

The Study of Modulo 2^n

A General Method To Calculate The Probability Or Correlation Coefficients For Most Statistical Property Of Modulo 2^n

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Abstract. In this paper, we present a new concept named the basic function. By the study of the basic function, we find the $O(n)$ -time algorithm to calculate the probability or correlation for some property of Modulo 2^n , including the difference-linear connective correlation coefficients, the linear approximation correlation coefficients, the differential probability, difference-boomerang connective probability, boomerang connective probability, boomerang-difference connective probability, etc.

Keywords: Modulo addition 2^n · Markov chain · basic function · difference-linear connective correlation coefficients · linear approximation correlation coefficients · difference-boomerang connective probability · boomerang connective probability · differential probability · boomerang-difference connective probability

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Introduction

Since the differential analysis and the linear analysis were proposed [1,2], statistical analysis, which makes use of some significant probability for the block cipher to distinguish from the random permutation, have become one of the research hotspots in the cryptanalysis of block cypher in the last 30 years. And most statistical analysis for block cypher is based on the statistical properties of nonlinear functions [1–5]. Specifically, most statistical properties of nonlinear functions is the probability that a series of boolean vector function with the form $\bigoplus_{i=1}^n f(x)_i$, where each $f(x)_i$ in the $\bigoplus_{i=1}^n f(x)_i$ is constituted by the composite

operation of \oplus , the nonlinear function and its inverse, are equal to some given values or the correlation coefficients of a single boolean vector function $\bigoplus_{i=1}^n f(x)_i$. In most of block cypher, the construction of nonlinear functions is based on the S-box that it can be regard as the nonlinear function with small scale, especially in the SPN structure and Fiestal structure, the calculation of statistical properties of nonlinear functions can be completed by calculating the statistical properties of S-box. Due to the S-box with tiny scale, the statistical properties of S-box can be got in a short time by trying all possible values. Thus, for the SPN cipher and the Fiestal cipher, it almost has no problem to do the statistical analysis, such as differential analysis, linear analysis, boomerang analysis, differential-linear analysis, etc. However, because of the modular addition 2^n with large-scale adopted as the nonlinear functions in ARX structure, it is infeasible for the ARX cipher to get the statistical properties of nonlinear functions by trying all possible value. In a word, compared with the SPN structure and the Fiestal structure, most studies of the utilize of statistical analysis method for the ARX cipher, especially the new method or improvement method proposed in the last decade, have made slow progress due to the above reason.

In order to overcome the above problem, we have to look for the polynomial time algorithms to calculate the statistical properties of modular addition 2^n . In 2001, Lipmaa.etc [6] and Johan Wallén [7] proposed the polynomial time algorithms to calculate the differential probability and the linear approximation correlation coefficients respectively, which was been about 10 years since the differential attack and the linear analysis were proposed. And in 2013, Schulte-Geers [8] showed that mod 2^n is CCZ-equivalent to a quadratic vectorial Boolean function. Based on it, he proposed the explicit formula to calculate the linear approximation correlation coefficients and the differential probability, of which time complex are polynomial. Since then, the CCZ-equivalent relation of mod 2^n had been acknowledged as the most powerful method to look for the polynomial time algorithms to calculate the statistical properties of modular addition 2^n . In addition, a series of new statistical properties, including the difference-linear connective correlation coefficients, difference-boomerang connective probability, boomerang connective probability and boomerang-difference connective probability, etc [1–5], were proposed in order to improve the previous methods. And those methods had better performance in the SPN cipher and the Fiestal cipher. Unfortunately, according to the study of difference-linear connective correlation coefficients for nonlinear function [9], for any two permutations in the same CCZ-equivalent class, their difference-linear connective correlation coefficients are not in general invariant. It means that the CCZ-equivalent relation can't be regard as the general method to look for polynomial time algorithms to calculate the statistical properties of modular addition 2^n . Then, the following questions may be asked naturally:

1. How to find the polynomial time algorithms to calculate such new proposed statistical properties?

2. Does there exists a general method that the explicit formula with polynomial time complex can be got for all of the current statistical properties of modular addition 2^n or even the properties that may come up in the future?

Our contribution

Firstly, this paper give the $O(n)$ -time algorithm to calculate the probability or correlation for newly proposed property of Modulo 2^n , including the difference-linear connective correlation coefficients, the linear approximation correlation coefficients, the differential probability, difference-boomerang connective probability, boomerang connective probability, boomerang-difference connective probability, etc.

Secondly, for all of the current statistical property used for statistical analysis in block cipher, such as difference-linear connective correlation coefficients, the linear approxima-

tion correlation coefficients, the differential probability, difference-boomerang connective probability, boomerang connective probability, boomerang-difference connective probability, etc, it can be summarized as the probability of $\bigoplus_{z=1}^{n_j} f(x)_{z,j} = A_j$, $1 \leq j \leq m$ or the correlation coefficients of $\lambda \cdot (\bigoplus_{i=1}^n f(x)_i)$, where $f(x)_{z,j}$, $1 \leq j \leq m$, $1 \leq z \leq n_j$; $f(x)_i$, $1 \leq i \leq n$ are composite operation of \oplus , the nonlinear function and its inverse. Based the above fact, we propose the concept of the basic function, which is a composite operation of \oplus and \boxplus . Then, based on the regular of the basic function, we construct the Aa-set and Df-set. According to the property of the Aa-set and Df-set, we can answer the question 2. The general method has been found, which is fit for all of current statistical properties of modular addition 2^n . As a result, the form of formula for all of current statistical properties of modular addition 2^n are similar with the Johan Wallén's work.

1 Preliminaries

In this section, we will introduce some basic knowledge that we will use in the following.

Definition 1 (Addition modulo 2^n): For $x, y \in F_2^n$, define $y \boxplus x = x \oplus y \oplus \text{carry}(x, y)$, where $\text{carry}(x, y) = [c_{n-1}, \dots, c_0]$. The i -th bit c_i is defined as

$$\begin{aligned} c_0 &= 0 \\ c_{i+1} &= (x_i \wedge y_i) \oplus (x_i \wedge c_i) \oplus (y_i \wedge c_i), 0 \leq i \leq n-1 \end{aligned}$$

Definition 2: Let $x, y \in F_2^n$, $e \in F_2$, define $y \boxplus x \boxplus e = y \boxplus x \boxplus (0, \dots, 0, e)$.

Let $\text{carry}_e^*(x, y) = [c_{n-1}^*, \dots, c_0^*]$, if we define i -th bit c_i^* as

$$\begin{aligned} c_0^* &= e \\ c_{i+1}^* &= (x_i \wedge y_i) \oplus (x_i \wedge c_i^*) \oplus (y_i \wedge c_i^*), 0 \leq i \leq t-1 \end{aligned}$$

And the purpose of defining $\text{carry}_e^*(x, y)$ is to make $y \boxplus x \boxplus e$ have the same form as $y \boxplus x$:

Theorems 1: For $x, y \in F_2^n$, $e \in F_2$, define $y \boxplus x \boxplus e = x \oplus y \oplus \text{carry}_e^*(x, y)$.

Proof. If $e = 0$, then the theorem holds.

If $e = 1$, let $c_i = \text{carry}(x, y)[i]$, $c_i^1 = \text{carry}(x \boxplus y, (0, \dots, 0, e))[i]$, for $0 \leq i \leq n$; then

$$\begin{aligned} c_0^1 &= 0 \\ c_1^1 &= x_0 \oplus y_0 \\ c_{i+1}^1 &= (x_i \oplus y_i \oplus c_i) \wedge c_i^1, 2 \leq i \leq n-1 \end{aligned}$$

Obviously, $c_0^1 \wedge c_0 = 0$, $c_1^1 \wedge c_1 = x_0 \wedge y_0 \wedge (x_0 \oplus y_0) = (x_0 \wedge y_0) \oplus (x_0 \wedge y_0) = 0$.

Next, we will proof $c_i^1 \wedge c_i = 0$, for $0 \leq i \leq n-1$, by induction. Supposed that $c_i^1 \wedge c_i = 0$ for $1 \leq i \leq k$. Then, for $c_{k+1}^1 \wedge c_{k+1}$, we have

$$\begin{aligned} &c_{k+1}^1 \wedge c_{k+1} \\ &= ((x_k \oplus y_k \oplus c_k) \wedge c_k^1) \wedge ((x_k \wedge y_k) \oplus (x_k \wedge c_k) \oplus (y_k \wedge c_k)) \\ &= ((x_k \oplus y_k) \wedge c_k^1) \wedge ((x_k \wedge y_k) \oplus ((x_k \oplus y_k) \wedge c_k)) \\ &= ((x_k \oplus y_k) \wedge (x_k \wedge y_k)) \wedge c_k^1 \\ &= ((x_k \wedge y_k) \oplus (x_k \wedge y_k)) \wedge c_k^1 = 0 \end{aligned}$$

Thus, for $2 \leq i \leq n-1$, we have $c_{i+1}^1 = (x_i \oplus y_i \oplus c_i) \wedge c_i^1 = (x_i \oplus y_i) \wedge c_i^1$, and

$$\begin{aligned} c_0^* &= c_0 \oplus c_0^1 \oplus 1 = 1 \\ c_1^* &= c_1 \oplus c_1^1 \\ c_{i+1}^* &= c_{i+1}^1 \oplus c_{i+1} \\ &= (x_i \wedge y_i) \oplus (x_i \wedge (c_i^1 \oplus c_i)) \oplus (y_i \wedge (c_i^1 \oplus c_i)), 2 \leq i \leq n-1 \end{aligned}$$

Notice that $\text{carry}_e^*(x, y) = \text{carry}(x, y) \oplus \text{carry}(x \boxplus y, (0, \dots, 0, e)) \oplus (0, \dots, 0, e)$.
Thus, $y \boxplus x \boxplus (0, \dots, 0, e) = x \oplus y \oplus \text{carry}_e^*(x, y)$. \square

From the definition of \boxplus , we can see that the subtraction modulo 2^n (\boxminus) can be converted into the addition modulo 2^n (\boxplus):

Theorems 2: $y \boxminus x = (x \oplus (1, \dots, 1)) \boxplus y \boxplus 1$.

Proof. Notice that $(1, \dots, 1) = -1 \text{ mod } 2^n$ and $\text{carry}(x \oplus (1, \dots, 1), x) = 0^n$, then

$$\begin{aligned} (x \oplus (1, \dots, 1)) \boxplus x &= (1, \dots, 1) = -1 \text{ mod } 2^n, \text{ which is equal to} \\ (x \oplus (1, \dots, 1)) \boxplus 1 &= -x \text{ mod } 2^n. \end{aligned}$$

Thus, $y \boxminus x = y - x \text{ mod } 2^n = (x \oplus (1, \dots, 1)) \boxplus y \boxplus 1$. \square

Combing with the **theorems 3**, the $y \boxminus x$ have the same form as $y \boxplus x$:

Corollary 1: $y \boxminus x = (x \oplus (1, \dots, 1)) \oplus y \oplus \text{carry}_1^*(x \oplus (1, \dots, 1), y)$.

