackrel{\text{def}}{=} \Pr[\mathbb{R}^6(X) \oplus \mathbb{R}^6(X') \in \mathcal{M}_{\mathcal{J}} | \mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}}] \\ & \simeq 2^{-30} + 2^{-51.985} - 2^{-61.415}. \end{aligned}$$

For $X \in (\mathbb{F}_{2^8})^{4 \times 4}$, let $\text{wt} : (\mathbb{F}_{2^8})^{4 \times 4} \rightarrow \mathbb{Z}_{16}$ represent the byte weight, which yields the number of non-zero cells in the input:

$$\text{wt}(X) \stackrel{\text{def}}{=} \sum_{i=0}^{15} \delta(X[i]) \quad \text{where} \quad \delta(X[i]) \stackrel{\text{def}}{=} \begin{cases} 1 & \text{if } X[i] \neq 0, \\ 0 & \text{otherwise.} \end{cases}$$

The probability for a difference with a single active byte after the first round is simply

$$\begin{aligned} p_{\text{AES}}^1 & \stackrel{\text{def}}{=} \Pr[\mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}}] \tag{14} \\ &= \Pr[\mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}} | \text{wt}(X \oplus X') = 4] \cdot \Pr[\text{wt}(X \oplus X') = 4] + \\ & \quad \Pr[\mathbb{R}(X) \oplus \mathbb{R}(X') \in \mathcal{B}_{\{r,c\}} | \text{wt}(X \oplus X') < 4] \cdot \Pr[\text{wt}(X \oplus X') < 4] \\ &= 4 \cdot \frac{255}{255^4} \cdot \frac{255^4}{2^{32} - 1} + 0 = \frac{4 \cdot 255}{2^{32} - 1}. \end{aligned}$$

Furthermore, we approximate

$$\begin{aligned} & \Pr[\mathbb{R}^6(X) \oplus \mathbb{R}^6(X') \in \mathcal{M}_{\mathcal{J}} | \mathbb{R}(X) \oplus \mathbb{R}(X') \notin \mathcal{B}_{\{r,c\}}] \\ & \simeq \Pr[\pi(X) \oplus \pi(X') \in \mathcal{M}_{\mathcal{J}}] \stackrel{\text{def}}{=} p_{\text{rand}} \simeq 2^{-30} - 2^{-61.415}. \end{aligned}$$

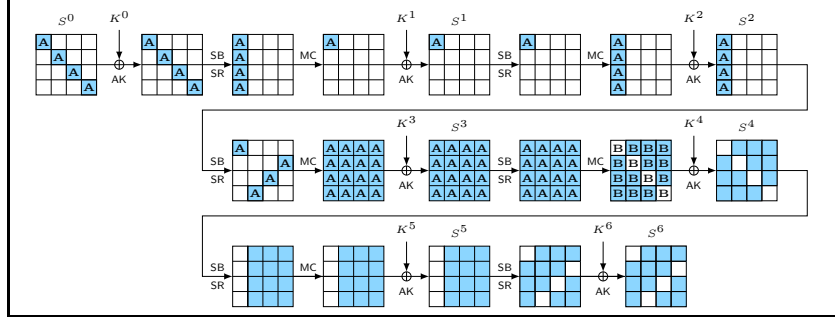


Figure 13: Six-round distinguisher. Darkened cells represent bytes with active (non-zero) difference, white cells represent bytes with zero difference.

Thus, we obtain from simple calculation that

$$p_{\text{AES}}^6 \simeq p_{\text{AES}}^5 \cdot \frac{4 \cdot 255}{2^{32} - 1} + p_{\text{rand}} \cdot \left(1 - \frac{4 \cdot 255}{2^{32} - 1}\right) \simeq 2^{-30} - 2^{-61.415} + 2^{-73.98856}. \quad (15)$$

This implies a difference of $|p_{\text{rand}} - p_{\text{AES}}^6| \simeq 2^{-73.989}$. This difference is very small, but may still be high enough to distinguish between the distributions. The setting is illustrated in Figure 13.

Equation (4) yields, again for a success probability of $P_S \geq 0.95$, that one would need approximately $N \geq 2^{120.43}$ pairs, and $N \geq 2^{121.43}$ pairs for $P_S \geq 0.99$. Assuming that a diagonal structure of 2^{32} texts contains $\binom{2^{32}}{2}$ pairs, this amount of pairs can be obtained by querying $2^{57.43}$ structures or $2^{89.43}$ chosen plaintexts.

Steps. The steps are as follows:

1. Initialize a collision counter.
2. For $i = 1..2^s$, collect a structure \mathcal{S} of 2^{32} texts that iterate over all values in $\mathcal{D}_{\{0\}}$ and leave the remaining bytes constant. Query the plaintexts of a structure and ask for their corresponding ciphertexts after six rounds. Invert the final MixColumns and ShiftRows operations to get the states $S^{6,\text{SB}}$ and store them in a list \mathcal{Q} .
3. For each structure and each state $S^{6,\text{SB}}$ therein:
 - 3.1 Initialize four lists \mathcal{L}_i , for $i = 0, 1, 2, 3$ of 2^{32} elements.
 - 3.2 For each column $i \in \{0, 1, 2, 3\}$ of $S^{6,\text{SB}}$, interpret the i -th column as 32-bit integer \mathbf{c} and append the text to $\mathcal{L}_i[\mathbf{c}]$ at the index corresponding to the column value \mathbf{c} , e.g., $\mathbf{c} = S^{6,\text{SB}}[0, 1, 2, 3]$ for Column 0.
4. For each list \mathcal{L}_i :
 - 4.1 Look for collisions, e.g., multiple values at the same index.
 - 4.2 For each collision, look up those pairs in other lists \mathcal{L}_j , for $i \neq j$, if they have already been counted. Otherwise, increment the corresponding key counter.
5. If the counter exceeds a given threshold θ , output real; otherwise, output random.

An approximation for θ can be $2^{57.43} \cdot \binom{2^{32}}{2} \cdot (p_{\text{AES}}^6 + p_{\text{rand}})/2$.

Instance	μ	σ
π		
Theory	131 067.000	362.02
Experiment	131 066.993	362.02
Small-AES		
Theory	131 067.137	362.02
Experiments	131 067.191	362.04

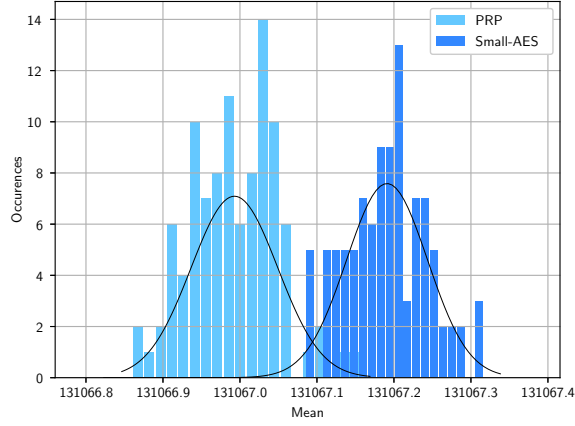


Figure 14: Distribution of the means from 100 experiments that counted the number of collisions in at least one ciphertext column per structure of 2^{16} texts of our six-round distinguisher on Small-AES and full Speck-64 as pseudorandom permutation π . The experimental values are listed in the table.

Complexity. The attack employs $2^{57.43}$ structures or $2^{89.43}$ chosen plaintexts for a success probability of about 0.95. The memory complexity is given by storing 2^{32} states in \mathcal{Q} and four lists \mathcal{L}_i of $4 \cdot 2^{32}$ columns at a time. The time complexity consists of

- $2^{89.43}$ encryptions and $1/6$ partial decryptions or $2^{89.7}$ encryption equivalents.
- $2^{89.43} \cdot 4 \simeq 2^{91.43}$ memory accesses for inserting the texts. Sorting requires approximately $2^{57.43} \cdot 4 \cdot 32 \cdot 2^{32} \simeq 2^{96.43}$ memory accesses.
- If a collision occurs when inserting for a column, at most $2 \cdot 2^{57.43} \cdot \binom{2^{32}}{2} \cdot 2^{-32} \cdot 3 \simeq 2^{91.02}$ additional memory accesses are needed to look up if other columns collide in the other lists.
- Approximately $2 \cdot 2^{57.43} \cdot \binom{2^{32}}{2} \cdot 2^{-32} \simeq 2^{89.43}$ memory accesses to increment the counter.

So, the computational complexity is given by approximately $2^{89.7}$ encryptions and $2^{91.43} + 2^{96.43} + 2^{91.02} + 2^{89.43} \simeq 2^{96.52}$ memory accesses.

Small Variant. For the small-scale variant of AES with four-bit S-boxes, we obtain from the adapted Equations (14) and (15) that

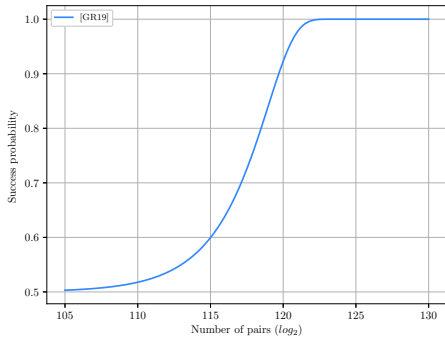
$$p_{\text{Small-AES}}^1 = 4 \cdot \frac{15}{15^4} \cdot \frac{15^4}{2^{16} - 1} = \frac{4 \cdot 15}{2^{16} - 1} \text{ and}$$

$$p_{\text{Small-AES}}^6 \simeq p_{\text{Small-AES}}^5 \cdot \frac{4 \cdot 15}{2^{16} - 1} + p_{\text{rand}} \cdot \left(1 - \frac{4 \cdot 15}{2^{16} - 1}\right) \simeq 2^{-14} - 2^{-29.415} + 2^{-33.869}.$$

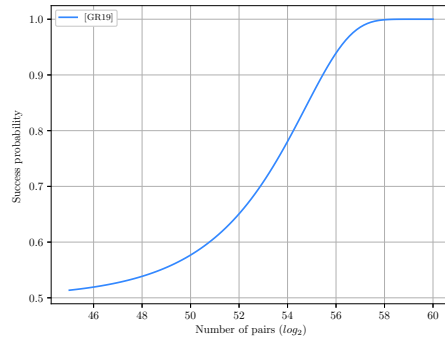
This yields a mean of $p_{\text{Small-AES}}^6 \cdot \binom{2^{16}}{2} \simeq 131\,067.137$ colliding pairs per structure. For a random permutation, the probability of a pair to collide is approximately

$$p_{\text{rand}} \simeq 2^{-14} - 2^{-29.415},$$

which gives a mean of $p_{\text{rand}} \cdot \binom{2^{16}}{2} \simeq 131\,067.000$ colliding pairs per structure. So, a similar distinguisher on the small-scale version of the AES would need $2^{56.18}$ experiments, which



(a) AES.



(b) Small-AES.

Figure 15: Success probability for the plausible range of the data complexity (number of pairs) for our six-round distinguisher, derived from Equation (4) from [GR19].

corresponds to $2^{25.18}$ structures of 2^{16} texts each, or $2^{41.18}$ chosen plaintexts. For both variants AES and Small-AES, we provide theoretical trade-offs of success probability and data complexity in Figure 15.

Experimental Verification. We implemented the distinguisher for Small-AES with 100 random keys. For each of them, we encrypted $37 \cdot 2^{20} \simeq 2^{25.21}$ structures of 2^{16} texts each. For Small-AES, we used five full rounds plus `SubBytes` and `AddRoundKeys` in the last round since `MixColumns` and `ShiftRows` are easily invertible. As a pseudorandom permutation π , we employed full Speck-64-96. For each primitive, the 100 experiments required about three CPU years of computations.

As random independent keys, we employed the most significant (leftmost) eight bytes from the first 100 output values from the NIST random Beacon service⁵. For each structure, we counted the number of pairs that collided in at least one column. For Small-AES, we observed a mean of $\mu = 131\,067.191$ pairs per structure, and $\mu = 131\,066.993$ for our pseudorandom permutation. Figure 14 illustrates the results of our experiments. Both illustrate that the difference in the distributions is even slightly higher in our experiments than in the theory. Most importantly, they show that both distributions can be distinguished well.

Transition-based Verification. Again, we employ the transition-based approach by Rønjom [Røn19] for a second form of verification. The results basically confirmed the biases of our six-round distinguishers as

$$\begin{aligned} p_{\text{AES}} &\simeq 0.533575144973 \dots \cdot 10^{-22} \simeq 2^{-30} + 2^{-73.98866} \\ p_{\text{Small-AES}} &\simeq 0.674785648935 \dots \cdot 10^{-10} \simeq 2^{-14} - 2^{-29.415} + 2^{-33.78678}, \end{aligned}$$

which have again slight deviations that can be due to rounding.

S-box and Index Dependencies. Again, we emphasize that our assumptions are an approximation of an ideal structures AES and Small-AES. Our six-round distinguisher does not directly employ a single cell as input and considers all four possible positions for δ -sets after the first round. So, the position-dependent considerations from our four- and five-round distinguishers do not apply here.

⁵See [38](https://beacon.nist.gov/beacon/2.0/chain/1/pulse/<i> for $1 \leq i \leq 100$.</p>
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8 Discussion and Conclusion

This work extends the well-known integral distinguisher on three-round AES to truncated-differential distinguishers. At the core, our attacks exploit a small difference in the average number of byte collisions between the sum of four permutations—which `MixColumns` approximates after three rounds—and the number of byte collisions of a truncated random permutation. The small but significant difference allowed to extend the integral attack to four rounds. By extending this approach to collisions in the four bytes of an inverse diagonal, we proposed a novel five-round distinguisher for the AES.

Compared to the five-round distinguisher by Grassi and Rechberger [GR19], our distinguishers start from a single active byte. Thus, they could benefit from the fact that one can easily prepend one round for an attack and start from a diagonal structure of 2^{32} texts. Thus, although inferior to previous key-recovery results on the AES, we could extend it to a key-recovery attack on six rounds in Section 6. Moreover, Section 7 showed that even this prepended round could still be included into a secret-key distinguisher, with considerably lower but still distinguishable bias.

Verifiability. We note that it is infeasible to verify the proposed distinguishers for the AES directly. Where possible, we tried where possible to implement small-scale variants with Cid et al.’s established version of Small-AES with four-bit cells. We employed the transition-distribution approach by Rønjom as second means of verification. We plan to make all our implementations freely available to the public domain.

Open Questions. We observe that the distinguishers ask several interesting follow-up questions:

- Which properties in the S-box and the cipher structure lead to the position-dependent diverse distribution of means in our four- and five-round distinguishers?
- Can we predict a-priori which input-output indices that yield particular strong distinguishers for the three-round trail in our four-round distinguisher, given the S-box? Can we predict the behavior for the S-box of the AES?
- Which property leads to the multiple-of-32 collision counts for Small-AES?
- Is it the same or a different property that leads to higher variances in our five-round distinguisher for Small-AES?
- Is it the same or a different property that leads to the multiple-of-four property when concerning a full diagonal space to a single cell through four full rounds of Small-AES?
- Which property leads to the significant deviation in the variance of Small-AES with the PRIDE S-box for our four-round distinguisher?

We think that the simple three-round truncated-differential trail that ignited this work seems particularly well-suited to further investigate those and related questions.

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A Five-round Multiple-of-Eight Property

Consider the five-round distinguisher from Setting (2), which takes pairs of plaintexts from δ -sets of a diagonal. This section shows briefly that this setting preserves the multiple-of-eight property of the full-diagonal distinguisher.

We define a cell-mixing function $\rho : \mathbb{F}_2^m \times \mathbb{F}_2^b \times \mathbb{F}_2^b \rightarrow \mathbb{F}_2^m$ similar to that in [RBH17]. Moreover, we define $\delta : \mathbb{F}_2 \times \mathbb{F}_2^b \times \mathbb{F}_2^b \rightarrow \mathbb{F}_2^b$. Let $X = (X[0], \dots, X[m-1]) \in \mathbb{F}_2^m$ be a vector of m bits, $P_i, P_j \in \mathbb{F}_2^b$ be two arbitrary elements. Then, the mixture $P = \rho(X, P_i, P_j)$ according to X is defined as

$$\delta(X[k], P^i[k], P^j[k]) \stackrel{\text{def}}{=} \begin{cases} P^i[k] & \text{if } X = 0, \\ P^j[k] & \text{otherwise,} \end{cases}$$

$$\rho(X, P^i, P^j) \stackrel{\text{def}}{=} (\delta(X[0], P^i[0], P^j[0]), \dots, \delta(X[m-1], P^i[m-1], P^j[m-1])).$$

We will slightly abuse the notation, by interpreting AES states as inputs as elements of \mathbb{F}_2^{16} . Moreover, we will consider also non-full states but also columns or diagonals.

Theorem 4. Let $\mathcal{I}, \mathcal{J} \subset \{0, 1, 2, 3\}$ and assume $|\mathcal{I}| = |\mathcal{J}| = 1$. Fix two cell indices $k \in \{0, \dots, 15\}$. Given a diagonal space $\mathcal{D}_{\mathcal{I}}$ and constants $A \in \mathbb{F}_2^{4 \times 4}$. Let \mathcal{P} denote the set of all plaintexts in a single-diagonal coset $P^i \in \mathcal{D}_{\mathcal{I}} \oplus A$, for $0 \leq i < 2^{4b}$. Let \mathcal{C} denote the set of their corresponding ciphertexts after almost five rounds without the final MixColumns operation containing $C^i = C^i = \text{MC}^{-1}(\mathbb{R}^4(P^i))$, for $0 \leq i < 2^{4b}$. Then,

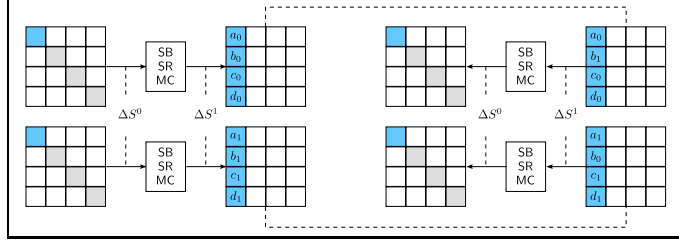


Figure 16: Left: A pair that maps a single-byte difference to a four-byte difference ΔS^1 through the first round. If this pair produces a desired difference after almost five rounds, then, there exist seven further mixture pairs (e.g., one is depicted **right**) of the variables after the first round that have the same difference ΔS^1 . Thus, they must have had the same difference before MixColumns, and therefore must also stem from δ -sets of the first diagonal.

the number of distinct pairs (C^i, C^j) for $0 \leq i < j < 2^{4b}$ such that $P^i \oplus P^j$ differ only in a single byte and $(C^i \oplus C^j)[k] = 0$ is a multiple of eight, i.e., $\exists N' \in \mathbb{N}$ such that

$$\mathcal{C} \stackrel{\text{def}}{=} |\{0 \leq i < j < 2^{4b} : P^i \oplus P^j \in \mathcal{D}_{\mathcal{I}} \cap \mathcal{C}_{\mathcal{J}} \wedge (C^i \oplus C^j)[k] = 0\}| = 8 \cdot N'.$$

Proof. Let P^i and P^j denote two plaintexts from a δ -set, i.e., that differ only in a single cell. W.l.o.g., we assume that they differ only in their first cell, $P^i[0] \neq P^j[0]$. Let S^i and S^j denote their corresponding states after the first round, respectively. P^i and P^j are encrypted to S^i and S^j through one round, which differ in four variables due to the branch number of MixColumns. We let $(a_0, b_0, c_0, d_0) \in \mathbb{F}_{2^b}^4$ denote the variables of the first column of S^i and $(a_1, b_1, c_1, d_1) \in \mathbb{F}_{2^b}^4$ denote those of S^j , as before. For this purpose, we also denote $S^i = S^{0000}$ and $S^j = S^{1111}$.

In general, let $I = (i_0, i_1, i_2, i_3) \in \mathbb{F}_2^4$. We define $S^I = \rho(I, S^i, S^j)$. For example, (S^{0001}, S^{1110}) is a mixture of (S^{0000}, S^{1111}) , which also differs only in the same column, in the sense that the leftmost column of S^{0001} consists of the values (a_0, b_0, c_0, d_1) , and that of S^{1110} of (a_1, b_1, c_1, d_0) . All other columns are equal in them and equal with the corresponding columns of S^{0000} and S^{1111} . We extend the definitions in the natural manner to

$$(S^{0010}, S^{1101}), (S^{0011}, S^{1100}), (S^{0100}, S^{1011}), \\ (S^{0101}, S^{1010}), (S^{0110}, S^{1001}), (S^{0111}, S^{1000}).$$

Note that for each of those eight mixtures, the difference between their pair after the first round is the same. Since $\text{SR}^{-1} \circ \text{MC}^{-1} \circ \text{AK}[K^1]$ is linear, this means that their difference $S_{\text{SB}}^{i_0 i_1 i_2 i_3} \oplus S_{\text{SB}}^{i'_0 i'_1 i'_2 i'_3}$ is the same for each pair. Since $S_{\text{SB}}^{0000} \oplus S_{\text{SB}}^{1111}$ is active in a single (the first) cell, so is each pair of all other seven mixtures. Hence, they are all guaranteed to be in δ -sets. Since $\text{AK}[K^0] \circ \text{SB}^{-1}$ is a permutation, they are guaranteed to be sixteen distinct texts. Though, note that the values S^1 differ in less than four bytes between each text of each pair. Thus, each of the eight pairs in a mixture is produced by a **different** δ -set.

It remains to show that if one pair of an eight-tuple of mixtures has an inactive diagonal after three rounds, the other seven mixture pairs also have. We denote by C^i the encryption of P^i after four rounds of AES without the final MixColumns operation. Assume, C_{0000} and C_{1111} lie in the coset of an inverse-diagonal space $\mathcal{ID}_{\mathcal{I}}$, for $\mathcal{I} \subset \{0, 1, 2, 3\}$ and $|\mathcal{I}| < 4$. Then, the differences of the seven further pairs

$$(C_{0001} \oplus C_{1110}), (C_{0010} \oplus C_{1101}), (C_{0011} \oplus C_{1100}), (C_{0100} \oplus C_{1011}), \\ (C_{0101} \oplus C_{1010}), (C_{0110} \oplus C_{1001}), (C_{0111} \oplus C_{1011})$$

lie also in $\mathcal{ID}_{\mathcal{I}}$. Though, this is exactly the well-known multiple-of-eight property that has already been shown by Grassi et al. in [GRR17]. Therefore, the multiple-of-eight property of our distinguisher in Setting (2) follows. \square

B Four-round Multiple-of-32 Property for Small-AES

Consider the four-round distinguisher in Setting (2), which takes pairs of plaintexts from δ -sets of a diagonal. This section outlines briefly that the distinguisher in this setting possesses a multiple-of-32 property.

Note the peaked plots that represent the experimental distribution in Setting (2) in Figure 18. Those peaks (can only) indicate that the number of collisions are a multiple of 32, e.g., the first ten random keys and δ -sets for the Small-AES yield the sequence of collisions: (123 456, 123 040, 120 512, 120 736, 123 744, 121 504, 122 368, 120 896, 123 136, 123 808). This phenomenon repeated for the other S-boxes as well, though, not for the AES with eight-bit S-boxes. 32 is also the largest power of two, e.g., the number of collisions are no multiple of 64 in general.

Clearly, the number of collisions has to be a multiple of eight since our four-round distinguisher can be seen as a subset of the five-round distinguisher that we studied in Section A. Our five-round distinguisher itself inherited that property from the well-known multiple-of-eight distinguisher by Grassi and Rechberger [GRR17]. Given that the other S-boxes yield similar results, we conjecture that there exists an additional multiple-of-four property for the AES when operating with four-bit S-boxes. Thus, we state the following claim, but cannot prove it at the moment.

Claim. Let $\mathcal{I}, \mathcal{J} \subset \{0, 1, 2, 3\}$ and assume $|\mathcal{I}| = |\mathcal{J}| = 1$. Fix a cell index $k \in \{0, \dots, 15\}$. Given a coset of $\mathcal{D}_{\mathcal{I}}$ containing all 2^{4b} plaintexts of a diagonal and their corresponding ciphertexts after four rounds (P^i, C^i) with $i \in \{0, \dots, 2^{4b} - 1\}$ and $C^i = \text{MC}^{-1}(\text{R}^4(P^i))$, that is after four rounds of AES with the final MixColumns operation omitted. Then, the number of distinct pairs (C^i, C^j) for $i \neq j$ and $0 \leq i < j < 2^{4b}$ such that $P^i \oplus P^j$ differ only in a single byte and $(C^i \oplus C^j)[k] = 0$ is a multiple of 32, i.e., $\exists N' \in \mathbb{N}$ such that

$$C \stackrel{\text{def}}{=} |\{0 \leq i < j < 2^{4b} : P^i \oplus P^j \in \mathcal{D}_{\mathcal{I}} \cap \mathcal{C}_{\mathcal{J}} \wedge (C^i \oplus C^j)[k] = 0\}| = 32 \cdot N'.$$

Four-round Distinguisher. It is straight-forward to formulate a distinguisher from the multiple-of-32 property with the following steps:

1. Initialize 16×16 collision counters $X_{i,j}$ for each byte and each value.
2. Collect a structure \mathcal{S} of 2^{16} texts that iterate over all δ -sets of a diagonal, e.g., using texts from $\mathcal{D}_{\{0\}} \oplus A$ for arbitrary $A \in \mathbb{F}_{2^4}^{4 \times 4}$, and leave the remaining bytes constants. Query the plaintexts of a structure and ask for their ciphertexts.
3. For each ciphertext, invert the final MixColumns operations to get the states $S^{4,SR}$. Increment its value j for Byte i , for $i \in \{0, \dots, 15\}$, in $X_{i,j}$.
4. For each byte $i \in \{0, \dots, 15\}$, test if $(\sum_{j=0}^{15} \binom{X_{i,j}}{2}) \bmod 32 = 0$. If yes, return “real”; return “random” otherwise.