2 The Basic Function

2.1 The motivation and property

In the cryptanalysis of block cypher, the scholar proposed many analysis method based on some statistics property of the nonlinear function in the cipher. And these statistics properties can be summarized as the probability of $\bigoplus_{z=1}^{n_j} f(x)_{z,j} = A_j$, $1 \leq j \leq m$ or the correlation coefficients of $\lambda \cdot (\bigoplus_{i=1}^n f(x)_i)$, where $f(x)_{z,j}$, $1 \leq j \leq m$, $1 \leq z \leq n_j$; $f(x)_i$, $1 \leq i \leq n$ are composite operation of \oplus , the nonlinear function and its inverse. In the ARX cipher, the sole nonlinear function is the \boxplus and its inverse can be converted into the composite operation of \oplus and \boxplus . Thus, we can define the basic function as composite operation of \oplus and \boxplus :

Definition 3(The Basic function): Supposed that $x, y \in F_2^n$, $E_k = (e_0, e_1, \dots, e_{k-1}) \in F_2^k$, $\alpha_0, \alpha_1, \dots, \alpha_{k-1} \in F_2^n$, $\beta_0, \beta_1, \dots, \beta_{k-1} \in F_2^n$. Let $A_k = (\alpha_0, \alpha_2, \dots, \alpha_{k-1})$, $B_k = (\beta_0, \beta_2, \dots, \beta_{k-1})$, where the element of A_k, B_k is n -dimension vector. Then the $f(x, y)_{E_k, A_k, B_k}$ is called basic function with k order, if $f(x, y)_{E_k, A_k, B_k}$ satisfies the follow the form:

$$\begin{aligned} (x_0, y_0) &= (x, y) \\ (x_{i+1}, y_{i+1}) &= (x_i \oplus \alpha_i, (x_i \oplus \alpha_i) \boxplus (y_i \oplus \beta_i) \boxplus e_i), 0 \leq i \leq k-1 \\ f(x, y)_{E_k, A_k, B_k} &= (x_k, y_k). \end{aligned}$$

Then, for $0 \leq i \leq k-1$, $(x_i \oplus \alpha_i, (x_i \oplus \alpha_i) \boxplus (y_i \oplus \beta_i) \boxplus e_i)$ is called the $(i+1)$ -th round function, and $\text{carry}_{e_i}^*(x_i \oplus \alpha_i, y_i \oplus \beta_i)$ is called the carry function of $(i+1)$ -th round function in $f(x, y)_{E_k, A_k, B_k}$.

According to the **definition 3**, we can see that the relation between the $f(x, y)_{E_{k-1}, A_{k-1}, B_{k-1}}$ and the $f(x, y)_{E_k, A_k, B_k}$:

Remark 3: According to the definition of $f(x, y)_{E_k, A_k, B_k}$, the $f(x, y)_{E_k, A_k, B_k}$ can be written as:

$$\begin{aligned} (x_{k-1}, y_{k-1}) &= f(x, y)_{E_{k-1}, A_{k-1}, B_{k-1}} \\ (x_k, y_k) &= (x_{k-1} \oplus \alpha_{k-1}, (x_{k-1} \oplus \alpha_{k-1}) \boxplus (y_{k-1} \oplus \beta_{k-1}) \boxplus e_{k-1}) \\ f(x, y)_{E_k, A_k, B_k} &= (x_k, y_k) \end{aligned}$$

where $E_{k-1} = E_k[0 : k-2] = (e_0, e_1, \dots, e_{k-2})$, $A_{k-1} = A_k[0 : k-2] = (\alpha_0, \alpha_2, \dots, \alpha_{k-2})$, $B_{k-1} = B_k[0 : k-2] = (\beta_0, \beta_2, \dots, \beta_{k-2})$.

Definition 4: For $x = X^2 || X^1 \in F_2^n$, where $X^2 \in F_2^q$, $X^1 \in F_2^p$, $p + q = n$. Define $x = X^2 || X^1 = X^2 \cdot 2^p + X^1$.

Then, according to the definition of Addition modulo 2^n , we can see that the basic function can be divided into two basic function:

Theorem 3: For any $X^{1,i+1}, Y^{1,i+1}, \alpha_0^{1,i+1}, \alpha_1^{1,i+1}, \dots, \alpha_{k-1}^{1,i+1}, \beta_0^{1,i+1}, \beta_1^{1,i+1}, \dots, \beta_{k-1}^{1,i+1} \in F_2^{(i+1) \cdot t}$ and $E_k = (e_0, e_1, \dots, e_{k-1}) \in F_2^k$. Let $Y^{1,i} = Y^{1,i+1}[i \cdot t - 1 : 0]$, $Y^{2,i} = Y^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t]$, $X^{1,i} = X^{1,i+1}[i \cdot t - 1 : 0]$, $X^{2,i} = X^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t]$, then the k order basic function $f(X^{1,i+1}, Y^{1,i+1})_{E_k, A_k^{i+1}, B_k^{i+1}}$ can be written as

$$f(X^{1,i+1}, Y^{1,i+1})_{E_k, A_k^{i+1}, B_k^{i+1}} = f(X^{2,i}, Y^{2,i})_{M_k^i, C_k^i, D_k^i} \cdot 2^{i \cdot t} + f(X^{1,i}, Y^{1,i})_{E_k, A_k^i, B_k^i}$$

where

$$\begin{aligned} \alpha_m^{2,i} &= \alpha_m^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t], 0 \leq m \leq k-1 \\ \beta_m^{2,i} &= \beta_m^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t], 0 \leq m \leq k-1 \\ \alpha_m^{1,i} &= \alpha_m^{1,i+1}[i \cdot t - 1 : 0], 0 \leq m \leq k-1 \\ \beta_m^{1,i} &= \beta_m^{1,i+1}[i \cdot t - 1 : 0], 0 \leq m \leq k-1 \\ c_m^i &= \text{carry}_{e_m}^*(X_m^{1,i} \oplus \alpha_m^{1,i}, Y_m^{1,i} \oplus \beta_m^{1,i})[i \cdot t], 0 \leq m \leq k-1 \\ s_m^i &= \text{carry}_{c_m^i}^*(X_m^{2,i} \oplus \alpha_m^{2,i}, Y_m^{2,i} \oplus \beta_m^{2,i})[t], 0 \leq m \leq k-1 \\ A_k^i &= [\alpha_0^{1,i}, \alpha_1^{1,i}, \dots, \alpha_{k-1}^{1,i}] \\ B_k^i &= [\beta_0^{1,i}, \beta_1^{1,i}, \dots, \beta_{k-1}^{1,i}] \\ C_k^i &= [\alpha_0^{2,i}, \alpha_1^{2,i}, \dots, \alpha_{k-1}^{2,i}] \\ D_k^i &= [\beta_0^{2,i}, \beta_1^{2,i}, \dots, \beta_{k-1}^{2,i}] \\ M_k^i &= [c_0^i, c_1^i, \dots, c_{k-1}^i] \\ S_k^i &= [s_0^i, s_1^i, \dots, s_{k-1}^i] \end{aligned}$$

Moreover, $S_k^i = M_k^{i+1}$.

Proof. Notice that when order $k = 1$, according to the **definition 2**, $(X^{1,i+1} \oplus \alpha_0^{1,i+1}) \boxplus (Y^{1,i+1} \oplus \beta_0^{1,i+1}) \boxplus e_0$ can be written as

$$\begin{aligned} &(X^{1,i+1} \oplus \alpha_0^{1,i+1}) \boxplus (Y^{1,i+1} \oplus \beta_0^{1,i+1}) \boxplus e_0 \\ &= ((X^{1,i} \oplus \alpha_0^{1,i}) \boxplus (Y^{1,i} \oplus \beta_0^{1,i}) \boxplus c_0^i) 2^{i \cdot t} + (X^{2,i} \oplus \alpha_0^{2,i}) \boxplus (Y^{2,i} \oplus \beta_0^{2,i}) \boxplus e_0 \end{aligned}$$

Thus, when order $k = 1$, the theorem holds.

Supposed that when order $m \leq k$, the theorem holds.

When order $m = k + 1$, according to the definition of $\text{carry}_e^*(x, y)$, the value of the $i \cdot t - th$

bit of $\text{carry}_{e_k}^*(X_k^{1,i+1} \oplus \alpha_k^{1,i+1}, Y_k^{1,i+1} \oplus \beta_k^{1,i+1})[i \cdot t]$ is only rely on the first $i \cdot t - 1$ bits of $X_k^{1,i+1} \oplus \alpha_k^{1,i+1}$ and $Y_k^{1,i+1} \oplus \beta_k^{1,i+1}$, namely,

$$\begin{aligned} & \text{carry}_{e_k}^*(X_k^{1,i+1} \oplus \alpha_k^{1,i+1}, Y_k^{1,i+1} \oplus \beta_k^{1,i+1})[i \cdot t] \\ &= \text{carry}_{e_k}^*(X_k^{1,i} \oplus \alpha_k^{1,i}, Y_k^{1,i} \oplus \beta_k^{1,i})[i \cdot t] \\ &= c_k^i \end{aligned}$$

Thus,

$$\begin{aligned} & f(X^{1,i+1}, Y^{1,i+1})_{E_{k+1}, A_{k+1}^{i+1}, B_{k+1}^{i+1}} \\ &= (X_k^{1,i+1} \oplus \alpha_k^{1,i+1}) \boxplus (Y_k^{1,i+1} \oplus \beta_k^{1,i+1}) \boxplus e_k \\ &= ((X_k^{2,i} \oplus \alpha_k^{2,i}) \boxplus (Y_k^{2,i} \oplus \beta_k^{2,i}) \boxplus c_k^i) 2^{i \cdot t} + (X_k^{1,i} \oplus \alpha_k^{1,i}) \boxplus (Y_k^{1,i} \oplus \beta_k^{1,i}) \boxplus e_k \end{aligned}$$

On the other hand, due to the assumption of induction, we have:

$$\begin{aligned} (X_k^{1,i+1}, X_k^{1,i+1}) &= (X_k^{2,i}, X_k^{2,i}) \cdot 2^{i \cdot t} + (X_k^{1,i}, X_k^{1,i}) \\ &= f(X^{2,i}, Y^{2,i})_{M_k^i, C_k^i, D_k^i} + f(X^{1,i}, Y^{1,i})_{E_k, A_k^i, B_k^i} \end{aligned}$$

Thus,

$$\begin{aligned} & f(X^{1,i+1}, Y^{1,i+1})_{E_{k+1}, A_{k+1}^{i+1}, B_{k+1}^{i+1}} = (X_k^{1,i+1} \oplus \alpha_k^{1,i+1}) \boxplus (Y_k^{1,i+1} \oplus \beta_k^{1,i+1}) \boxplus e_k \\ &= ((X_k^{2,i} \oplus \alpha_k^{2,i}) \boxplus (Y_k^{2,i} \oplus \beta_k^{2,i}) \boxplus c_k^i) \cdot 2^{i \cdot t} + (X_k^{1,i} \oplus \alpha_k^{1,i}) \boxplus (Y_k^{1,i} \oplus \beta_k^{1,i}) \boxplus e_k \\ &= f(X^{2,i}, Y^{2,i})_{M_{k+1}^i, C_{k+1}^i, D_{k+1}^i} \cdot 2^{i \cdot t} + f(X^{1,i}, Y^{1,i})_{E_{k+1}, A_{k+1}^i, B_{k+1}^i} \end{aligned}$$

When $m = k + 1$, the theorem holds. \square

Remark 4: Obviously, for any $E_k, A_k^{i+1}, B_k^{i+1}$, when $X^{1,i+1}, Y^{1,i+1}$ are given, then M_k^{i+1} are uniquely identified.

3 Notion Description About The Basic Function Series

In order to reduce the redundancy of the article, we will introduce some notion description about the given z basic function series $\{f(X^{1,i}, Y^{1,i})_{E_{k_m}, A_{k_m}^i, B_{k_m}^i}; 1 \leq m \leq z\}$, which will be frequently adopted in the following proof.

Supposed that $X, Y \in F_2^{q \cdot t}$. For $1 \leq m \leq z$, let $\alpha_{0,m}^{1,q}, \alpha_{1,m}^{1,q}, \dots, \alpha_{k_m-1,m}^{1,q}, \beta_{0,m}^{1,q}, \beta_{1,m}^{1,q}, \dots, \beta_{k_m-1,m}^{1,q} \in F_2^{q \cdot t}$, $E_m = (e_0, e_1, \dots, e_{k_m-1}) \in F_2^{k_m}$.

In addition, let

$$\begin{aligned} X^{1,q} &= X \\ Y^{1,q} &= Y \\ A_m^q &= [\alpha_{0,m}^{1,q}, \alpha_{1,m}^{1,q}, \dots, \alpha_{k_m-1,m}^{1,q}], \\ B_m^q &= [\beta_{0,m}^{1,q}, \beta_{1,m}^{1,q}, \dots, \beta_{k_m-1,m}^{1,q}]; \end{aligned}$$

And for $1 \leq i \leq q-1$, define:

$$\begin{aligned} Y^{1,i} &= Y^{1,i+1}[i \cdot t - 1 : 0] = Y^{1,q}[i \cdot t - 1 : 0] \\ Y^{2,i} &= Y^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t] = Y^{1,q}[(i+1) \cdot t - 1 : i \cdot t] \\ X^{1,i} &= X^{1,i+1}[i \cdot t - 1 : 0] = X^{1,q}[i \cdot t - 1 : 0] \\ X^{2,i} &= X^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t] = X^{1,q}[(i+1) \cdot t - 1 : i \cdot t] \\ \alpha_{j,m}^{2,i} &= \alpha_{j,m}^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t] = \alpha_{j,m}^{1,q}[(i+1) \cdot t - 1 : i \cdot t], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \\ \beta_{j,m}^{2,i} &= \beta_{j,m}^{1,i+1}[(i+1) \cdot t - 1 : i \cdot t] = \beta_{j,m}^{1,q}[(i+1) \cdot t - 1 : i \cdot t], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \\ \alpha_{j,m}^{1,i} &= \alpha_{j,m}^{1,i+1}[i \cdot t - 1 : 0] = \alpha_{j,m}^{1,q}[i \cdot t - 1 : 0], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \\ \beta_{j,m}^{1,i} &= \beta_{j,m}^{1,i+1}[i \cdot t - 1 : 0] = \beta_{j,m}^{1,q}[i \cdot t - 1 : 0], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \end{aligned}$$

Beside this, by the ,for $1 \leq i \leq q$, $1 \leq m \leq z$, let

$$\begin{aligned} (X_{0,m}^{1,i}, y_{0,m}^{1,i}) &= (X^{1,i}, Y^{1,i}) \\ (X_{j+1,m}^{1,i}, Y_{j+1,m}^{1,i}) &= (X_{j,m}^{1,i} \oplus \alpha_{j,m}^{1,i}, (X_{j,m}^{1,i} \oplus \alpha_{j,m}^{1,i}) \boxplus (Y_{j,m}^{1,i} \oplus \beta_{j,m}^{1,i}) \boxplus e_{j,m}), 0 \leq j \leq k_m - 1 \end{aligned}$$

and for $1 \leq i \leq q$, let

$$\begin{aligned} A_m^i &= [\alpha_{0,m}^{1,i}, \alpha_{1,m}^{1,i}, \dots, \alpha_{k_m-1,m}^{1,i}], 1 \leq m \leq z \\ B_m^i &= [\beta_{0,m}^{1,i}, \beta_{1,m}^{1,i}, \dots, \beta_{k_m-1,m}^{1,i}], 1 \leq m \leq z \end{aligned}$$

then we have $(X_{k_m,j}^{1,i}, Y_{k_m,j}^{1,i}) = f(X^{1,i}, Y^{1,i})_{E_m, A_m^{i+1}, B_m^{i+1}}$, for $1 \leq m \leq z$.
Secondly, we can define:

$$\begin{aligned} c_{j,m}^i &= \text{carry}_{e_{j,m}}^*(X_j^{1,i} \oplus \alpha_{j,m}^{1,i}, Y_j^{1,i} \oplus \beta_{j,m}^{1,i})[i \cdot t], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \\ M_m^i &= [c_{0,m}^i, c_{1,m}^i, \dots, c_{k_m,j-1,m}^i], 1 \leq m \leq z \end{aligned}$$

where $1 \leq i \leq q$.

$$\begin{aligned} s_{j,m}^i &= \text{carry}_{c_{j,m}^i}^*(X_j^{2,i} \oplus \alpha_{j,m}^{2,i}, Y_j^{2,i} \oplus \beta_{j,m}^{2,i})[t], 0 \leq j \leq k_m - 1, 1 \leq m \leq z \\ S_{j,m}^i &= [s_{0,m}^i, s_{1,m}^i, \dots, s_{k_m,j-1,m}^i], 1 \leq m \leq z. \end{aligned}$$

where $1 \leq i \leq q-1$.

According to the **Remark 4**:, we have

Remark 6: For $1 \leq m \leq r_j$, $1 \leq i \leq q-1$, $M_{k_m,j}^{i+1} = S_{k_m,j}^i$.

For convenience to the follow following discussion, for $1 \leq i \leq q-1$, let

$$\begin{aligned} T_m^0 &= A_m^1, 1 \leq m \leq z \\ D_m^0 &= B_m^1, 1 \leq m \leq z \\ T_m^i &= [\alpha_{0,m}^{2,i}, \alpha_{1,m}^{2,i}, \dots, \alpha_{k_m-1,m}^{2,i}], 1 \leq m \leq z \\ D_m^i &= [\beta_{0,m}^{2,i}, \beta_{1,m}^{2,i}, \dots, \beta_{k_m-1,m}^{2,i}], 1 \leq m \leq z \end{aligned}$$

4 The Aa-set

Definition 6: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. We define the Aa-set with Out , In as

$$\begin{aligned} Aa_{Out, In}^{i+1} &(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \{(X^{1,i+1}, Y^{1,i+1}) | X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}, Out = (M_1^{i+1}, \dots, M_z^{i+1})\} \end{aligned}$$

where $Out, In \in F_2^d$, $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$.

Let $d = \sum_{i=0}^z k_{r_i}$, then there are 2^d possible results for $(M_1^{i+1}, \dots, M_z^{i+1})$, thus

Property 3: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Then, for all $Out \in F_2^d$, the Aa-set with Out, In satisfies:

$$\bigcup_{Out \in F_2^d} Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \left\{ (X_{1,i+1}, Y_{1,i+1}) \mid X_{1,i+1}, Y_{1,i+1} \in F_2^{(i+1)t} \right\}$$

According to the **Property 2**, we have:

Property 4: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. For any $Out_1, Out_2 \in F_2^d$ satisfied $Out_1 \neq Out_2$, then

$$Aa_{Out_1, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Aa_{Out_2, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

holds.

From **theorem 3**, we can see that the Aa-set with Out, In can be divided into many subset, and each disjoint subset has the following recursive structure:

Lemma 1: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. We define the Aa-set with Out, Mi, In as

$$\begin{aligned} & Aa_{Out, Mi, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ (X^{2,i} \| X^{1,i}, Y^{2,i} \| Y^{1,i}) \mid \begin{array}{l} X^{1,i}, Y^{1,i} \in Aa_{Mi, In}^i(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z); \\ Mi = (M_1^i, \dots, M_z^i); \\ X^{2,i}, Y^{2,i} \in Aa_{Out, Mi}^1(T_j^i, D_j^i, 1 \leq j \leq z). \end{array} \right\}. \end{aligned}$$

where $Mi, Out, In \in F_2^d$, $d = \sum_{i=0}^z k_{r_i}$, $In = (E_1, \dots, E_z)$. Then,

$$Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \bigcup_{Mi \in F_2^d} Aa_{Out, Mi, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$$

holds, and for $Mi_1 \neq Mi_2$, $Mi_1, Mi_2 \in F_2^d$, satisfy

$$Aa_{Out, Mi_1, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Aa_{Out, Mi_2, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

Proof. Firstly, due to A_j^i, B_j^i can be decided according to the definition from the A_j^{i+1}, B_j^{i+1} . And from the **property 4**, we know that for $Mi_1 \neq Mi_2$,

$$Aa_{Mi_1, In}^i(A_j^i, B_j^i, 1 \leq j \leq z) \cap Aa_{Mi_2, In}^i(A_j^i, B_j^i, 1 \leq j \leq z) = \emptyset$$

Thus,

$$Aa_{Out, Mi_1, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Aa_{Out, Mi_2, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

Secondly, from **theorem 3**, for $1 \leq j \leq z$, we have:

$$f(X^{1,i+1}, Y^{1,i+1})_{E_j, A_j^{i+1}, B_j^{i+1}} = f(X^{2,i}, Y^{2,i})_{M_j^i, T_j^i, D_j^i} \cdot 2^{i \cdot t} + f(X^{1,i}, Y^{1,i})_{E_j, A_j^i, B_j^i}.$$