Complexity. The distinguisher requires 2^{16} chosen plaintexts, 2^{16} encryptions and $2^{16} \cdot 16 = 2^{20}$ memory accesses. The success probability is one for the real construction; the probability is $(2^{-5})^{16} = 2^{-80}$ to wrongly identify a random permutation as the real four-round Small-AES. The memory is negligible, 256 16-bit values or 2^6 states of memory plus some constant overhead suffices.

Table 7: Means, standard deviations, and maximal k for multiple-of- k for the number of collisions in at least one anti-diagonal after four rounds of 100 000 experiments with random key each of 2^{16} texts from a full single-diagonal space. k is the maximum factor.

S-box	MC Matrix	μ	σ	k
Small-AES	circ(2, 3, 1, 1)	122 921.676	1 329.731	2^5
Small-AES	circ(1, 3, 2, 2)	122 946.734	1 347.866	2^5
Small-AES	circ(0, 1, 1, 1)	123 980.022	47 151.760	2^{14}
PRESENT	circ(2, 3, 1, 1)	122 908.033	1 209.539	2^5
PRESENT	circ(1, 3, 2, 2)	122 911.852	1 226.596	2^5
PRESENT	circ(0, 1, 1, 1)	124 763.341	47 429.148	2^{14}

Alternative MixColumns Matrices. We further tested whether the multiple-of- k property depends on the employed MixColumns matrix. We denote circulant matrices by the elements of their topmost row $\text{circ}(a_1, \dots, a_m)$, e.g., the MixColumns matrix can be described as $\text{circ}(2, 3, 1, 1)$. We employed variants of Small-AES with alternative MixColumns matrices and for comparison with different S-box:

- Small-AES with its original S-box and the circulant matrix generated by $\text{circ}(1, 3, 2, 2)$. This is also an MDS matrix.
- Small-AES with its original S-box and the matrix generated by $\text{circ}(0, 1, 1, 1)$ from Midori, which is not MDS.
- Small-AES with the S-box of PRESENT and the circulant matrix generated by $\text{circ}(1, 3, 2, 2)$.
- Small-AES with the S-box of PRESENT and the matrix $\text{circ}(0, 1, 1, 1)$.

All other components remained unchanged. Table 7 visualizes our results and compares the result with the original matrix. It illustrates that the multiple-of- k property depends strongly on the employed MixColumns matrix; the variants with the Midori matrix yielded even a multiple-of-16 384 property. The variance differs greatly as a direct consequence.

C Expectation of The Five-round Distinguisher

This section shows Theorem 2 in Setting (2). It is rooted on the existing rigorous study by Grassi and Rechberger [GR18] of the number of solutions for the transition of pairs in a mixed space $\mathcal{M}_{\{k\}}$ to a diagonal space $\mathcal{D}_{\mathcal{J}}$ with $\mathcal{J} \subset \{0, 1, 2, 3\}$ and $|\mathcal{J}| = 3$.

Notation. The subspace notation does not always express easily that certain cells are zero or not. Therefore, we can adapt the notation by [RBH17]. Though, we will use the active-byte pattern, whereas [RBH17] employ the zero-byte pattern convention. We decide for the former only since it is more intuitive that the hamming weight represents the number of active cells from our point of view.

We use bit vectors $\mathbf{u} = (u_0, u_1, u_2, u_3)$ and $\mathbf{v} = (v_0, v_1, v_2, v_3)$ with $\mathbf{u}, \mathbf{v} \in \mathbb{F}_2^4$ to represent the active-byte patterns of a column or diagonal. For a column or diagonal vector $\mathbf{x} = (x_0, x_1, x_2, x_3) \in \mathbb{F}_{2^8}^4$, we further define the function $\rho : \mathbb{F}_{2^8}^4 \rightarrow \mathbb{F}_2^4$ to map \mathbf{x} to its active-byte pattern: $\mathbf{u} = (u_0, u_1, u_2, u_3) = \rho((x_0, x_1, x_2, x_3))$, where $u_i = 1$ iff x_i is active and 0 otherwise. For a state $S \in \mathbb{F}_{2^8}^{4 \times 4}$, we define $\nu : \mathbb{F}_{2^8}^{4 \times 4} \rightarrow \mathbb{F}_2^{4 \times 4}$ to produce its active-column pattern:

$$\nu(S) \stackrel{\text{def}}{=} (\rho(S_0), \rho(S_1), \rho(S_2), \rho(S_3)).$$

Thus, we search for all pairs $\mathbf{u} \xrightarrow{R} \mathbf{v}$ with $\text{wt}(\mathbf{u}) > 1$ and $\text{wt}(\mathbf{v}) = 4$. Below, we briefly restate Theorem 2 for readability.

Theorem 2. Let $\mathcal{X} = \{X^i \in \mathbb{F}_{2^8}^{4 \times 4}\}$, for $0 \leq i < 2^{32}$, denote a set of all texts in a coset $\mathcal{M}_{\{k\}} \oplus A$ for some $A \in \mathbb{F}_{2^8}^{4 \times 4}$, $A \in \mathcal{X}^\perp$, where all columns are active. Let $\mathcal{Y} = \{Y^i \in \mathbb{F}_{2^8}^{4 \times 4} : Y^i = R(X^i)\}$, $0 \leq i < 2^{32}$, denote the set of corresponding outputs after one round of AES. Let Y be a random variable for the number of different unordered ciphertext pairs (Y^i, Y^j) s. t. $Y^i, Y^j \in \mathcal{D}_{\mathcal{J}} \oplus B$, for some $B \in \mathbb{F}_{2^8}^{4 \times 4}$, $\mathcal{J} \subset \{0, 1, 2, 3\}$, and $|\mathcal{J}| = 3$. Then, under the assumption of the uniform distribution of non-trivial solutions of differential transitions through the S-box,

$$\Pr \left[\mathcal{M}_{\{k\}} \oplus A \xrightarrow{R} \mathcal{D}_{\mathcal{J}} \oplus B \right] \simeq 2^{-30} + 2^{-51.985}.$$

The analysis of Theorem 2 will employ two Lemmas. Lemma 1 will reconsider the number of solutions, Lemma 2 the fraction of solutions that start from a single active byte.

Lemma 1. Under the assumption of the uniform distribution of non-trivial solutions of differential transitions through the S-box, the number of solutions from a mixed space with no equal generating variables is approximately $2\,114\,125\,822.5 \simeq 2^{30.977}$.

Proof. The proof of Lemma 1 is given from the third case of the proof of Theorem 2 by Grassi and Rechberger [GR19, Sect. 5.3.3]. We restate it only for the sake of completeness. In the following, we consider $\mathcal{J} = \{1, 2, 3\}$ as the set of active diagonals after the round, or $\ell = 0$ as the index of the zero-difference diagonal after the round. the cases for other values of ℓ are analogous. We denote $X^1, X^2 \in \mathcal{M}_{\{k\}} \oplus A$. Then, there exist $x^i, y^i, z^i, w^i \in \mathbb{F}_{2^8}$, for $1 \leq i \leq 2$, s. t.

$$X^i = \begin{bmatrix} 2x^i & w^i & z^i & 3y^i \\ x^i & w^i & 3z^i & 2y^i \\ x^i & 3w^i & 2z^i & y^i \\ 3x^i & 2w^i & z^i & y^i \end{bmatrix} \oplus A.$$

In the context here, all variables are distinct, i.e. $x^1 \neq x^2, y^1 \neq y^2, z^1 \neq z^2$, and $w^1 \neq w^2$. Let $\Delta Y = Y^1 \oplus Y^2$. We consider the following four equations:

$$\Delta Y[0] = \bigoplus_{i=1}^2 2S(2x^i \oplus A[0]) \oplus 3S(w^i \oplus A[5]) \oplus S(2z^i \oplus A[10]) \oplus S(y^i \oplus A[15]) = 0 \quad (16)$$

$$\Delta Y[5] = \bigoplus_{i=1}^2 S(w^i \oplus A[4]) \oplus 2S(3z^i \oplus A[9]) \oplus 3S(y^i \oplus A[14]) \oplus S(3x^i \oplus A[3]) = 0$$

$$\Delta Y[10] = \bigoplus_{i=1}^2 S(z^i \oplus A[8]) \oplus S(2y^i \oplus A[13]) \oplus 2S(x^i \oplus A[2]) \oplus 3S(2w^i \oplus A[7]) = 0$$

$$\Delta Y[15] = \bigoplus_{i=1}^2 3S(3y^i \oplus A[12]) \oplus S(x^i \oplus A[1]) \oplus S(3w^i \oplus A[6]) \oplus 2S(z^i \oplus A[11]) = 0.$$

All four equations possess a form of

$$\bigoplus_{i=1}^2 A \cdot S(B \cdot x^i \oplus b) \oplus C \cdot S(D \cdot w^i \oplus d) \oplus E \cdot S(F \cdot z^i \oplus f) \oplus G \cdot S(H \cdot y^i \oplus h) = 0.$$

We rewrite $\hat{x} = B \cdot x^i \oplus b$, $\hat{w} = D \cdot w^i \oplus d$, $\hat{z} = F \cdot z^i \oplus f$, and $\hat{y} = H \cdot y^i \oplus h$. We denote the input differences between the variable pairs as $\Delta_I^x = x^1 \oplus x^2$, and $\Delta_I^y = y^1 \oplus y^2$, $\Delta_I^z = z^1 \oplus z^2$, and $\Delta_I^w = w^1 \oplus w^2$. We can reformulate the equation systems to

$$\begin{aligned} S(\hat{x}) \oplus S(\hat{x} \oplus \Delta_I^x) &= \Delta_O^x \\ S(\hat{w}) \oplus S(\hat{w} \oplus \Delta_I^w) &= \Delta_O^w \\ S(\hat{y}) \oplus S(\hat{y} \oplus \Delta_I^y) &= \Delta_O^y \\ S(\hat{z}) \oplus S(\hat{z} \oplus \Delta_I^z) &= \Delta_O^z. \end{aligned}$$

From $\Delta_I^x, \Delta_I^w, \Delta_I^y, \Delta_I^z \neq 0$, it follows that $\Delta_O^x, \Delta_O^w, \Delta_O^y, \Delta_O^z \neq 0$. Moreover, one variable is defined from the equation system of the others as

$$\Delta_O^z = G^{-1} \cdot (A \cdot \Delta_O^x \oplus C \cdot \Delta_O^w \oplus E \cdot \Delta_O^y).$$

Since they cannot be zero, there exist in total 255^3 solutions for $\Delta_O^x, \Delta_O^w, \Delta_O^y$. Though, having fixed Δ_O^x , $255 \cdot 254$ pairs Δ_O^w, Δ_O^y solutions would lead to $\Delta_O^z = 0$. So, there are $255^3 - 255 \cdot 254$ possible tuples $(\Delta_O^x, \Delta_O^w, \Delta_O^y, \Delta_O^z)$. The next step is to obtain the number of different non-trivial solutions $[(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)]$ to the four equations from Equation (16). For $\Delta_I, \Delta_O \in \mathbb{F}_2^8$, we denote the number of solutions by

$$\delta_S(\Delta_I, \Delta_O) \stackrel{\text{def}}{=} |\{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \Delta_I) = \Delta_O\}|.$$

Independently of the S-box details, the expected number of solutions x for arbitrary non-zero input-output differences (Δ_I, Δ_O) is simply the mean over all non-zero columns or rows in the DDT:

$$\mathbb{E}[\delta_S(\Delta_I, \Delta_O)] = \frac{256}{255}.$$

Given a fixed Δ_I , the average number of solutions for one equation is $256/255$. There exist 255 values of Δ_I , and four pairs $(\Delta_I^x, \Delta_O^x), \dots, (\Delta_I^z, \Delta_O^z)$, which yields

$$\left(255 \cdot \frac{256}{255}\right)^4 = 2^{32}$$

solutions of values $(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)$ for fixed differences $(\Delta_I^x, \Delta_I^y, \Delta_I^z, \Delta_I^w)$ on average. To prevent double counting solutions $[(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)]$ and $[(x^2, y^2, z^2, w^2), (x^1, y^1, z^1, w^1)]$, we employ a partial order, e.g., require that $w^1 < w^2$. The total number of different solutions $(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)$, for the order $w^1 < w^2$, to the four Equations 16 is then given by

$$\frac{1}{2} \cdot (255^3 - 255 \cdot 254) \cdot \left(255 \cdot \frac{256}{255}\right)^4 = (255^3 - 255 \cdot 254) \cdot 2^{31}.$$

So, we have the number of solutions for one (arbitrary equation) of the four equations in Equation 16. Next, we have to determine the probability that the solutions for the remaining three equations match the fixed equation. Fix $[(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)]$ for the first equation. Consider a second equation $[(\hat{x}^1, \hat{y}^1, \hat{z}^1, \hat{w}^1), (\hat{x}^2, \hat{y}^2, \hat{z}^2, \hat{w}^2)]$. The probability that $(x^1, y^1, z^1, w^1) = (\hat{x}^1, \hat{y}^1, \hat{z}^1, \hat{w}^1)$ is

$$\left(\frac{1}{256}\right)^4.$$

Having fixed (x^1, y^1, z^1, w^1) , there remain only 255 possible values for x^2 and \hat{x}^2 since they must be distinct. So, the probability that $(x^2, y^2, z^2, w^2) = (\hat{x}^2, \hat{y}^2, \hat{z}^2, \hat{w}^2)$ is $1/255^4$.