Thus,

$$\begin{aligned}
& Aa_{Out,In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= \{(X^{1,i+1}, Y^{1,i+1}) \mid X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1)\cdot t}, Out = (M_1^{i+1}, \dots, M_z^{i+1})\} \\
&= \left\{ (X^{2,i} \parallel X^{1,i}, Y^{2,i} \parallel Y^{1,i}) \mid \begin{array}{l} X^{1,i}, Y^{1,i} \in F_2^{i\cdot t}, X^{2,i}, Y^{2,i} \in F_2^t, Mi \in F_2^d, \\ Mi = (M_1^i, \dots, M_z^i), Out = (S_1^i, \dots, S_z^i). \end{array} \right\} \\
&= \bigcup_{Mi \in F_2^d} \left\{ (X^{2,i} \parallel X^{1,i}, Y^{2,i} \parallel Y^{1,i}) \mid \begin{array}{l} X^{1,i}, Y^{1,i} \in F_2^{i\cdot t}; Mi = (M_1^i, \dots, M_z^i); \\ X^{2,i}, Y^{2,i} \in F_2^t; Out = (S_1^i, \dots, S_z^i). \end{array} \right\} \\
&= \bigcup_{Mi \in F_2^d} \left\{ (X^{2,i} \parallel X^{1,i}, Y^{2,i} \parallel Y^{1,i}) \mid \begin{array}{l} X^{1,i}, Y^{1,i} \in Aa_{Mi,In}^i(A_j^i, B_j^i, 1 \leq j \leq z), \\ Mi = (M_1^i, \dots, M_z^i); \\ X^{2,i}, Y^{2,i} \in Aa_{Out,Mi}^1(T_j^i, D_j^i, 1 \leq j \leq z). \end{array} \right\} \\
&= \bigcup_{Mi \in F_2^d} Aa_{Out,Mi,In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)
\end{aligned}$$

□

Definition 7: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1)\cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $\gamma, \lambda, v, w \in F_2^{(i+1)\cdot t}$, $Out, In, Middle \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Define the correlation coefficients of $h(X^{1,i+1}, Y^{1,i+1})$ as

$$\begin{aligned}
& Cor_{In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\
&= \frac{1}{2^{2(i+1)\cdot t}} \sum_{\substack{X_{1,i+1} \in F_2^{(i+1)\cdot t}, \\ Y_{1,i+1} \in F_2^{(i+1)\cdot t}}} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}}
\end{aligned}$$

where $h(X^{1,i+1}, Y^{1,i+1}) = \bigoplus_{m=1}^z f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}$.

Additionally define the correlation coefficients of $h(X^{1,i+1}, Y^{1,i+1})$ over the Aa-set with Out, In and the correlation coefficients of $h(X^{1,i+1}, Y^{1,i+1})$ over the Aa-set with Out, Mi, In respectively as

$$\begin{aligned}
& Cor_{Out,In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\
&= \frac{1}{2^{2(i+1)\cdot t}} \sum_{(X_{1,i+1}, Y_{1,i+1}) \in Aa_{Out,In}^{i+1}(A_j^{i+1}, B_j^{i+1})} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}} \\
& Cor_{Out,Mi,In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\
&= \frac{1}{2^{2(i+1)\cdot t}} \sum_{(X_{1,i+1}, Y_{1,i+1}) \in Aa_{Out,Mi,In}^{i+1}(A_j^{i+1}, B_j^{i+1})} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}}
\end{aligned}$$

From the property of the Aa-set, we can see that the correlation coefficients of $h(X^{1,i+1}, Y^{1,i+1})$ is equal to the sum of the $h(X^{1,i+1}, Y^{1,i+1})$'s correlation coefficients over the Aa-set with Out, In for all $Out \in F_2^d$. And the correlation coefficients of

$h(X^{1,i+1}, Y^{1,i+1})$ over the Aa-set with Out, In have the following recursive structure :

Theorem 4: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $\gamma, \lambda, v, w \in F_2^{(i+1) \cdot t}$, $Out, In, Middle \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Then

$$\begin{aligned} & Cor_{In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= \sum_{Out \in F_2^d} Cor_{Out, In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= Cor_{Out, In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= \sum_{Mi \in F_2^d} Cor_{Out, Mi, In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= \sum_{Mi \in F_2^d} Cor_{Out, Mi}^1 \left(\begin{array}{c} \gamma^2, \lambda^2, v^2, w^2, \\ T_j^i, D_j^i, G_j, \quad 1 \leq j \leq z \end{array} \right) \times Cor_{Mi, In}^i \left(\begin{array}{c} \gamma^1, \lambda^1, v^1, w^1, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \end{aligned}$$

where

$$\begin{aligned} \gamma^1 &= \gamma[0 : i \cdot t - 1], \quad \lambda^1 = \lambda[0 : i \cdot t - 1], \quad v^1 = v[0 : i \cdot t - 1], \quad w^1 = w[0 : i \cdot t - 1], \\ \gamma^2 &= \gamma[i \cdot t : (i+1) \cdot t - 1], \quad \lambda^2 = \lambda[i \cdot t : (i+1) \cdot t - 1], \quad v^2 = v[i \cdot t : (i+1) \cdot t - 1], \\ w^2 &= w[i \cdot t : (i+1) \cdot t - 1]. \end{aligned}$$

Proof. According to the **property 3** and **property 4**, we have:

$$\begin{aligned} & Cor_{In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= \frac{1}{2^{2(i+1) \cdot t}} \sum_{(X_{1,i+1}, Y_{1,i+1}) \in F_2^{(i+1)t}} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}} \\ &= \frac{1}{2^{2(i+1) \cdot t}} \sum_{(X_{1,i+1}, Y_{1,i+1}) \in \bigcup_{Out \in F_2^d} Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}} \\ &= \frac{1}{2^{2(i+1) \cdot t}} \sum_{Out \in F_2^d} \sum_{(X_{1,i+1}, Y_{1,i+1}) \in Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)} (-1)^{(\gamma, \lambda) \cdot h(X^{1,i+1}, Y^{1,i+1}) \oplus v \cdot X^{1,i+1} \oplus w \cdot Y^{1,i+1}} \\ &= \sum_{Out \in F_2^d} Cor_{In, Out}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \end{aligned}$$

Likely, from **lemma 1**, we can also get

$$\begin{aligned} & Cor_{In, Out}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \\ &= \sum_{Mi \in F_2^d} Cor_{In, Mi, Out}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1} \quad 1 \leq j \leq z \end{array} \right) \end{aligned}$$

Secondly, from **theorem 3**, for $1 \leq j \leq z$, we have

$$f(X^{1,i+1}, Y^{1,i+1})_{E_j, A_j^{i+1}, B_j^{i+1}} = f(X^{2,i}, Y^{2,i})_{M_j^i, T_j^i, D_j^i} \cdot 2^{i \cdot t} + f(X^{1,i}, Y^{1,i})_{E_j, A_j^i, B_j^i}.$$

Then

$$\begin{aligned}
h(X^{1,i+1}, Y^{1,i+1}) &= \bigoplus_{j=1}^z h(X^{1,i+1}, Y^{1,i+1})_{E_j, A_j^{i+1}, B_j^{i+1}} \\
&= \bigoplus_{j=1}^z f(X^{2,i}, Y^{2,i})_{M_j^i, T_j^i, D_j^i} \cdot 2^{i-t} + \bigoplus_{j=1}^z f(X^{1,i}, Y^{1,i})_{E_j, A_j^i, B_j^i} \\
&= h(X^{1,i}, Y^{1,i}) \cdot 2^{i-t} + h(X^{2,i}, Y^{2,i})
\end{aligned}$$

It can be concluded that:

$$\begin{aligned}
&Cor_{In, Mi, Out}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z \end{array} \right) \\
&= \frac{1}{2^{2i-t}} \sum_{(X^{1,i}, Y^{1,i}) \in Aa_{Mi, In}^i(A_j^i, B_j^i, 1 \leq j \leq z)} (-1)^{(\gamma^1, \lambda^1) \cdot h(X^{1,i}, Y^{1,i}) \oplus V^1 \cdot X^{1,i} \oplus W^1 \cdot Y^{1,i}} \\
&\times \frac{1}{2^{2t}} \sum_{(X^{2,i}, Y^{2,i}) \in Aa_{Out, Mi}^1(T_j^i, D_j^i, 1 \leq j \leq z)} (-1)^{(\gamma^2, \lambda^2) \cdot h(X^{2,i}, Y^{2,i}) \oplus V^2 \cdot X^{2,i} \oplus W^2 \cdot Y^{2,i}} \\
&= Cor_{Out, Mi}^1 \left(\begin{array}{c} \gamma^2, \lambda^2, v^2, w^2, \\ T_j^i, D_j^i, G_j, 1 \leq j \leq z \end{array} \right) \times Cor_{Mi, In}^i \left(\begin{array}{c} \gamma^1, \lambda^1, v^1, w^1, \\ A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z \end{array} \right)
\end{aligned}$$

Thus,

$$\begin{aligned}
&Cor_{Out, In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z \end{array} \right) \\
&= \sum_{Mi \in F_2^d} Cor_{Out, Mi, In}^{i+1} \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z \end{array} \right) \\
&= \sum_{Mi \in F_2^d} Cor_{Out, Mi}^1 \left(\begin{array}{c} \gamma^2, \lambda^2, v^2, w^2, \\ T_j^i, D_j^i, G_j, 1 \leq j \leq z \end{array} \right) \times Cor_{Mi, In}^i \left(\begin{array}{c} \gamma^1, \lambda^1, v^1, w^1, \\ A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z \end{array} \right)
\end{aligned}$$

□

Given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. For $0 \leq i \leq q-1$, supposed that matrix $Ma^i \in R^{d \times d}$, of which elements define as :

$$Ma^i[Out][In_1] = Cor_{Out, In_1}^1 \left(\begin{array}{c} \gamma_i^2, \lambda_i^2, v_i^2, w_i^2, \\ T_j^i, D_j^i, 1 \leq j \leq z \end{array} \right),$$

where $0 \leq Out, In_1 \leq 2^d - 1$, $\gamma, \lambda, v, w \in F_2^{q-t}$, $\gamma_i^2 = \gamma[i \cdot t : (i+1) \cdot t - 1]$, $\lambda_i^2 = \lambda[i \cdot t : (i+1) \cdot t - 1]$, $v_i^2 = v[i \cdot t : (i+1) \cdot t - 1]$, $w_i^2 = w[i \cdot t : (i+1) \cdot t - 1]$.

Then, we have:

Theorem 6: Given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$, where $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$. For $\gamma, \lambda, v, w \in F_2^{q-t}$, $In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Then,

$$Cor_{In}^q \left(\begin{array}{c} \gamma, \lambda, v, w, \\ A_j^q, B_j^q, G_j, 1 \leq j \leq z \end{array} \right) = L \prod_{i=0}^{q-1} Ma^i Q_{In}^T$$

where $1 \leq i \leq q-1$, $L = (1, 1, \dots, 1) \in F_2^d$, $Q_{In} \in F_2^d$ of which the sole nonzero component satisfies $Q[In] = 1$.

Proof. For $1 \leq i \leq q$, define the vector $Base_{In}^i \in F_2^d$ as follow:

$$Base_{In}^i[Out] = Cor_{Out, In}^i \left(\begin{array}{c} \gamma_i^1, \lambda_i^1, v_i^1, w_i^1, \\ A_j^i, B_j^i, 1 \leq j \leq z \end{array} \right)$$

where $0 \leq Out \leq 2^d - 1$, $\gamma^1 = \gamma[0 : i \cdot t - 1]$, $\lambda^1 = \lambda[0 : i \cdot t - 1]$, $v^1 = v[0 : i \cdot t - 1]$, $w^1 = w[0 : i \cdot t - 1]$.

According to the **lemma 1**, we have:

$$Cor_{In}^q \left(\begin{array}{c} \gamma, \lambda, V, W, \\ A_j^q, B_j^q, 1 \leq j \leq z \end{array} \right) = L \cdot Base_{In}^q$$

And by the definition, we see that $Base^1 = Ma^0 \cdot Q^T$.

In addition, according to the **lemma 1**, for $1 \leq i \leq q - 1$, $0 \leq Out \leq 2^d - 1$, we have

$$Base_{In}^{i+1}[Out] = \sum_{Out_1 \in F_2^d} Ma^i[Out][Out_1] \cdot Base_{In}^i[Out_1]$$

Namely,

$$Base_{In}^{i+1} = Ma^i \cdot Base_{In}^i$$

Thus, the theorem holds. \square

5 The Df -set

Recall the law of total probability, given a series of subset $\{A_i, 1 \leq i \leq n\}$ of the total space, where $\bigcup_{i=1}^n A_i$ is to equal the total space and each A_i are disjoint with each other, then for any event B , its probability $P(B)$ is equal to $\sum_{i=0}^n P(B \cap A_i)$. Notice that the the Aa-set with Out, In satisfy the above property, thus the same idea can be adopted to do the following study.

Definition 8: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q - 1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_m^{i+1}, B_m^{i+1}}; 1 \leq m \leq z\}$. Supposed that $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times (i+1) \cdot t}, F_2^{n \times (i+1) \cdot t})$, $Out, In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. We define the Df -set with In and the Df -set with Out, In respectively:

$$\begin{aligned} & Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ (X^{1,i+1}, Y^{1,i+1}) \left| \begin{array}{c} X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}, \\ h_s(X^{1,i+1}, Y^{1,i+1}) = C[s], 1 \leq s \leq n. \end{array} \right. \right\} \\ & Df_{Out, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ (X^{1,i+1}, Y^{1,i+1}) \left| \begin{array}{c} X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}, Out = (M_1^{i+1}, \dots, M_z^{i+1}), \\ h_s(X^{1,i+1}, Y^{1,i+1}) = C[s], 1 \leq s \leq n. \end{array} \right. \right\} \\ & Df_{Out, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)^* \\ &= (\{F_2^{(i+1) \cdot t} \times F_2^{(i+1) \cdot t}\} - Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)) \cap Aa_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \end{aligned}$$

where

$$h_s(X^{1,i+1}, Y^{1,i+1}) = \bigoplus_{m=1}^z G[s, m] * f(X^{1,i+1}, Y^{1,i+1})_{E_m, A_m^{i+1}, B_m^{i+1}}, 1 \leq s \leq n.$$

According to the **property 3** and **property 4** we have:

Property 5: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times (i+1) \cdot t}, F_2^{n \times (i+1) \cdot t})$, $Out, In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Then

$$\bigcup_{Out \in F_2^d} Df_{Out, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$$

According to the **Property 2:**, we have:

Property 6: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times (i+1) \cdot t}, F_2^{n \times (i+1) \cdot t})$, $Out, In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. If any two $Out_1, Out_2 \in F_2^d$ satisfy $Out_1 \neq Out_2$, then

$$Df_{Out_1, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Df_{Out_2, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

As same as the law of total probability, we can get the relation between the Df-set with In and the Df-set with Out , In :

Theorem 5: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times (i+1) \cdot t}, F_2^{n \times (i+1) \cdot t})$, $Out, In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Then

$$\begin{aligned} & \sum_{Out \in F_2^d} \#Df_{Out, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \#Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \end{aligned}$$

5.1 Direct method

According to the **lemma 1** and **theorem 3**, it can be concluded that the Df-set with Out , In have the following recursive structure:

Lemma 2: For $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$, $1 \leq i \leq q-1$, given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. Supposed that $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times (i+1) \cdot t}, F_2^{n \times (i+1) \cdot t})$, $Out, In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$, $J^1 = J[0 : i \cdot t - 1]$, $J^2 = J[i \cdot t : (i+1) \cdot t - 1]$, $N^1 = N[0 : i \cdot t - 1]$, $N^2 = N[i \cdot t : (i+1) \cdot t - 1]$, $C^1 = (J^1, N^1)$, $C^2 = (J^2, N^2)$. Define

$$\begin{aligned} & Df_{Out, Mi, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ (X^{2,i} || X^{1,i}, Y^{2,i} || Y^{1,i}) \left| \begin{array}{l} X^{1,i}, Y^{1,i} \in Df_{Mi, In}^{i, n}(C^1, G, A_j^i, B_j^i, 1 \leq j \leq z) \\ X^{2,i}, Y^{2,i} \in Df_{Out, Mi}^{1, n}(C^2, G, T_j^i, D_j^i, 1 \leq j \leq z) \end{array} \right. \right\} \end{aligned}$$

Then,

$$Df_{Out, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \bigcup_{Mi \in F_2^d} Df_{Out, Mi, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$$

holds. And for any two $Mi_1, Mi_2 \in F_2^d$, where $Mi_1 \neq Mi_2$, satisfy

$$Df_{Out, Mi_1, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Df_{Out, Mi_2, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

Proof. According to the **Theorem 3**, we have

$$f(X^{1,i+1}, Y^{1,i+1})_{E_j, A_j^{i+1}, B_j^{i+1}} = f(X^{2,i}, Y^{2,i})_{M_j^i, T_j^i, D_j^i} \cdot 2^{i \cdot t} + f(X^{1,i}, Y^{1,i})_{E_j, A_j^i, B_j^i}.$$

Thus, for $1 \leq s \leq n$:

$$\begin{aligned} h_s(X^{1,i+1}, Y^{1,i+1}) &= \bigoplus_{m=1}^z G[s, m] * f(X^{1,i+1}, Y^{1,i+1})_{E_m, A_m^{i+1}, B_m^{i+1}} \\ &= \bigoplus_{M=1}^z G[s, m] * f(X^{2,i}, Y^{2,i})_{M_m^i, T_m^i, D_m^i} \cdot 2^{i \cdot t} + \bigoplus_{m=1}^z G_j[s, m] * f(X^{1,i}, Y^{1,i})_{E_m, A_m^i, B_m^i} \\ &= h_s(X^{1,i}, Y^{1,i}) \cdot 2^{i \cdot t} + h_s(X^{2,i}, Y^{2,i}) \end{aligned}$$

It means that:

$$\begin{aligned} &Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ (X^{2,i} || X^{1,i}, Y^{2,i} || Y^{1,i}) \left| \begin{array}{l} X^{1,i}, Y^{1,i} \in F_2^{i \cdot t}, X^{2,i}, Y^{2,i} \in F_2^t, Mi \in F_2^d, Mi = (M_1^{i+1}, \dots, M_z^{i+1}), \\ h_s(X^{1,i}, Y^{1,i}) = C_i^1[s], \\ h_s(X^{2,i}, Y^{2,i}) = C_i^2[s], 1 \leq s \leq n. \end{array} \right. \right\} \\ &= \left\{ (X^{2,i} || X^{1,i}, Y^{2,i} || Y^{1,i}) \left| \begin{array}{l} X^{1,i}, Y^{1,i} \in Df_{Mi}^{i, n}(C^1, G, A_j^i, B_j^i, 1 \leq j \leq z) \\ Mi \in F_2^d, Mi = (M_1^{i+1}, \dots, M_z^{i+1}); \\ X^{2,i}, Y^{2,i} \in Df_{In}^{1, n}(C^2, G, T_j^i, D_j^i, 1 \leq j \leq z) \end{array} \right. \right\} \end{aligned}$$

Then,

$$\begin{aligned} &Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap A_{Out, Mi, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= \left\{ \left(\begin{array}{l} X^{2,i} || X^{1,i} \\ Y^{2,i} || Y^{1,i} \end{array} \right) \left| \begin{array}{l} Mi = (M_1^{i+1}, \dots, M_z^{i+1}), \\ X^{1,i}, Y^{1,i} \in A_{Mi, In}^i(A_j^i, B_j^i, 1 \leq j \leq z) \cap Df_{Mi}^{i, n}(C^1, G, A_j^i, B_j^i, 1 \leq j \leq z) \\ X^{2,i}, Y^{2,i} \in A_{Out, Mi}^i(T_j^i, D_j^i, 1 \leq j \leq z) \cap Df_{In}^{1, n}(C^2, G, T_j^i, D_j^i, 1 \leq j \leq z) \end{array} \right. \right\} \end{aligned}$$

It can be concluded that

$$\begin{aligned} &Df_{In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap A_{Out, Mi, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\ &= Df_{Out, Mi, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \end{aligned}$$

From Lemma 1, for any two $Mi_1, Mi_2 \in F_2^d$, where $Mi_1 \neq Mi_2$, satisfy

$$A_{Out, Mi_1, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap A_{Out, Mi_2, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

Thus,

$$Df_{Out, Mi_1, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \cap Df_{Out, Mi_2, In}^{i+1, n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \emptyset$$

Secondly, due to

$$A_{Out, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) = \bigcup_{Mi \in F_2^d} A_{Out, Mi, In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$$

holds in Lemma 1, we can get:

$$\begin{aligned}
& Df_{Out,In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= Df_{In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \bigcap A_{Out,In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= Df_{In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \bigcap \bigcup_{Mi \in F_2^d} A_{Out,Mi,In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= \bigcup_{Mi \in F_2^d} Df_{In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \bigcap A_{Out,Mi,In}^{i+1}(A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= \bigcup_{Mi \in F_2^d} Df_{Out,Mi,In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)
\end{aligned}$$

□

According to the above recursive structure, we can calculate the order of the Df-set with Out, In based on the following recurrence relation:

Corollary 2: For any positive integer q, t, z, n and $0 \leq i \leq q-1$, then

$$\begin{aligned}
& \#Df_{Out,In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) \\
&= \sum_{Mi \in F_2^d} \#Df_{Out,Mi}^{1,n}(C_i^2, G, T_j^i, D_j^i, 1 \leq j \leq z) \\
&\quad \times \#Df_{Mi,In}^{i,n}(C_i^1, G, A_j^i, B_j^i, 1 \leq j \leq z).
\end{aligned}$$

holds.

Given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. For $0 \leq i \leq q-1$, supposed that matrix $Md^i \in R^{d \times d}$, of which elements define as :

$$Md^i[Out][In_1] = \frac{1}{2^{2t}} \#Df_{Out,In_1}^{1,n}(C_i^2, G, T_j^i, D_j^i, 1 \leq j \leq z)$$

where $0 \leq Out, In_1 \leq 2^d - 1$, $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times q \cdot t}, F_2^{n \times q \cdot t})$, $J_i^2 = J[i \cdot t : (i+1) \cdot t - 1]$, $N_i^2 = N[i \cdot t : (i+1) \cdot t - 1]$, $C_i^2 = (J_i^2, N_i^2)$. Then, we have:

Theorem 7: Given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$, where $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$. For $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times q \cdot t}, F_2^{n \times q \cdot t})$, $In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$.