The same can apply if $(x^1, y^1, z^1, w^1) = (\hat{x}^2, \hat{y}^2, \hat{z}^2, \hat{w}^2)$ and $(x^2, y^2, z^2, w^2) = (\hat{x}^1, \hat{y}^1, \hat{z}^1, \hat{w}^1)$. So, there are two ways to combine each pair of equations. Thus, the probability that two equations are equal is

$$2 \cdot \left(\frac{1}{256}\right)^4 \cdot \left(\frac{1}{255}\right)^4. \quad (17)$$

We obtain that

$$\left((255^3 - 255 \cdot 254) \cdot 2^{31}\right)^4 \cdot \left(2 \cdot \left(\frac{1}{256}\right)^4 \cdot \left(\frac{1}{255}\right)^4\right)^3 \simeq 2\,114\,125\,822.5 \simeq 2^{30.977}, \quad (18)$$

which yields the claim of Lemma 1. \square

Equation (18) uses all pairs that come from a mixed space $\mathcal{M}_{\{k\}}$ with four different variables. Clearly, such a mixed space must have resulted from a diagonal space $\mathcal{D}_{\{k\}}$ two rounds earlier. Here, we do not consider all pairs from the diagonal space $\mathcal{D}_{\{k\}}$. Among the $\binom{2^{32}}{2}$ pairs from a full diagonal space, we restrict our focus on δ -sets. Assume, we consider all δ -sets from a diagonal space. This yields $2^{24} \cdot 4$ δ -sets, each of which contains $\binom{2^8}{2}$ pairs. We have to determine which fraction of solutions is covered by the δ -sets. For this purpose, we have to count the number of pairs that translate from a single active diagonal to a mixed space with two restrictions: (1) their plaintext pairs have come from a diagonal space and (2) their mixed space after two rounds has four active columns, which implies that after the first round, all bytes in Column k are active.

Lemma 2. Let $\mathcal{X} = \{X^i \in \mathbb{F}_2^{4 \times 4}\}$, for $0 \leq i < 2^{32}$, denote a set of all texts in a coset $\mathcal{D}_{\{k\}} \oplus A$ for some $A \in \mathbb{F}_2^{4 \times 4}$. Then, the fraction of pairs X^i, X^j whose difference $\Delta X^{i,j}$ lies in $\mathcal{C}_{\{c\}} \cap \mathcal{D}_{\{k\}}$, for any $c \in \{0, 1, 2, 3\}$, among all pairs in \mathcal{X} , is $4/255^3$.

Proof of Lemma 2. Table 8 lists the number of differences that follow the transitions from $\mathbf{u} \xrightarrow{\mathbb{R}} \mathbf{v}$. The rightmost column of that table lists all input differences that lead to four active bytes in $\mathcal{C}_{\{k\}}$ after the first round. In sum, there are 255^4 such differences. In total, there are $2^{32} - 1$ differences. So, from $2^{32} \cdot (2^{32} - 1)/2$ pairs, we can multiply each entry in Table 8 by $2^{32}/2 = 2^{31}$ to obtain the number of pairs with such a difference in a diagonal. We obtain that a diagonal contains $4 \cdot 2^{24} \cdot \binom{2^8}{2}$ pairs in δ -sets that differ in a single plaintext byte. This is a fraction of

$$\frac{\binom{4}{1} \binom{4}{4} \cdot 255 \cdot 2^{31}}{255^4 \cdot 2^{31}} = \frac{4}{255^3}$$

pairs, which yields the claim in Lemma 2. \square

It remains to combine the results of Lemma 1 and 2 to show Theorem 2.

Proof of Theorem 2. A diagonal space has 2 114 125 822.5 solutions, on average, when we consider pairs generated from a mixed space with four pair-wise different variables $x^1 \neq x^2$, $y^1 \neq y^2$, $w^1 \neq w^2$, $z^1 \neq z^2$ after two rounds. Among all the pairs with four pair-wise different variables, our expectation distinguisher considers a fraction of $4/255^3$ of the pairs, namely only those that were mapped from a diagonal space of one active plaintext byte among to a column of four active bytes after the first round. Thus, we have on average

$$2\,114\,125\,822.5 \cdot \frac{4}{255^3} \simeq 510.000123$$

Table 8: Number of differences that translate from a diagonal vector \mathbf{x} with active-byte pattern \mathbf{u} to a column vector \mathbf{y} with active-byte pattern \mathbf{v} through one round of the AES, depending on the weight of \mathbf{u} and \mathbf{v} .

wt(\mathbf{u})	wt(\mathbf{v})			
	1	2	3	4
1	–	–	–	$\binom{4}{1} \binom{4}{4} 255$
2	–	–	$\binom{4}{2} \binom{4}{3} 255$	$\binom{4}{2} \binom{4}{4} 64\,005$
3	–	$\binom{4}{3} \binom{4}{2} 255$	$\binom{4}{3} \binom{4}{3} 64\,005$	$\binom{4}{3} \binom{4}{4} 16\,323\,825$
4	$\binom{4}{4} \binom{4}{1} 255$	$\binom{4}{4} \binom{4}{2} 64\,005$	$\binom{4}{4} \binom{4}{3} 16\,323\,825$	$\binom{4}{4} \binom{4}{4} 4\,162\,570\,275$
Σ	$\binom{4}{1} \cdot 255$	$\binom{4}{2} \cdot 255^2$	$\binom{4}{3} \cdot 255^3$	$\binom{4}{4} \cdot 255^4$

solutions if we consider all interesting pairs among a diagonal space of 2^{32} plaintexts. There exist $4 \cdot 2^{24} \cdot \binom{2^8}{2}$ such pairs in such a diagonal space. Thus, we obtain a probability of

$$\frac{2\,114\,125\,822.5 \cdot \frac{4}{255^3}}{4 \cdot 2^{24} \cdot \binom{2^8}{2}} \simeq 2^{-32} + 2^{-53.9831}.$$

Since there are four diagonals, the probability to have one inactive diagonal after three rounds is on average

$$p_{\text{AES}} \simeq 1 - (1 - (2^{-32} + 2^{-53.9831}))^4 \simeq 2^{-30} + 2^{-51.98515}.$$

This yields the claim of our five-round distinguisher. \square

D Expectation of The Four-round Distinguisher

This section shows Theorem 1 in Setting (2). We restate it below for readability. Our goal in the remainder is to obtain an approximate for the probability p for the transition of the third round, under the condition that all columns after two rounds are active.

Theorem 1. Let $\mathcal{X} = \{X^i \in \mathbb{F}_{2^8}^{4 \times 4}\}$, for $0 \leq i < 2^{32}$, denote a set of all texts in a coset $\mathcal{M}_{\{k\}} \oplus A$ for some $A \in \mathbb{F}_{2^8}^{4 \times 4}$, $A \in \mathcal{X}^\perp$, where all columns are active. Let $\mathcal{Y} = \{Y^i \in \mathbb{F}_{2^8}^{4 \times 4} : Y^i = R(X^i)\}$, $0 \leq i < 2^{32}$, denote the set of corresponding outputs after one round of AES. Let Y be a random variable for the number of different ciphertext pairs (Y^i, Y^j) s. t. $Y^i, Y^j \in \mathcal{D}_{\mathcal{J}} \cap \mathcal{C}_{\{c\}} \oplus A'$, for some $A' \in \mathbb{F}_{2^8}^{4 \times 4}$, $\mathcal{J} \subset \{0, 1, 2, 3\}$, $|\mathcal{J}| = 3$, and $c \in \{0, 1, 2, 3\}$. Then, under the assumption of the uniform distribution of non-trivial solutions of S-box differential transitions in Equation (5),

$$\Pr \left[\mathcal{M}_{\{k\}} \oplus A \xrightarrow{R} \mathcal{D}_{\mathcal{J}} \cap \mathcal{C}_{\{c\}} \oplus A' \right] \simeq 2^{-8} + 2^{-31.983}.$$

Note that the analysis of Theorem 1 is a simpler variant of the proof strategy of our five-round distinguisher. To show Theorem 1, we will employ two Lemmas. Lemma 3 will reconsider the number of solutions, Lemma 2 provides us with the fraction of solutions that start from a single active byte.

Lemma 3. Under the assumption of the uniform distribution of non-trivial solutions of S-box differential transitions, the number of solutions from a mixed space with no equal generating variables is approximately $2^{54.9774}$.

Proof. Again, we can prove Lemma 3 from adapting the third case of the proof of Theorem 2 by Grassi and Rechberger [GR19, Sect. 5.3.3]. Let $\mathcal{J} = \{1, 2, 3\}$ as the set of active diagonals after the round, and let $\ell = 0$ and $c = 0$; the cases for other values of ℓ and c are analogous. We denote $X^1, X^2 \in \mathcal{M}_{\{k\}} \oplus A$. Then, there exist $x^i, y^i, z^i, w^i \in \mathbb{F}_{2^8}$, for $1 \leq i \leq 2$, s. t.

$$X^i = \begin{bmatrix} 2x^i & w^i & z^i & 3y^i \\ x^i & w^i & 3z^i & 2y^i \\ x^i & 3w^i & 2z^i & y^i \\ 3x^i & 2w^i & z^i & y^i \end{bmatrix} \oplus A.$$

All variables are distinct, i.e. $x^1 \neq x^2, y^1 \neq y^2, z^1 \neq z^2$, and $w^1 \neq w^2$. We consider the following equation:

$$\Delta Y[0] = \bigoplus_{i=1}^2 2S(2x^i \oplus A[0]) \oplus 3S(w^i \oplus A[5]) \oplus S(2z^i \oplus A[10]) \oplus S(y^i \oplus A[15]) = 0, \quad (19)$$

which can be generalized to

$$\bigoplus_{i=1}^2 A \cdot S(\hat{x}) \oplus C \cdot S(\hat{w}) \oplus E \cdot S(\hat{z}) \oplus G \cdot S(\hat{y}) = 0.$$

We denote the input differences between the variable pairs as $\Delta_I^x = x^1 \oplus x^2$, and $\Delta_I^y = y^1 \oplus y^2$, $\Delta_I^z = z^1 \oplus z^2$, and $\Delta_I^w = w^1 \oplus w^2$. We can reformulate the equation systems to

$$\begin{aligned} S(\hat{x}) \oplus S(\hat{x} \oplus \Delta_I^x) &= \Delta_O^x \\ S(\hat{w}) \oplus S(\hat{w} \oplus \Delta_I^w) &= \Delta_O^w \\ S(\hat{y}) \oplus S(\hat{y} \oplus \Delta_I^y) &= \Delta_O^y \\ S(\hat{z}) \oplus S(\hat{z} \oplus \Delta_I^z) &= \Delta_O^z. \end{aligned}$$

From $\Delta_I^x, \Delta_I^w, \Delta_I^y, \Delta_I^z \neq 0$, it follows that $\Delta_O^x, \Delta_O^w, \Delta_O^y, \Delta_O^z \neq 0$. Moreover, one variable is defined from the equation system of the others as

$$\Delta_O^z = G^{-1} \cdot (A \cdot \Delta_O^x \oplus C \cdot \Delta_O^w \oplus E \cdot \Delta_O^y).$$

Again, there exist 255^3 solutions for $\Delta_O^x, \Delta_O^w, \Delta_O^y$. Though, having fixed Δ_O^x , $255 \cdot 254$ pairs Δ_O^w, Δ_O^y solutions would lead to $\Delta_O^z = 0$. So, there are $255^3 - 255 \cdot 254$ possible tuples $(\Delta_O^x, \Delta_O^w, \Delta_O^y, \Delta_O^z)$. Given a fixed Δ_I , the average number of solutions for one equation is $256/255$. There exist 255 values of Δ_I , and four pairs $(\Delta_I^x, \Delta_O^x), \dots, (\Delta_I^z, \Delta_O^z)$, which yields

$$\left(255 \cdot \frac{256}{255}\right)^4 = 2^{32}$$

solutions of values $(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)$ for fixed differences $(\Delta_I^x, \Delta_I^y, \Delta_I^z, \Delta_I^w)$ on average. To prevent double counting solutions $[(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)]$ and $[(x^2, y^2, z^2, w^2), (x^1, y^1, z^1, w^1)]$, we employ a partial order, e.g., require that $w^1 < w^2$. The total number of different solutions $(x^1, y^1, z^1, w^1), (x^2, y^2, z^2, w^2)$, for the order $w^1 < w^2$, to the four Equations 16 is then given by

$$\frac{1}{2} \cdot (255^3 - 255 \cdot 254) \cdot \left(255 \cdot \frac{256}{255}\right)^4 = (255^3 - 255 \cdot 254) \cdot 2^{31} \quad (20)$$

$$= 35469139157975040 \simeq 2^{54.9774}. \quad (21)$$

which yields the claim of Lemma 3. \square

Equation (20) uses all pairs that come from a mixed space $\mathcal{M}_{\{k\}}$ with four different variables. Clearly, such a mixed space must have resulted from a diagonal space $\mathcal{D}_{\{k\}}$ two rounds earlier. Here, we do not consider all pairs from the diagonal space $\mathcal{D}_{\{k\}}$. Among the $\binom{2^{32}}{2}$ pairs from a full diagonal space, we restrict our focus on δ -sets. Assume, we consider all δ -sets from a diagonal space. Then, Lemma 2 showed that the fraction of pairs X^i, X^j whose difference $\Delta X^{i,j}$ lies in $\mathcal{C}_{\{c\}} \cap \mathcal{D}_{\{k\}}$, for any $c \in \{0, 1, 2, 3\}$, among all pairs in \mathcal{X} , is $4/255^3$. This yields $2^{24} \cdot 4$ δ -sets, each of which contains $\binom{2^8}{2}$ pairs.