Let $C = (J, N) = C_q^1 = (J_q^1, N_q^1)$, then,

$$\frac{1}{2^{2qt}} \#Df_{In}^{q,n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z) = L \prod_{i=0}^{q-1} Md^i Q_{In}^T$$

where $1 \leq i \leq q-1$, $L = (1, 1, \dots, 1) \in F_2^d$, $Q_{In} \in F_2^d$ of which the sole nonzero component satisfies $Q[In] = 1$.

Proof. For $1 \leq i \leq q$, define vector $Num_{In}^i \in F_2^d$ as follow:

$$Num_{In}^i[Out] = \frac{1}{2^{2it}} Df_{Out,In}^{i,n}(C_i^1, G_0, G_j, A_j^i, B_j^i, 1 \leq j \leq z)$$

where $0 \leq Out \leq 2^d - 1$.

According to the **theorem 5** we have:

$$\frac{1}{2^{2qt}} Df_{In}^{q,n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z) = L \cdot Num_{In}^q$$

And by the definition, we see that $Num_{In}^1 = Md^0 \cdot Q_{In}^T$.

In addition, according to the **corollary 2**, for $1 \leq i \leq q-1$, $0 \leq Out \leq 2^d - 1$, we have

$$Num_{In}^{i+1}[Out] = \sum_{Out_1 \in F_2^d} Md^i[Out][Out_1] \cdot Num_{In}^i[Out_1]$$

Namely,

$$Num_{In}^{i+1} = Md^i \cdot Num_{In}^i$$

Thus, the theorem holds. \square

5.2 Indirect method

In this part, we will use the Markov chain to complete the proof of the **Theorem 7**:

Theorem 7: Given any z basic function $\{f(X^{1,i+1}, Y^{1,i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$, where $X^{1,i+1}, Y^{1,i+1} \in F_2^{(i+1) \cdot t}$. For $G \in F_2^{n \times z}$, $C = (J, N) \in (F_2^{n \times q \cdot t}, F_2^{n \times q \cdot t})$, $In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$.

Let $C = (J, N) = C_q^1 = (J_q^1, N_q^1)$. If $X^{1,i+1}, Y^{1,i+1}$ are chosen uniformly at random in $F_2^{(i+1) \cdot t}$, then,

$$\begin{aligned} Pr((X^{1,i+1}, Y^{1,i+1}) \in Df_{In}^{q,n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z)) \\ = \frac{1}{2^{2qt}} \# Df_{In}^{q,n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z) = L \prod_{i=0}^{q-1} Md^i Q_{In}^T \end{aligned}$$

where $1 \leq i \leq q-1$, $L = (1, 1, \dots, 1) \in F_2^d$, $Q_{In} \in F_2^d$ of which the sole nonzero component satisfies $Q[In] = 1$.

Proof. Supposed that $(X^{1,i+1}, Y^{1,i+1}) \in Df_{Out, In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$ represents the state $T_{Out, In}^{i+1,n}(C_{i+1}^1, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$;

$(X^{1,i+1}, Y^{1,i+1}) \in Df_{Out, In}^{i+1,n}(C, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)^*$ represents the state $F_{Out, In}^{i+1,n}(C_{i+1}^1, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z)$.

Thus, we can find that any two of the above states are disjoint and the union of all the above state is Ω according to the Property 5-6.

Secondly, for $1 \leq s \leq n$, due to

$$\begin{aligned} h_s(X^{1,i+1}, Y^{1,i+1}) &= \bigoplus_{m=1}^z G[s, m] * f(X^{1,i+1}, Y^{1,i+1})_{E_m, A_m^{i+1}, B_m^{i+1}} \\ &= \bigoplus_{M=1}^z G[s, m] * f(X^{2,i}, Y^{2,i})_{M_m^i, T_m^i, D_m^i} \cdot 2^{i \cdot t} + \bigoplus_{m=1}^z G_j[s, m] * f(X^{1,i}, Y^{1,i})_{E_m, A_m^i, B_m^i} \\ &= h_s(X^{1,i}, Y^{1,i}) \cdot 2^{i \cdot t} + h_s(X^{2,i}, Y^{2,i}) \end{aligned}$$

If we know the value of $X^{1,i}, Y^{1,i}$, then the probability of which state $(X^{1,i+1}, Y^{1,i+1})$ belongs to is depend on the state that $X^{1,i}, Y^{1,i}$ belongs to. Thus, we can use it to build a Markov chain.

For $1 \leq i \leq q-1$, $O_1, O_2 \in F_2^d$:

$$Pr(T_{O_1, In}^{i+1,n}(C_{i+1}^1, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) | T_{O_2, In}^{i,n}(C_i^1, G, A_j^i, B_j^i, 1 \leq j \leq z)) = Md^i[O_1, O_2]$$

$$Pr(T_{O_1, In}^{i+1,n}(C_{i+1}^1, G, A_j^{i+1}, B_j^{i+1}, 1 \leq j \leq z) | F_{O_2, In}^{i,n}(C_i^1, G, A_j^i, B_j^i, 1 \leq j \leq z)) = 0$$

Thus, the one step transform matrix form i to $i + 1$ is:

$$\begin{Bmatrix} Md^i, O \\ *, * \end{Bmatrix}$$

Then, the $q - 1$ step transform matrix is

$$\begin{Bmatrix} \prod_{i=1}^{q-1} Md^i, O \\ *, * \end{Bmatrix}$$

Define vector $V_{In} \in F_2^d$ as follow:

$$V_{In}[Out] = Df_{Out, In}^{1, n}(C_1^1, G_0, G_j, A_j^1, B_j^1, 1 \leq j \leq z)$$

where $0 \leq Out \leq 2^d - 1$. Then $V_{In} = Md^0 * Q_{In}^T$. According to the property of Markov chain, we get

$$Pr((X^{1, i+1}, Y^{1, i+1}) \in Df_{In}^{q, n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z)) = L \prod_{i=1}^{q-1} Md^i V_{In} = L \prod_{i=0}^{q-1} Md^i Q_{In}^T$$

□

Likely, if we use the Markov chain, we get another form of **Theorem 6**:

Given any z basic function $\{f(X^{1, i+1}, Y^{1, i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$. For $0 \leq i \leq q-1$, supposed that matrix $M0^i, M1^i \in R^{d \times d}$, of which elements define as :

$$\begin{aligned} M0^i[Out][In_1] &= \frac{\#Aa_{Out, In_1}^1(T_j^i, D_j^i, 1 \leq j \leq z)}{2^{t+1}} + \frac{1}{2} Cor_{Out, In_1}^1 \left(\begin{matrix} \gamma_i^2, \lambda_i^2, V_i^2, W_i^2 \\ T_j^i, D_j^i, 1 \leq j \leq z \end{matrix} \right) \\ M1^i[Out][In_1] &= \frac{\#Aa_{Out, In_1}^1(T_j^i, D_j^i, 1 \leq j \leq z)}{2^{t+1}} - \frac{1}{2} Cor_{Out, In_1}^1 \left(\begin{matrix} \gamma_i^2, \lambda_i^2, V_i^2, W_i^2 \\ T_j^i, D_j^i, 1 \leq j \leq z \end{matrix} \right) \end{aligned}$$

where $0 \leq Out, In_1 \leq 2^d - 1$, $\gamma, \lambda, v, w \in F_2^{q \cdot t}$, $\gamma_i^2 = \gamma[i \cdot t : (i+1) \cdot t - 1]$, $\lambda_i^2 = \lambda[i \cdot t : (i+1) \cdot t - 1]$, $v_i^2 = v[i \cdot t : (i+1) \cdot t - 1]$, $w_i^2 = w[i \cdot t : (i+1) \cdot t - 1]$.

Then, for $0 \leq i \leq q-1$, define the matrix M_i as follow:

$$\begin{Bmatrix} M0^i, M1^i \\ M1^i, M0^i \end{Bmatrix}$$

Theorem 8: Given any z basic function $\{f(X^{1, i+1}, Y^{1, i+1})_{E_{k_m}, A_{k_m}^{i+1}, B_{k_m}^{i+1}}; 1 \leq m \leq z\}$,

where $X^{1, i+1}, Y^{1, i+1} \in F_2^{(i+1) \cdot t}$. For $\gamma, \lambda, v, w \in F_2^{q \cdot t}$, $In \in F_2^d$, where $d = \sum_{i=0}^z k_i$, $In = (E_1, \dots, E_z)$. Then,

$$Cor_{In}^q \left(\begin{matrix} \gamma, \lambda, v, w \\ A_j^q, B_j^q, G_j, 1 \leq j \leq z \end{matrix} \right) = (L, O) \cdot \prod_{i=0}^{q-1} M_i \cdot (Q_{In}, Q_{In})^T$$

where $1 \leq i \leq q-1$, $L = (1, 1, \dots, 1)$, $O = (0, 0, \dots, 0) \in F_2^d$, $Q_{In} \in F_2^d$ of which the sole nonzero component satisfies $Q[In] = 1$.

Remark 7: Supposed that $m = q$ and $t = 1$ in **theorem 6** and **theorem 7**, if $2^m \gg 2^d$ and $m \gg 2^{d^2}$, then the time complex of calculating the $Cor_{In}^q \left(\begin{matrix} \gamma, \lambda, V, W \\ A_j^q, B_j^q, 1 \leq j \leq z \end{matrix} \right)$ and $Df_{In}^{m, n}(C, G, A_j^q, B_j^q, 1 \leq j \leq z)$ is about $O(m)$.