Proof of Theorem 1. A diagonal space has $(255^3 - 255 \cdot 254) \cdot 2^{31}$ solutions when we consider pairs generated from a mixed space with four pair-wise different variables $x^1 \neq x^2$, $y^1 \neq y^2$, $w^1 \neq w^2$, $z^1 \neq z^2$ after two rounds. Among all the pairs with four pair-wise different variables, our expectation distinguisher considers a fraction of $4/255^3$ of the pairs, namely only those that were mapped from a diagonal space of one active plaintext byte among to a column of four active bytes after the first round. Among a diagonal space of 2^{32} plaintexts, there exist $4 \cdot 2^{24} \cdot \binom{2^8}{2}$ such pairs in such a diagonal space. Thus, we obtain a probability of

$$\frac{35469139157975040 \cdot \frac{4}{255^3}}{4 \cdot 2^{24} \cdot \binom{2^8}{2}} \simeq 2^{-8} + 2^{-31.983}.$$

which yields the claim of our four-round distinguisher. \square

E Variance of The Five-round Distinguisher

E.1 Recap of the Variance from [GR19]

Grassi and Rechberger [GR19] considered all pairs in a diagonal space $\mathcal{D}_{\mathcal{I}}$ to an inverse-diagonal space after five rounds. They showed that their distribution of collisions could be modeled by the sum of three binomial distributions. Let Y be the random variable for the number of collisions after five rounds. Then, the probability of k collisions from N trials follows

$$\Pr[Y = k] = \begin{cases} 0 & \text{if } k \bmod 8 \neq 0 \\ \sum_{k_3, k_{10}, k_{17} \in \mathcal{K}_N} \left(\prod_{i \in \{3, 10, 17\}} \binom{N_i}{k_i} \cdot (p_i)^{k_i} \cdot (1 - p_i)^{N_i - k_i} \right) & \text{otherwise,} \end{cases}$$

where

$$K_N \stackrel{\text{def}}{=} \{(k_3, k_{10}, k_{17}) \in \mathbb{N}^3 \mid 0 \leq k_i \leq N_i \wedge 2^3 \cdot k_3 + 2^{10} \cdot k_{10} + 2^{17} \cdot k_{17} = k\},$$

and

$$\begin{aligned} N_3 &= 1 \cdot 2^{28} \cdot 255^4 & p_3 &= 2^{-32} + 2^{-53.983} \\ N_{10} &= 4 \cdot 2^{21} \cdot 255^3 & p_{10} &= 2^{-32} - 2^{-45.989} \\ N_{17} &= 6 \cdot 2^{14} \cdot 255^2 & p_{17} &= 2^{-32} + 2^{-37.986}. \end{aligned}$$

Note that, in contrast to p_i and N_i , the tuples $(k_3, k_{10}, \text{ and } k_{17})$ are no single values but can take potentially many combinations

Or said easier: the distribution of the number of a fixed single anti-diagonal collisions after five-round AES from a diagonal space \mathcal{Y} is described by

$$\Pr[\mathcal{Y}] = 2^3 \cdot \mathcal{X}_3 + 2^{10} \cdot \mathcal{X}_{10} + 2^{17} \cdot \mathcal{X}_{17},$$

where $\mathcal{X}_i \sim \mathcal{B}(N_i, p_i)$ are binomial distributions for $i \in \{3, 10, 17\}$. The factors 2^3 , 2^{10} , and 2^{17} represent the number of pairs in the input space that follow this distribution.

They correspond to the pairs with four, three, and two (in that order) different generating variables in the expectation proof.

[GR19] determined the variance $\text{Var}[Y]$ as

$$\text{Var}[Y] \stackrel{\text{def}}{=} \text{Var}[2^3 \cdot X_3] + \text{Var}[2^{10} \cdot X_{10}] + \text{Var}[2^{17} \cdot X_{17}] .$$

Since $\text{Var}[\alpha \cdot X] = \alpha^2 \cdot \text{Var}[X]$ holds for any constant α , it followed that

$$\begin{aligned} \text{Var}[Y] &= (2^3)^2 \cdot \text{Var}[X_3] + (2^{10})^2 \cdot \text{Var}[X_{10}] + (2^{17})^2 \cdot \text{Var}[X_{17}] \\ &= 2^6 \cdot (N_3 \cdot p_3 \cdot (1 - p_3)) + 2^{20} \cdot (N_{10} \cdot p_{10} \cdot (1 - p_{10})) + 2^{34} \cdot (N_{17} \cdot p_{17} \cdot (1 - p_{17})) . \end{aligned}$$

[GR19] employed a full diagonal space with up to four active bytes in the input. To obtain a zero-difference diagonal after three rounds (i.e., at most three active bytes per column), they needed at least two active bytes per column, i.e., at least two active bytes after the first round. Under other output restrictions, the sum would also contain another term X_{24} for the number of pairs with a single different generating variable.

E.2 Variance of Our Five-round Distinguisher

The five-round distinguisher in this work does not employ a full diagonal space, but starts from a single active byte. Therefore, all pairs after one round are guaranteed to have four active bytes, and four different generating variables. In the following, we derive the variance for both settings.

Setting (2). Here, we claim that the probability distribution of collisions of a single inverse diagonal after five rounds is given by

$$\Pr[Y = k] = \begin{cases} 0 & \text{if } k \bmod 8 \neq 0 \\ \binom{N_3}{k_3} \cdot (p_3)^{k_3} \cdot (1 - p_3)^{N_3 - k_3} & \text{otherwise.} \end{cases}$$

So, we obtain actually one of the three sub-cases from [GR19]. We can determine the variance as

$$\begin{aligned} \text{Var}[Y] &= (2^3)^2 \cdot \text{Var}[X_3] = 2^6 \cdot (N_3 \cdot p_3 \cdot (1 - p_3)) \\ &= 2^6 \cdot \frac{4 \cdot 255 \cdot 2^{31}}{8} \cdot (2^{-30} + 2^{-51.985} - 2^{-61.415}) \cdot (1 - (2^{-30} + 2^{-51.985} - 2^{-61.415})) \\ &\simeq 4080 . \end{aligned}$$

Note that the variance of a random binomial distribution is given by

$$\begin{aligned} \text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{24} \cdot \binom{2^8}{2} \cdot (2^{-30} - 2^{-61.415}) \cdot (1 - (2^{-30} - 2^{-61.415})) \\ &\simeq 509.999 \dots \simeq 510 . \end{aligned}$$

So, the variance is about eight times higher in Setting (2).

Setting (1). This setting does not have the strong multiple-of-eight relation, The distribution after N trials follows approximately a binomial distribution:

$$\Pr[Y = k] = \binom{N}{k} \cdot p^k \cdot (1 - p)^{N - k} .$$

We can determine the variance as

$$\begin{aligned}\text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{24} \cdot \binom{2^8}{2} \cdot (2^{-30} + 2^{-51.985} - 2^{-61.415}) \cdot (1 - (2^{-30} + 2^{-51.985} - 2^{-61.415})) \\ &\simeq 510.000122 \dots\end{aligned}$$

So, the variance is marginally higher than that for a random permutation.

E.3 Variance of Our Five-round Distinguisher for Small-AES

For Small-AES, we derive the variance similarly.

Setting (2). Here, we can determine the variance as

$$\begin{aligned}\text{Var}[Y] &= (2^3)^2 \cdot \text{Var}[X_3] = 2^6 \cdot (N_3 \cdot p_3 \cdot (1 - p_3)) \\ &= 2^6 \cdot \frac{4 \cdot 15 \cdot 2^{15}}{8} \cdot (2^{-14} + 2^{-23.748} - 2^{-29.415}) \cdot (1 - (2^{-14} + 2^{-23.748} - 2^{-29.415})) \\ &\simeq 961.094 \simeq 2^{9.9085}.\end{aligned}$$

Note that the variance of a random binomial distribution is given by

$$\begin{aligned}\text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{12} \cdot \binom{2^4}{2} \cdot (2^{-14} - 2^{-29.415}) \cdot (1 - (2^{-14} - 2^{-29.415})) \\ &\simeq 119.9899 \dots \simeq 2^{6.907}.\end{aligned}$$

So, the variance is again about eight times higher in the Setting (2).

Setting (1). Again, this setting lacks the multiple-of-eight relation, and the distribution of collisions after N trials follows roughly a binomial distribution:

$$\Pr[Y = k] = \binom{N}{k} \cdot p^k \cdot (1 - p)^{N-k}.$$

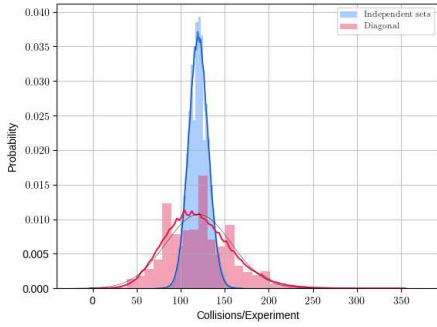
We can determine the variance as

$$\begin{aligned}\text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{12} \cdot \binom{2^4}{2} \cdot (2^{-14} + 2^{-23.748} - 2^{-29.415}) \cdot (1 - (2^{-14} + 2^{-23.748} - 2^{-29.415})) \\ &\simeq 120.1295 \dots \simeq 2^{6.908}.\end{aligned}$$

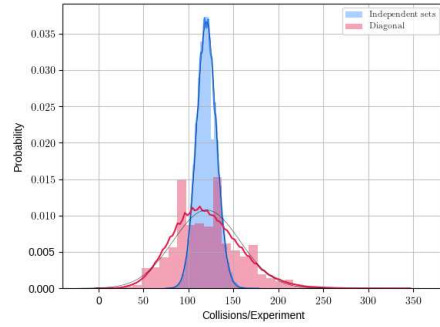
So, the variance in Setting (1) is again marginally higher than that for a random permutation.

E.4 Experimental Verification of the Variance of Our Five-round Distinguisher for Small-AES

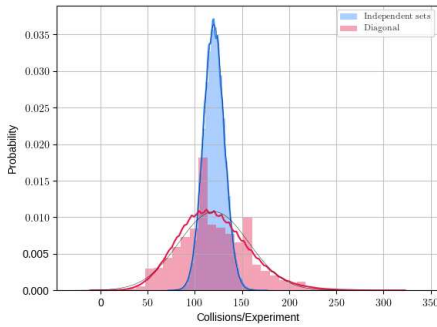
We verified the variance experimentally for Small-AES. We chose 100 000 random keys. For Setting (1), we chose 2^{14} random independent δ -sets per key, which matches the number of δ -sets in a diagonal of Setting (2). For each such experiment, we counted the number of collisions in at least one of the inverse diagonals after five rounds without the final MixColumns operation. Figure 17 visualizes the distributions for both settings.



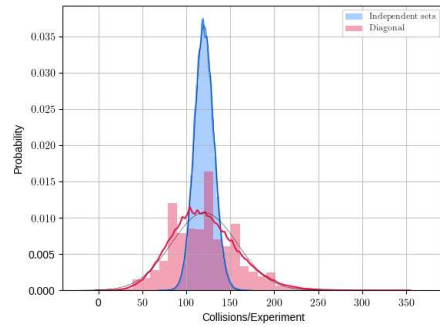
(a) Small-AES



(b) Small-AES with PRESENT S-box



(c) Small-AES with PRINCE S-box



(d) Small-AES with PRIDE S-box

Figure 17: Distribution of the number of collisions in the first cell after five rounds in both settings, with the final MixColumns operation omitted. Both distributions used 100 000 random keys and counted the collisions from 2^{14} δ -sets per key that were independent in Setting (1) or the 2^{14} δ -sets from $\mathcal{D}_{\{0\}}$. Thin lines are fitted normal distributions, thick lines the best fitted distributions.

E.5 Discussion

As a result, we conclude that there has to be another property in Small-AES beyond the multiple-of-eight property that yields the significant deviation. In Setting (2), the variance differs significantly from our expectation. One could expect that it roughly is eight times higher than the variance in Setting (1) due to the multiple-of-eight property. Though, for Small-AES, it is only roughly 11.4 times larger as outlined in Table 10, i.e., roughly 12 times larger than the variance of Setting (1). The reason for this gap is still unknown.

F Variance of The Four-round Distinguisher

The variance of the four-round distinguisher can be derived similarly as for our five-round distinguisher.

Table 9: Summary of the variances for our five-round distinguisher. π = random permutation. Exp. = experiments from 100 000 random keys with 2^{14} δ -sets each.

Construction	Setting (1)		Setting (2)	
	Theory	Experiments	Theory	Experiments
π	509.999	–	509.999	–
AES	510.000	–	4080.000	–
π	119.9899	–	119.9899	–
Small-AES	120.130	119.601	961.094	1 388.735
PRESENT S-box	120.130	118.796	961.094	1 375.978
PRINCE S-box	120.130	120.299	961.094	1 370.575
PRIDE S-box	120.130	119.996	961.094	1 391.445

F.1 Variance of Our Four-round Distinguisher for The AES

Setting (2). Here, we claim that the probability distribution of collisions of a single inverse diagonal after five rounds is given by

$$\Pr[Y = k] = \begin{cases} 0 & \text{if } k \bmod 8 \neq 0 \\ \binom{N_3}{k_3} \cdot (p_3)^{k_3} \cdot (1 - p_3)^{N_3 - k_3} & \text{otherwise.} \end{cases}$$

So, we obtain actually one of the three sub-cases from [GR19]. We can determine the variance as

$$\begin{aligned} \text{Var}[Y] &= (2^3)^2 \cdot \text{Var}[X_3] = 2^6 \cdot (N_3 \cdot p_3 \cdot (1 - p_3)) \\ &= 2^6 \cdot \frac{4 \cdot 255 \cdot 2^{31}}{8} \cdot (2^{-8} + 2^{-31.983}) \cdot (1 - (2^{-8} + 2^{-31.983})) \\ &\simeq 68\,183\,658\,496.108 \simeq 2^{35.9887}. \end{aligned}$$

Note that the variance of a random binomial distribution is given by

$$\begin{aligned} \text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{24} \cdot \binom{2^8}{2} \cdot (2^{-8}) \cdot (1 - (2^{-8})) \\ &\simeq 8\,522\,956\,800 \simeq 2^{32.9887}. \end{aligned}$$

So, the variance should be about eight times higher in Setting (2).

F.2 Variance of Our Four-round Distinguisher for Small-AES

For Small-AES, we derive the variance similarly. However, due to the multiple-of-32 property, we have a different computation. We define p_5 , X_5 and N_5 similar as p_3 , X_3 and N_3 before.

Setting (2). Here, we can determine the variance as

$$\begin{aligned} \text{Var}[Y] &= (2^5)^2 \cdot \text{Var}[X_5] = 2^{10} \cdot (N_5 \cdot p_5 \cdot (1 - p_5)) \\ &= 2^{10} \cdot \frac{4 \cdot 15 \cdot 2^{15}}{32} \cdot (2^{-4} + 2^{-15.721}) \cdot (1 - (2^{-4} + 2^{-15.721})) \\ &\simeq 3\,687\,419.195 \simeq 2^{21.8142}. \end{aligned}$$

Table 10: Summary of the variances for our four-round distinguisher. RP = random permutation. Exp. = experiments from 100 000 random keys with 2^{14} δ -sets each.

Construction	Setting (1)		Setting (2)	
	Theory	Experiments	Theory	Experiments
π	$2^{32.9887}$	–	$2^{32.9887}$	–
AES	$2^{32.9887}$	–	$2^{35.9887}$	–
π	115 200	–	115 200	–
Small-AES	115 231.850	115 080.346	3 687 419.195	1 768 183.658
PRESENT S-box	115 231.850	116 507.699	3 687 419.195	1 462 985.130
PRINCE S-box	115 231.850	115 049.655	3 687 419.195	1 886 523.688
PRIDE S-box	115 231.850	114 859.203	3 687 419.195	4 709 431.259

Note that the variance of a random binomial distribution is given by

$$\begin{aligned} \text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{12} \cdot \binom{2^4}{2} \cdot (2^{-4} - 2^{-64.093}) \cdot (1 - (2^{-4} - 2^{-64.093})) \\ &\simeq 115\,200 \simeq 2^{16.8138}. \end{aligned}$$

So, the variance should be about 32 times higher in the Setting (2).

Setting (1). Again, this setting lacks the multiple-of-32 relation, and the distribution of collisions after N trials follows roughly a binomial distribution:

$$\Pr[Y = k] = \binom{N}{k} \cdot p^k \cdot (1 - p)^{N-k}.$$

We can determine the variance as

$$\begin{aligned} \text{Var}[Y] &= N \cdot p \cdot (1 - p) \\ &= 4 \cdot 2^{12} \cdot \binom{2^4}{2} \cdot (2^{-4} + 2^{-15.721}) \cdot (1 - (2^{-4} + 2^{-15.721})) \\ &\simeq 115\,231.850 \dots \simeq 2^{16.8142}. \end{aligned}$$

So, the variance in Setting (1) is again marginally higher than that for a random permutation.

F.3 Experimental Verification of the Variance of Our Four-round Distinguisher for Small-AES

We verified the variance experimentally for Small-AES. We chose 100 000 random keys. For Setting (1), we chose 2^{14} random independent δ -sets per key, which matches the number of δ -sets in a diagonal of Setting (2). For each such experiment, we counted the number of collisions in the first byte after four rounds without the final MixColumns operation. Figure 18 visualizes the distributions for both settings.

F.4 Discussion

In Setting (2), the variance differs significantly from our expectation. First, one can expect that it roughly is 32 times higher than the variance in Setting (1) due to the multiple-of-32

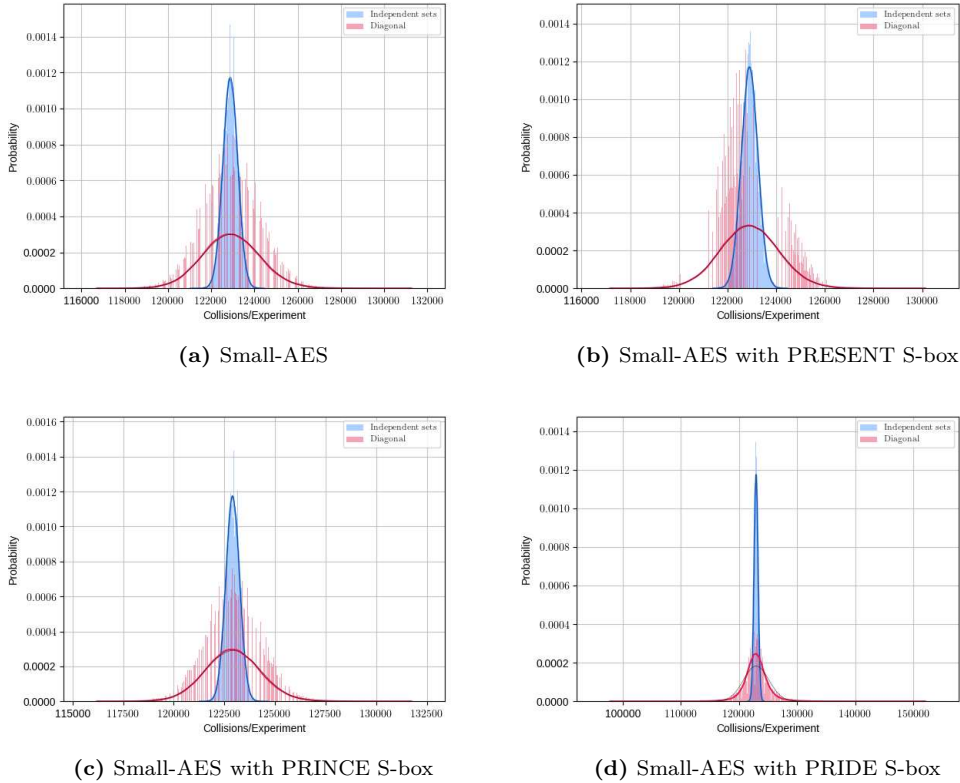


Figure 18: Distribution of the number of collisions in the first cell after four rounds in both settings, with the final MixColumns operation omitted. Both distributions used 100 000 random keys and counted the collisions from 2^{14} δ -sets per key that were independent in Setting (1) or the 2^{14} δ -sets from $\mathcal{D}_{\{0\}}$. Thin lines are fitted normal distributions, thick lines the best fitted distributions.

property. From the experiments, we observed the multiple-of-32 property in Section B. However, the variance of Small-AES in Setting (2) can then be expected to be roughly a 32 times the variance of Setting (1). Though, for Small-AES, it is only roughly 15.4 times larger as outlined in Table 10, i.e., roughly 16 times larger than the variance of Setting (1). The reason for this gap is still unknown.

For the five-round distinguisher, the tested S-boxes behaved quite similarly. For the four-round distinguisher, they differ hugely, e.g., the variance for Small-AES with the S-box of PRESENT has a factor of 12.7 instead of roughly 16, and that of PRIDE a factor of roughly 40.9. It is equally unknown which properties cause those S-boxes to possess such huge differences.

We conjecture that the multiple-of-four property behaves in a multiplicative fashion in the four-round distinguisher and in an additive fashion in our five-round distinguisher. This means, instead of 16-tuples of texts that form a multiple-of-eight property, the property builds 64-tuples texts that yield the multiple-of-32 property. So, it is included in the multiple-of-eight tuples. For the five-round distinguisher, we conjecture that the property behaves additively, i.e., there exist 16-tuples of texts that yield the multiple-of-eight property and separate 8-tuples that yield an additional multiple-of-four property that is hidden but notable in the variance. It is imaginable that the same property causes the five-round multiple-of-four distinguisher in Appendix H.

G Four-round Distinguisher of Small-AES with Alternative S-boxes

This section studies our four-round distinguisher with versions of Small-AES instantiated with different S-boxes. First, we employed six S-boxes that are a relevant subset of the experiments from [GR18, GR19]: three-real world S-boxes from PRESENT, PRINCE, and PRIDE, as well as three toy S-boxes, Toy-6, Toy-8, Toy-10 [GR19], named after their maximal differential probability. The results of those experiments are listed in Table 11 and Figure 19. Second, we employed 20 random 4-bit S-boxes. For the sake of reproducibility, we used the first four bytes from the first NIST random beacon value as seed. For illustration purpose, the results are illustrated as their fitted normal curves of their in Figure 2 and listed in Table 12. The following python snippet was used to produce the S-boxes.

```

1 from numpy import random
2 random.seed(0x7665F054) # First 4 bytes of first NIST random-beacon value
3
4 for i in range(20):
5     perm = random.permutation(16)

```

Table 11: Results of the four-round distinguisher on Small-AES with different S-boxes and 100 random keys. Speck-64-96 has been used as PRP. Difference = to the theoretical value; \mathbf{DP}_{\max} = the maximal number of pairs for a differential through the S-box; Var = variance; hom. = homogeneous; \checkmark /- = yes/no.

Instance	μ	σ	Difference to Theory				
			π	Small-AES	DU	Var	Hom.
Theory	62 914 000.	7 680.	-	-19 200.			
PRP	62 915 004.	7 820.	+1 004.	-18 196.			
Theory	62 933 200.	7 681.	+19 200.		-		
Small-AES	62 967 244.	7 030.	+53 244.	+34 044.	4 344/225		\checkmark
PRESENT	62 928 210.	8 321.	+14 210.	-4 990.	4 416/225		-
PRINCE	62 928 742.	8 387.	+14 742.	-4 458.	4 344/225		-
PRIDE	62 940 902.	7 761.	+26 902.	+7 682.	4 416/225		-
Toy-6	62 892 051.	8 151.	-21 949.	-41 148.	6 392/225		-
Toy-8	62 913 848.	7 443.	-152.	-19 352.	8 464/225		-
Toy-10	62 913 322.	7 779.	-678.	-19 878.	10 608/225		-

H Multiple-of-four Distinguisher and Attack on Small-AES

Let \mathcal{S} be a structure that iterates over all 2^{16} values of one diagonal space of Small-AES and fixes all other cells to a constant. Then, after four rounds, the number of pairs that collide in any fixed byte is a multiple of four. This implies that, after five rounds without the final MixColumns operation (or inverted), the number of pairs that collide in any fixed byte is a multiple of four. This property is absent in the real AES with eight-bit S-boxes; it seems independent, however, from the S-box. It has been found also when using e.g., in variants of Small-AES with the S-boxes of PRESENT or Toy-10.