5.3 Instance

For $F : (x, y) \xrightarrow{F} (x, x \boxplus y)$, it can be treated as 1-order basic function. Besides this, its inverse function F^{-1} is $(x, y) \xrightarrow{F^{-1}} (x, x \boxminus y)$. According to the **corollary 1**, it can be conversed into $(x, y) \xrightarrow{F^{-1}} (x, (x \oplus (1, 1, \dots, 1)) \boxplus y \boxplus 1)$. For $\alpha, \beta, \gamma, \lambda \in F_2^n$, let $E = [1, 0]$, $B_1 = [0, \lambda]$, $A_1 = [0, \gamma]$, $B_2 = [\beta, \lambda]$, $A_2 = [\alpha, \gamma]$. Then, for 2-order basic function $f(x, y)_{E, A_1, B_1}$ and $f(x, y)_{E, A_2, B_2}$, the following equation holds.

$$\begin{aligned} f(x, y)_{E, A_1, B_1} &= F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \\ f(x, y)_{E, A_2, B_2} &= F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) \end{aligned}$$

Thus, according to the **theorem 6** and the **theorem 7**, we have:

1. The formula for calculating the boomerang connective probability and its variant :

Corollary 3(BCT): Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of BCT [4] defined by

$$BCT(\alpha, \beta, \gamma, \lambda) = \#\{(x, y) \mid x, y \in F_2^n, F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\alpha, \beta)\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda \in F_2^{2n}$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i])$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$\begin{aligned} &M_{d[i]}[Out][In] \\ &= Df_{Out, In}^{1,1}((\alpha[i], \beta[i]), (0, 0), \{(0, 1), (0, \overline{\gamma[i]})\}, \{(0, 1), (\alpha[i], \overline{\gamma[i]})\}, (\beta[i], \lambda[i])\}) \\ &= \# \left\{ (x, y) \left| \begin{array}{l} e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \lambda[i], x \oplus 1 \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ &\text{where } 0 \leq Out, In \leq 15. \end{aligned}$$

Thus, we have

$$BCT(\alpha, \beta, \gamma, \lambda) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

Corollary 4 (BCT¹): Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of BCT¹ [10] defined as

$$BCT^1(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = \#\{(x, y) \mid x, y \in F_2^n, F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\theta, \zeta)\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta \in F_2^{2n}$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i]) \in F_2^6$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$\begin{aligned} &M_{d[i]}[Out][In] \\ &= Df_{Out, In}^{1,1}((\theta[i], \zeta[i]), (0, 0), \{(0, 1), (0, \overline{\gamma[i]})\}, \{(0, 1), (\alpha[i], \overline{\gamma[i]})\}, (\beta[i], \lambda[i])) \\ &= \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] = \theta[i], \zeta[i] \oplus \beta[i] \oplus e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \lambda[i], x \oplus 1 \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ &\text{where } 0 \leq Out, In \leq 15. \end{aligned}$$

Thus, we have:

$$BCT^1(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

Corollary 5(BCT^2): Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of BCT^2 [10] defined as

$$BCT^2(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = \#\{(x, y) \mid x, y \in F_2^n, F^{-1}(F(x, y) \oplus (\theta, \zeta)) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\alpha, \beta)\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta \in F_2^{2n}$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i]) \in F_2^6$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$\begin{aligned} & M_{d[i]}[Out][In] \\ &= Df_{Out, In}^{1,1}((\alpha[i], \beta[i]), (0, 0), \{(0, 1), (0, \overline{\theta[i]})\}, \{(0, 1), (\alpha[i], \overline{\gamma[i]})\}, (\beta[i], \lambda[i])) \\ &= \# \left\{ (x, y) \left| \begin{array}{l} \lambda[i] \oplus \theta[i] \oplus \gamma[i] \oplus \zeta[i] \oplus e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \\ \theta[i] \oplus \gamma[i] = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \zeta[i], x \oplus 1 \oplus \theta[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ & \text{where } 0 \leq Out, In \leq 15. \end{aligned}$$

Thus, we have:

$$BCT^1(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

3. The formula for calculating the difference-boomerange connective probability and the inverse difference-boomerange probability, respectively:

Corollary 6-1(**DBT**) Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of DBT [5] defined by

$$DBT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = \# \left\{ (x, y) \left| \begin{array}{l} x, y \in F_2^n, F(x, y) \oplus F(x \oplus \alpha, y \oplus \beta) = (\theta, \zeta), \\ F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\alpha, \beta). \end{array} \right. \right\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta \in F_2^{2n}$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i]) \in F_2^6$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$\begin{aligned} & M_{d[i]}[Out][In] \\ &= Df_{Out, In}^{1,1} \left(\begin{pmatrix} \alpha[i] & \beta[i] \\ \theta[i] & \zeta[i] \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \left\{ \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, (0, \overline{\gamma[i]}), (0, \lambda[i]) \right\}, \left\{ \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, (\alpha[i], \overline{\gamma[i]}), (\beta[i], \lambda[i]) \right\} \right) \\ &= \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] = \theta[i], \alpha[i] \oplus \beta[i] \oplus \zeta[i] \oplus e_1 \oplus e_3 = 0, \\ e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \lambda[i], x \oplus 1 \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ & \text{where } 0 \leq Out, In \leq 15. \end{aligned}$$

Thus, we have

$$DBT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

Corollary 6-2:(IDBT) Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of DBT [5] defined by

$$DBT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = \# \left\{ (x, y) \left| \begin{array}{l} x, y \in F_2^n, F^{-1}(x, y) \oplus F^{-1}(x \oplus \alpha, y \oplus \beta) = (\theta, \zeta), \\ F(F^{-1}(x, y) \oplus (\gamma, \lambda)) \oplus F(F^{-1}(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\alpha, \beta). \end{array} \right. \right\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i]) \in F_2^6$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$\begin{aligned} & M_{d[i]}[Out][In] \\ = & Df_{Out, In}^{1,1} \left(\begin{pmatrix} \alpha[i] \\ \theta[i] \end{pmatrix}, \begin{pmatrix} \beta[i] \\ \zeta[i] \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \left\{ \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, (1, \gamma[i]), (0, \lambda[i]) \right\}, \left\{ \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, (\overline{\alpha[i]}, \gamma[i]), (\beta[i], \lambda[i]) \right\} \right) \\ = & \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] = \theta[i], \alpha[i] \oplus \beta[i] \oplus \zeta[i] \oplus e_1 \oplus e_3 = 0, \\ e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x \oplus 1)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus 1 \oplus e_1 \oplus \lambda[i], x \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus 1 \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus 1 \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ & \text{where } 0 \leq Out, In \leq 15. \end{aligned}$$

Then, we have:

$$DBT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

3. The formula for calculating the difference probability :

Corollary 7:(DDT) Let S be $S(x, y) = x \boxplus y$, an element of DDT [1] defined by

$$DDT(\alpha, \beta, \Delta) = \#\{(x, y) \mid x, y \in F_2^n, S(x \oplus \alpha, y \oplus \beta) \oplus S(x, y) = \Delta\} \cdot 2^{-2n}$$

where $\alpha, \beta, \Delta \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_3) \in F_2^2$, $In = (e_1, e_3), L = (1, 1, 1, 1), Q = (1, 0, 0, 0) \in F_2^4$, $d[i] = (\alpha[i], \beta[i], \Delta[i]) \in F_2^3$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{4 \times 4}$, of which elements define as :

$$\begin{aligned} & M_{d[i]}[Out, In] \\ = & Df_{Out, In}^{1,1} \left((\alpha[i], \Delta[i]), (0, 0), \{(1), (0), (0)\}, \{(1), (\alpha[i]), (\beta[i])\} \right) \\ = & \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] \oplus \beta[i] \oplus \Delta[i] \oplus e_1 \oplus e_3 = 0, \\ \text{carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\} \\ & \text{where } 0 \leq Out, In \leq 3. \end{aligned}$$

Thus, we have:

$$DDT(\alpha, \beta, \Delta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

4. The formula for calculating difference-linear connective correlation coefficients:

Corollary 8:(DLCT) Let S be a $S(x, y) = x \boxplus y$, an element of DLCT [4] defined by

$$DLCT(\alpha, \beta, \lambda) = 2^{-2n} \cdot \sum_{x, y \in F_2^n} (-1)^{\lambda \cdot (S(x \oplus \alpha, y \oplus \beta) \oplus S(x, y))}$$

where $\alpha, \beta, \lambda \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_3) \in F_2^2$, $In = (e_1, e_3), L = (1, 1, 1, 1), Q = (1, 0, 0, 0) \in F_2^4$, $a[i] = (\alpha[i], \beta[i], \lambda[i]) \in F_2^3$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{a[i]} \in F_2^{4 \times 4}$, of which elements define as :

$$\begin{aligned} & M_{a[i]}[Out][In] \\ &= Cor_{Out, In}^1 \left(\begin{array}{c} 0, \lambda[i], 0, 0, \\ \{(1), (0), (0)\}, \{(1), (\alpha[i]), (\beta[i])\} \end{array} \right) \\ &= \sum_{x, y \in Set_{a[i], Out, In}} (-1)^{\lambda[i] \cdot (S(x \oplus \alpha[i], y \oplus \beta[i]) \oplus S(x, y))} \\ & \text{where } Set_{a[i], Out, In} = \left\{ (x, y) \left| \begin{array}{l} carry_{e_1}(y, x)[1] = o_1, \\ carry_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{where } carry_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\}, \\ & 0 \leq Out, In \leq 3. \end{aligned}$$

Thus, we have:

$$DLCT(\alpha, \beta, \lambda) = L \prod_{i=0}^{n-1} M_{a[i]} Q^T$$

5. The formula for calculating the linear approximation correlation coefficients(LAT):

Corollary 9:(LAT) Let S be $S(x, y) = x \boxplus y$, an element of LAT [2] defined by

$$LAT(\mu, \omega, \lambda) = 2^{-2n} \cdot \sum_{x, y \in F_2^n} (-1)^{\mu \cdot x \oplus \omega \cdot y \oplus \lambda \cdot S(x, y)}$$

where $\mu, \omega, \lambda \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out \in F_2, In \in F_2, L = (1, 1), Q = (1, 0), a[i] = (\mu[i], \omega[i], \lambda[i]) \in F_2^3$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{a[i]} \in F_2^{2 \times 2}$, of which elements define as :

$$\begin{aligned} & M_{a[i]}[Out][In] = Cor_{Out, In}^1 \left(\begin{array}{c} 0, \lambda[i], \mu[i], \omega[i], \\ \{(1), (0), (0)\}, \end{array} \right) \\ &= \sum_{x, y \in Set_{a[i], Out, In}} (-1)^{\mu[i] \cdot x \oplus \omega[i] \cdot y \oplus \lambda[i] \cdot S(x, y)} \\ & \text{where } Set_{a[i], Out, In} = \left\{ (x, y) \left| \begin{array}{l} carry_{In}(y, x)[1] = Out, \\ \text{where } carry_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\}, \\ & 0 \leq Out, In \leq 1. \end{aligned}$$

Thus, we have:

$$LAT(\mu, \omega, \lambda) = L \prod_{i=0}^{n-1} M_{a[i]} Q^T$$

According to the **Theorem 8** in Appendix-A, after reducing the redundancy of matrix, we have the formulas for the calculating the boomerang-difference connective probability and the variant of difference-boomerang connective probability, respectively:

Corollary 10(BDT): Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of BDT [10] defined as

$$BDT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = \# \left\{ (x, y) \left| \begin{array}{l} x, y \in F_2^n, (x, y) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta)) = (\theta, \zeta) \\ F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \oplus F^{-1}(F(x \oplus \alpha, y \oplus \beta) \oplus (\gamma, \lambda)) = (\alpha, \beta). \end{array} \right. \right\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4, o_5) \in F_2^5$, $In = (e_1, e_2, e_3, e_4, e_5) \in F_2^5$, $L = (1, 1, \dots, 1) \in F_2^{32}$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i]) \in$

F_2^6 , $Q = (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, \dots, 0) \in F_2^{32}$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{32 \times 32}$, of which elements define as :

$$M_{d[i]}[Out][In] = \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] = \theta[i], 1 \oplus \beta[i] \oplus \zeta[i] \oplus e_5 \oplus e_3 = 0, \\ e_4 \oplus e_3 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \lambda[i], x \oplus 1 \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3 \oplus \lambda[i], x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_4; \\ \text{carry}_{e_5}(y \oplus \beta[i] \oplus x \oplus \alpha[i] \oplus e_3, x \oplus 1 \oplus \alpha[i] \oplus \gamma[i])[1] = o_5; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\}$$

where $0 \leq Out, In \leq 31$.