H.1 Five-round Distinguisher

The distinguisher is depicted in Figure 22.

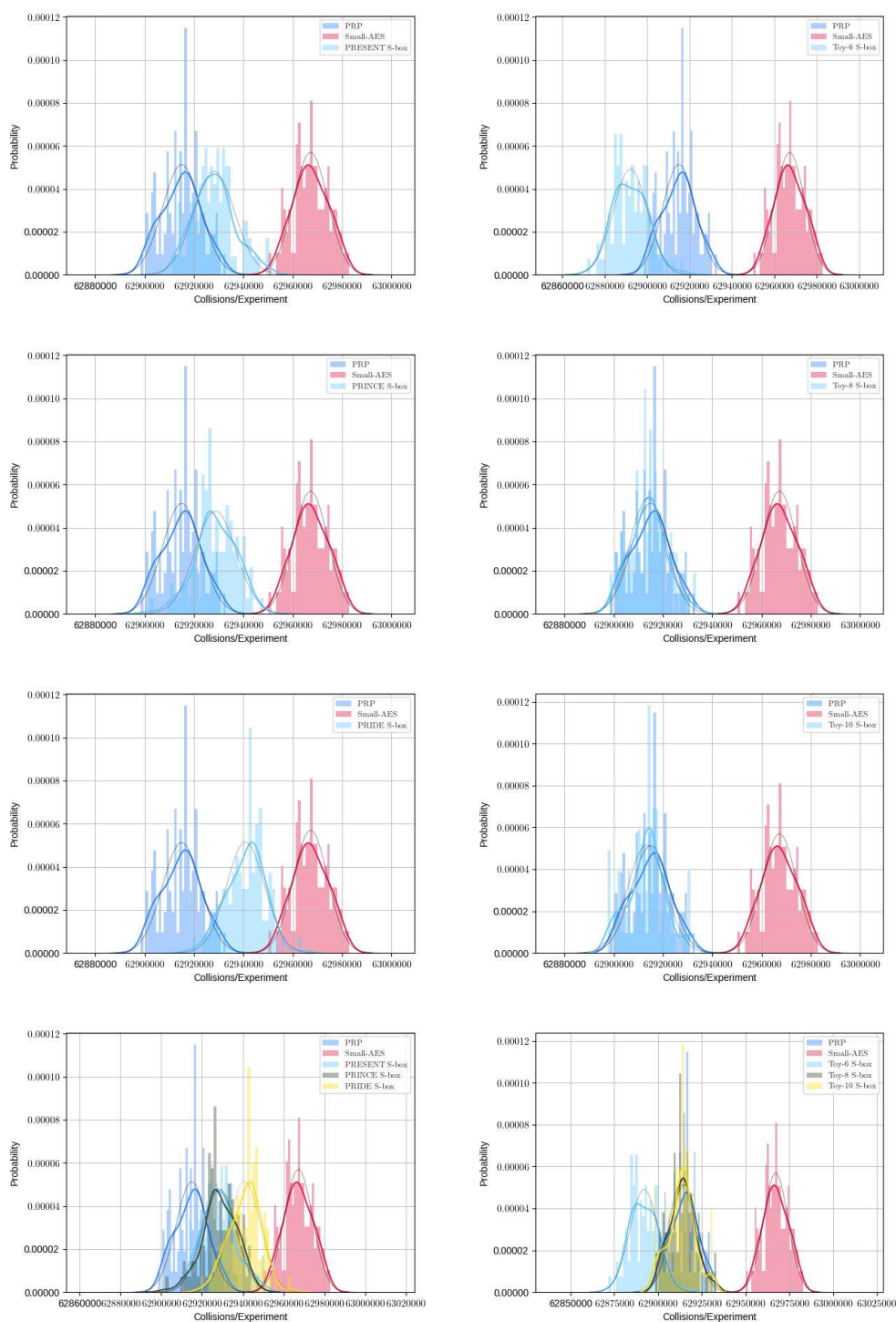


Figure 19: Distribution of means from 100 experiments with random keys for our four-round distinguisher on variants of Small-AES with different S-boxes. Each experiment used 2^{23} independent δ -sets.

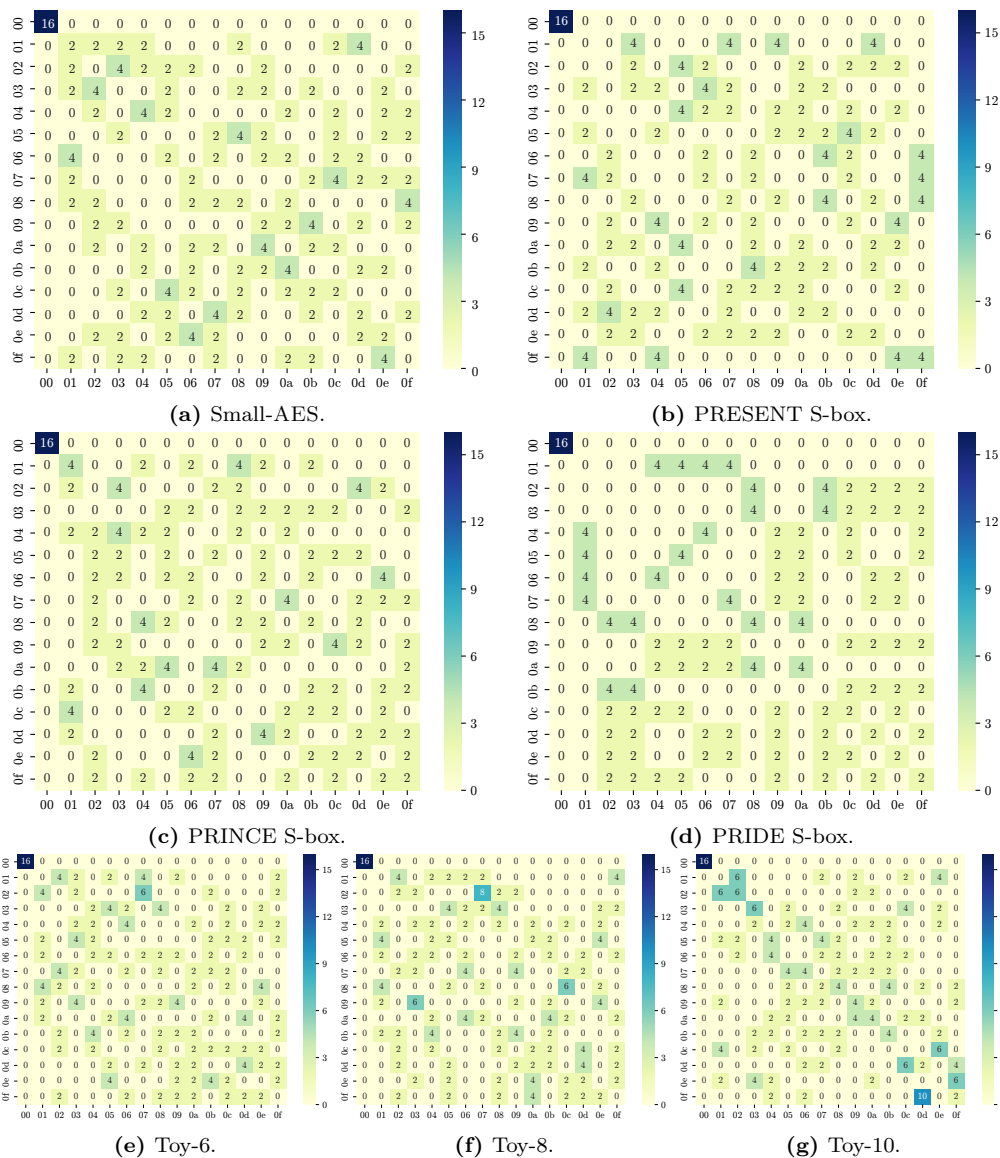
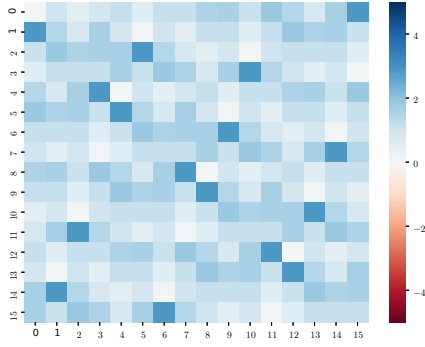
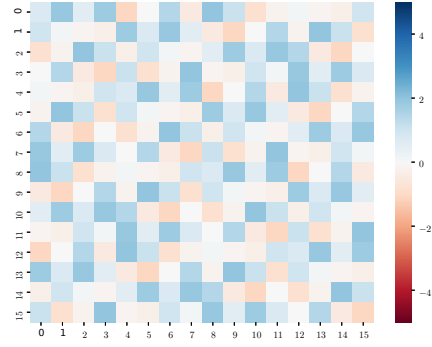


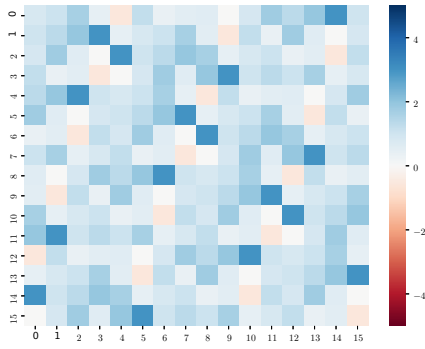
Figure 20: DDTs of the S-boxes.



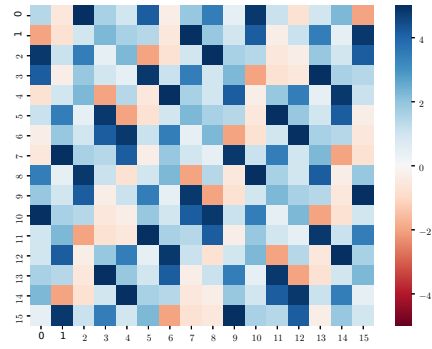
(a) Small-AES.



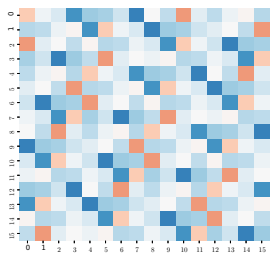
(b) PRESENT S-box.



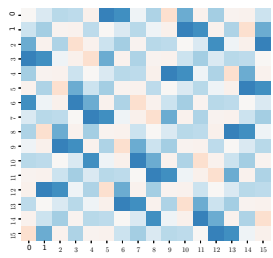
(c) PRINCE S-box.



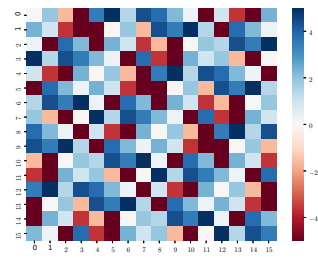
(d) PRIDE S-box.



(e) Toy-6 S-box.



(f) Toy-8 S-box.



(g) Toy-10 S-box.

Figure 21: Theoretical index-dependent means from the equation system for our four-round distinguisher with different S-boxes, in $\times |p_{\text{Small-AES}} - p_{\text{rand}}|$ deviations from the mean of the PRP. Each cell represents the result for the active input cell in the plaintexts (y-axis) and the tested inactive cell (x-axis) after four rounds without the final MixColumns operation. Darker colors (ignorant of sign) imply better distinguishers.

Table 12: Results of the four-round distinguisher on Small-AES with random S-boxes S_i and 100 random keys each. Speck-64-96 has been used as PRP π . **Var** = variance of the S-box.

	μ	σ	Var	S-box	μ	σ	Var
π	62 915 003.93	7 780.640		S_9	62 930 817.02	7 624.145	392/225
Small-AES	62 967 244.05	6 994.377	344/225	S_{10}	62 905 740.98	7 978.434	560/225
S_0	62 926 515.68	7 098.528	416/225	S_{11}	62 908 755.19	7 099.473	464/225
S_1	62 927 513.71	7 285.990	416/225	S_{12}	62 923 518.15	7 953.181	464/225
S_2	62 905 364.53	7 669.595	512/225	S_{13}	62 953 510.92	6 936.930	416/225
S_3	62 626 909.73	7 895.898	608/225	S_{14}	62 957 733.27	8 593.392	416/225
S_4	62 927 881.83	7 183.728	464/225	S_{15}	62 931 409.20	7 116.738	416/225
S_5	62 927 149.39	8 073.709	536/225	S_{16}	62 925 348.42	7 007.317	416/225
S_6	62 952 815.06	7 639.192	464/225	S_{17}	63 054 909.76	7 140.571	656/225
S_7	62 931 818.25	8 027.718	440/225	S_{18}	62 932 745.27	6 866.083	368/225
S_8	62 934 690.34	8 301.992	464/225	S_{19}	62 876 977.05	7 870.852	488/225

Steps. The steps of our distinguisher are as follows:

1. Initialize 16×16 collision counters $X_{i,j}$ for each byte and each value.
2. Collect a structure \mathcal{S} of 2^{16} texts that iterate over all values in four cells of a diagonal, e.g., using texts from $\mathcal{D}_{\{0\}} \oplus A$ for arbitrary $A \in \mathbb{F}_{2^4}^{4 \times 4}$, and leave the remaining bytes constants. Query the plaintexts of a structure and ask for their ciphertexts.
3. For each ciphertext, invert the final MixColumns operations to get the states $S^{5,SR}$. Increment its value j for Byte i , for $i \in \{0, \dots, 15\}$, in $X_{i,j}$.
4. For each byte $i \in \{0, \dots, 15\}$, test if $(\sum_{j=0}^{15} \binom{X_{i,j}}{2}) \bmod 4 = 0$. If yes, return real; return random otherwise.

Complexity. The distinguisher requires 2^{16} chosen plaintexts, 2^{16} encryptions and $2^{16} \cdot 16 = 2^{20}$ memory accesses. The success probability is one for the real construction; the probability is $(2^{-2})^{16} = 2^{-32}$ to wrongly identify a random permutation as the real five-round Small-AES. The memory is negligible, 256 16-bit values or 2^6 states of memory plus some constant overhead suffices.

H.2 Six-round Key-recovery Attack

The five-round distinguisher can be extended to a six-round key-recovery attack in a straight-forward manner.

Steps. The steps are as follows:

1. Initialize $4 \cdot 2^{16} \cdot 16 \times 16$ collision counters $X_{i,j}$ for each candidate of $K^6[0, 7, 10, 13]$, $K^6[3, 4, 9, 14]$, $K^6[2, 7, 8, 13]$, $K^6[1, 6, 11, 12]$, each byte, and each byte value. Initialize four lists \mathcal{K}_ℓ , for $\ell \in \{0, 1, 2, 3\}$.
2. Collect a structure \mathcal{S} of 2^{16} texts that iterate over all values in four cells of a diagonal, e.g., using texts from $\mathcal{D}_{\{0\}} \oplus A$ for arbitrary $A \in \mathbb{F}_{2^4}^{4 \times 4}$, and leave the remaining bytes constants. Query the plaintexts of a structure and ask for their ciphertexts.
3. For each candidate of $K^6[0, 7, 10, 13]$:

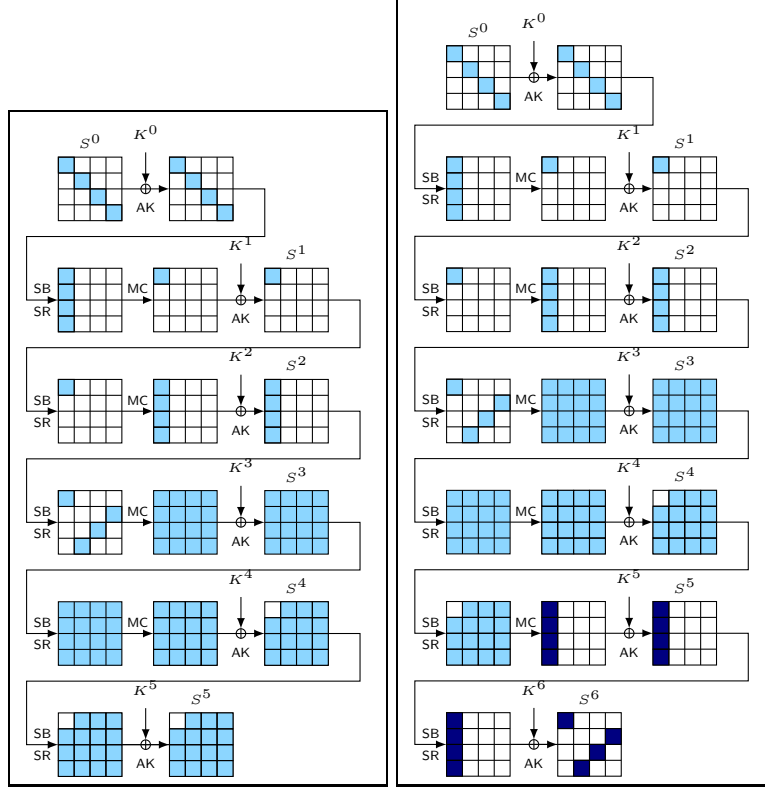
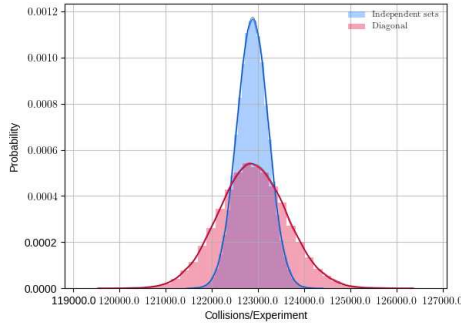


Figure 22: Multiple-of-four distinguisher on five-round Small-AES (**left**), and the key-recovery attack on six rounds (**right**).

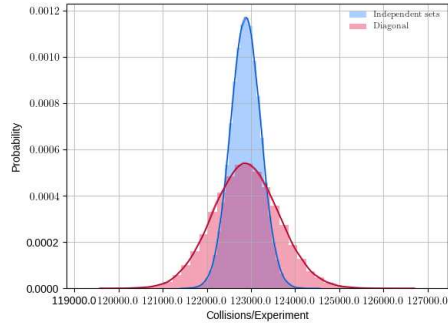
- 3.1 For each ciphertext, decrypt to the state $\widehat{S}^{5,SR}[0, 1, 2, 3]$ and increment its value j for Byte i , for $i \in \{0, \dots, 15\}$, in $X_{i,j}$.
- 3.2 For each counter, test if $(X_{2,j}) \bmod 4 = 0$, for $i \in \{0, 1, 2, 3\}$. If yes, add the surviving inverse-diagonal key candidate to list \mathcal{K}_0 .
- 3.3 Repeat the same procedure for the other inverse diagonals of K^6 , and add the surviving lists \mathcal{K}_ℓ for the corresponding values of ℓ .

4. Test the remaining key candidates.

Complexity. The distinguisher requires 2^{16} chosen plaintexts, 2^{16} encryptions, and $2^{16} \cdot 2^{16} \cdot 4 \cdot \frac{1}{4} \cdot \frac{1}{6}$ decryptions of one column through one of six rounds for each of four anti-diagonal key candidate and each of 2^{16} texts. This yields $2^{29.415}$ encryption equivalents. Each value of each inverse diagonal of the key has a chance of $2^{16} \cdot (2^{-2})^4 = 2^{-8}$ to survive. Thus, we can expect about $2^{4-8} = 2^{32}$ key candidates that can be tested with exhaustive search. Thus, the computational complexity is upper bounded by $2^{32} + 2^{29.415} \simeq 2^{32.223}$ Encryptions. The attack needs $4 \cdot 2^{16} \cdot 16 \cdot 16 = 2^{26}$ memory accesses. The memory is $4 \cdot 2^{16} \cdot 16 \cdot 16 = 2^{26}$ 64-bit values or states, and $4 \cdot 2^8 = 2^{10}$ values for the key candidates. We note that the time complexity of the attack can be decreased further with the meet-in-the-middle technique by Sasaki and Wang [SW12] that was adapted by Bar-On et al. to attacks on final-round AES [BODK⁺18] to about $2^{16} \cdot 2 \cdot 2^8 \simeq 2^{25}$ last-round decryptions. However, the property is the relevant thing in this section, and we provide distinguisher and key-recovery attack only for the sake of completeness.

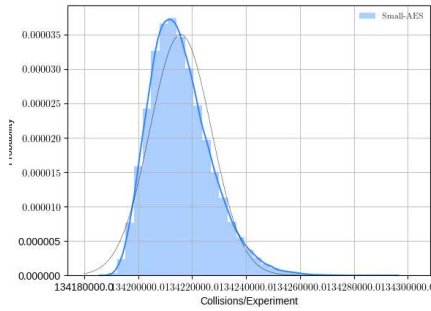


(a) Small-AES

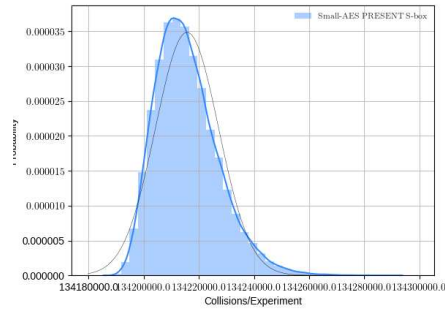


(b) Small-AES with PRESENT S-box

Figure 23: Distribution of the number of collisions in the first byte after five rounds in both settings, with the final MixColumns operation omitted. Both distributions used 100 000 random keys and counted the collisions from 2^{14} δ -sets per key that were independent in Setting (1) or the 2^{14} δ -sets from $\mathcal{D}_{\{0\}}$. Thin lines are fitted normal distributions, thick lines the best fitted distributions.



(a) Small-AES



(b) Small-AES with PRESENT S-box

Figure 24: Distribution of the number of collisions in the first byte after five rounds taking all $\binom{2^{16}}{2}$ pairs from a single plaintext diagonal. All distributions used 100 000 random keys and counted the collisions for one diagonal per key. Thin lines are fitted normal distributions, thick lines the best fitted distributions.

I Adapted Heys-Liu Probabilistic Integral Distinguisher on Four-round Small-AES

The Heys-Liu Integral Distinguisher on BSPN. In 2014, Heys [Hey14] conducted a similar analysis BSPN [YTH96], a byte-wise variant of the involutory SPN of [HT96]. BSPN transforms a vector of eight bytes through SubBytes, a mixing layer, and AddRoundKeys, similar to the AES and with the same S-box, but with a mixing layer that multiplies the state with the circulant matrix $\circ(0, 1, \dots, 1)$. The round function of BSPN is therefore simpler and a similar integral as for the AES applies. Heys and Liu studied a probabilistic extension of the three-round integral. More precisely, they studied the probability of obtaining a zero sum in the individual bytes after three rounds that naturally allows to be traced through the subsequent round. In contrast to our approach, they studied the probability that the sum of an entire δ -set of 256 texts is zero with probability significantly deviating from random.

(a) Mean and standard deviation of the Heys-Liu experiments adapted to four rounds of Small-AES with 100 experiments with random keys.

δ -sets s	Random permutation		Small-AES	
	μ	σ	μ	σ
24	1 048 573.120	964.778	1 049 134.150	1 079.363
26	4 194 297.190	2 182.769	4 196 155.600	2 382.283
28	16 776 997.485	3 808.268	16 785 415.790	5 787.742
30	67 109 083.490	8 114.173	67 143 877.360	16 479.876

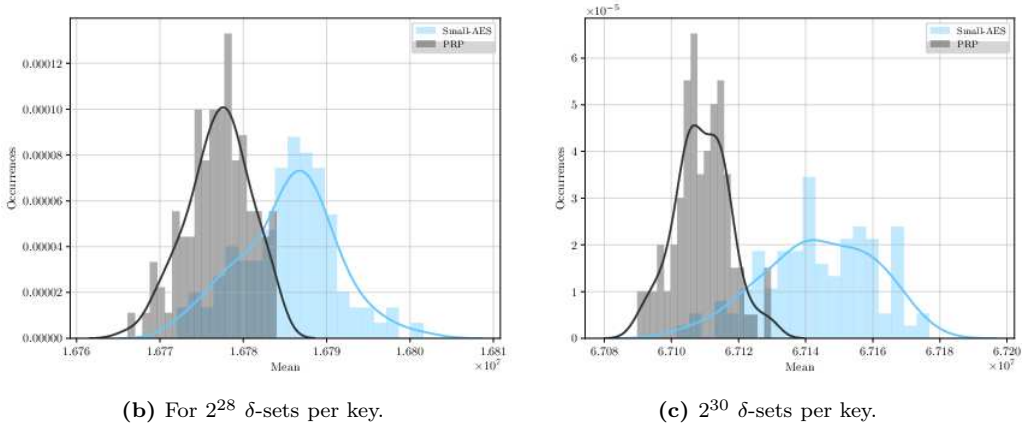


Figure 25: Results of the Heys-Liu experiments adapted to four rounds of Small-AES with 100 experiments with random keys.

Adaption to Small-AES. Inspired by their approach, we conducted a similar study for round-reduced AES. Since we expected the result to be inapplicable for the AES, we considered Small-AES instead. More precisely, we employed 100 random keys and 2^s δ -sets that iterated over all 2^4 values of the first input byte for each. We counted the number of times that we obtained a zero sum in the first output byte after almost four rounds, i.e., without the final MixColumns operation. Our results are illustrated in Figure 25 for $s \in \{28, 30\}$; detailed means and standard deviations are provided for those as well as for $s \in \{24, 26\}$ in Table 25a.

Interpretation. We see considerable differences in the distribution starting from $s = 28$, although still a broad overlap of both distributions. For $s = 30$, the different distributions become apparent. From our point of view, the distinguisher is a consequence of ours, but significantly more sophisticated to explore. We leave its more detailed foundations as an interesting future work.