Thus, we have

$$BDT(\alpha, \beta, \gamma, \lambda, \theta, \zeta) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

Corollary 11(DBT¹): Let F be $(x, y) \xrightarrow{F} (x, x \boxplus y)$, an element of DBT^1 [10] defined as

$$DBT^1(\alpha, \beta, \gamma, \lambda, \theta, \zeta, \eta, \psi) = \# \left\{ (x, y) \left|_{F^{-1}(F(x, y) \oplus (\gamma, \lambda)) \oplus F^{-1}(F(x, y) \oplus (\gamma, \lambda) \oplus (\theta, \zeta)) = (\eta, \psi)} \begin{array}{l} x, y \in F_2^n, F(x, y) \oplus F(x \oplus \alpha, y \oplus \beta) = (\theta, \zeta), \end{array} \right. \right\} \cdot 2^{-2n}$$

where $\alpha, \beta, \gamma, \lambda, \theta, \zeta, \eta, \psi \in F_2^n$. Then, for $0 \leq i \leq n-1$, let $Out = (o_1, o_2, o_3, o_4) \in F_2^4$, $In = (e_1, e_2, e_3, e_4) \in F_2^4$, $L = (1, 1, \dots, 1) \in F_2^{16}$, $Q = (0, 0, 0, 0, 0, 1, 0, \dots, 0) \in F_2^{16}$, $d[i] = (\alpha[i], \beta[i], \gamma[i], \lambda[i], \theta[i], \zeta[i], \eta[i], \psi[i]) \in F_2^8$. Beside this, for $0 \leq i \leq n-1$, supposed that the matrix $M_{d[i]} \in F_2^{16 \times 16}$, of which elements define as :

$$M_{d[i]}[Out][In] = \# \left\{ (x, y) \left| \begin{array}{l} \alpha[i] = \theta[i], \alpha[i] \oplus \beta[i] \oplus \zeta[i] \oplus e_1 \oplus e_3 = 0, \eta[i] = \theta[i], \\ \psi[i] \oplus \theta[i] \oplus \zeta[i] \oplus e_4 \oplus e_1 \oplus e_2 = 0, \text{ carry}_{e_1}(y, x)[1] = o_1, \\ \text{carry}_{e_2}(y \oplus x \oplus e_1 \oplus \lambda[i], x \oplus 1 \oplus \gamma[i])[1] = o_2, \\ \text{carry}_{e_3}(y \oplus \beta[i], x \oplus \alpha[i])[1] = o_3, \\ \text{carry}_{e_4}(y \oplus x \oplus e_1 \oplus \lambda[i] \oplus \zeta[i], x \oplus 1 \oplus \gamma[i] \oplus \theta[i])[1] = o_4; \\ \text{where } \text{carry}_e(x, y)[1] = (x \wedge y) \oplus (x \wedge e) \oplus (e \wedge y); x, y, \in F^2. \end{array} \right. \right\}$$

where $0 \leq Out, In \leq 15$.

Thus, we have:

$$DBT^1(\alpha, \beta, \gamma, \lambda, \theta, \zeta, \eta, \psi) = 2^{-2n} \cdot L \prod_{i=0}^{n-1} M_{d[i]} Q^T$$

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A Reduce The Redundancy of Matrix

Definition 9: For any 2 basic function series $\{f(X, Y)_{E_{k_{m,1}}, A_{k_{m,1}}^{i+1}, B_{k_{m,1}}^{i+1}}; 1 \leq m \leq r_1\}_1$ and $\{f(X, Y)_{E_{k_{m,2}}, A_{k_{m,2}}^{i+1}, B_{k_{m,2}}^{i+1}}; 1 \leq m \leq r_2\}_2$, if there exist a k_1 , such that $k_1 \geq 0$, $E_{k_{r_1,1}}[0 : k_1] = E_{k_{r_2,2}}[0 : k_1]$, $A_{k_{r_1,1}}^{i+1}[0 : k_1] = A_{k_{r_2,2}}^{i+1}[0 : k_1]$, $B_{k_{r_1,1}}^{i+1}[0 : k_1] = B_{k_{r_2,2}}^{i+1}[0 : k_1]$; then we called that the two basic function series are similar. And the degree of similarity *deg* for the two basic function series is defined as

$$deg = \max \{k + 1 | k \geq 0, E_{k_{m,1}}[0 : k] = E_{k_{m,2}}[0 : k], A_{k_{m,1}}^{i+1}[0 : k] = A_{k_{m,2}}^{i+1}[0 : k], B_{k_{m,1}}^{i+1}[0 : k] = B_{k_{m,2}}^{i+1}[0 : k]\}.$$

Beside this, if two basic function series are not similar, we define $deg = 0$.

Definition 10: Given any z number r_1, \dots, r_z , supposed that $1 \leq j_1 \leq j_2 \leq z$, $deg \leq r_{j_1}, r_{j_2}$, where $d = \sum_{i=1}^z r_i$, then we can define

$$Sim_{deg}^{j_1, j_2} = \{V \in F_2^d | V = (V_1, V_2, \dots, V_z), V_{r_{j_1}}[0 : d] = V_{r_{j_2}}[0 : d], V_i \in F_2^{r_i}, 1 \leq i \leq z\}$$

Define the bijection $F : Sim_{deg}^{j_1, j_2} \xrightarrow{F} F_2^{d-deg}$ as:
 $F(V_1, V_2, \dots, V_{j_1}, \dots, V_{j_2}, \dots, V_z) = (V_1, V_2, \dots, V_{j_1}, \dots, V_{j_2}[deg : r_{j_2} - 1], \dots, V_z)$.

Theorem 8: For any positive integer q, t, z, n , given any z basic function series $\{f(X^{1,q}, Y^{1,q})_{E_{k_{m,j}}, A_{k_{m,j}}^q, B_{k_{m,j}}^q}; 1 \leq m \leq r_j\}_j (X^{1,q}, Y^{1,q} \in F_2^{q \cdot t}), 1 \leq j \leq z$; then we can get $T_{k_{r_j, j}}^i, D_{k_{r_j, j}}^i \in (F_2^t)^{r_j}$ from $A_{k_{m, j}}^q, B_{k_{m, j}}^q$, where $1 \leq j \leq z, 0 \leq i \leq q-1$. And there are two basic function series, of which code are j_1, j_2 respectively, are similar. Then we define the degree of similarity for the two basic function series is *deg*. Supposed that $G_0 \in F_2^{n \times 2}$, $G_j \in F_2^{n \times r_j}, 1 \leq j \leq z$; $C = (J, N) \in (F_2^{n \times q \cdot t}, F_2^{n \times q \cdot t})$, $Out, In_1 \in F_2^d$, $In \in Sim_{deg}^{j_1, j_2}$, where $d = \sum_{i=0}^z k_{r_i}$, $In = (E_{k_{r_1,1}}, \dots, E_{k_{r_z,z}})$. Let $C = (J, N) = C_q^1 = (J_q^1, N_q^1)$. We have:

$$Df_{In}^{q,n}(C, G_0, G_j, A_{k_{r_j, j}}^q, B_{k_{r_j, j}}^q, 1 \leq j \leq z) = L \prod_{i=0}^{q-1} Md^i Q^T$$

where $L = (1, 1, \dots, 1) \in F_2^{d-deg}$, $Q \in F_2^{d-deg}$, of which the sole nonzero component satisfies $Q[F(In)] = 1$, and $Md^i \in R^{(d-deg) \times (d-deg)}$ satisfying $Md^i[out, in] =$

$$\#Df_{F^{-1}(out), F^{-1}(in)}^{1,n}(C_i^2, G_0, G_j, T_{k_{r_j,j}}^i, D_{k_{r_j,j}}^i, 1 \leq j \leq z), 0 \leq i \leq q-1, 0 \leq out, in \leq 2^{d-deg} - 1.$$

Proof. According to the **remark 5**, for $1 \leq i \leq q$, we have :

$$Df_{Out, In}^{i,n}(C_i^1, G_0, G_j, A_{k_{r_j,j}}^i, B_{k_{r_j,j}}^i, 1 \leq j \leq z) = 0, \text{ when } in \notin Sim_{deg}^{j_1, j_2} \text{ and } out \in Sim_{deg}^{j_1, j_2}.$$

$$Df_{Out, In}^{i,n}(C_i^1, G_0, G_j, A_{k_{r_j,j}}^i, B_{k_{r_j,j}}^i, 1 \leq j \leq z) = 0, \text{ when } in \in Sim_{deg}^{j_1, j_2} \text{ and } out \notin Sim_{deg}^{j_1, j_2}.$$

Thus, for $1 \leq i \leq q$, when $In \in Sim_{deg}^{j_1, j_2}$

$$\begin{aligned} & \sum_{Out \in Sim_{deg}^{j_1, j_2}} \#Df_{Out, In}^{i,n}(C_{i+1}^1, G_0, G_j, A_{k_{r_j,j}}^{i+1}, B_{k_{r_j,j}}^{i+1}, 1 \leq j \leq z) \\ &= \#Df_{In}^{i,n}(C_{i+1}^1, G_0, G_j, A_{k_{r_j,j}}^{i+1}, B_{k_{r_j,j}}^{i+1}, 1 \leq j \leq z) \end{aligned}$$

And for $1 \leq i \leq q-1$, when $In \in Sim_{deg}^{j_1, j_2}$,

$$\begin{aligned} & \#Df_{Out, In}^{i+1,n}(C_{i+1}^1, G_0, G_j, A_{k_{r_j,j}}^{i+1}, B_{k_{r_j,j}}^{i+1}, 1 \leq j \leq z) \\ &= \sum_{Mi \in Sim_{deg}^{j_1, j_2}} \#Df_{Out, Mi}^{1,n}(C_i^2, G_0, G_j, T_{k_{r_j,j}}^i, D_{k_{r_j,j}}^i, 1 \leq j \leq z) \\ & \quad \times \#Df_{Mi, In}^{i,n}(C_i^1, G_0, G_j, A_{k_{r_j,j}}^i, B_{k_{r_j,j}}^i, 1 \leq j \leq z). \end{aligned}$$

Next, use the bijection F to transform $Sim_{deg}^{j_1, j_2}$ into F_2^{d-deg} and do the similar way like the **theorem 7**. Then, the theorem holds. \square

Theorem 9: For any positive integer q, t, z, n , given any z basic function series $\{f(X^{1,q}, Y^{1,q})_{E_{k_{m,j}}, A_{k_{m,j}}^q, B_{k_{m,j}}^q}; 1 \leq m \leq r_j\}_j (X^{1,q}, Y^{1,q} \in F_2^{q \cdot t}), 1 \leq j \leq z$; then we can get $T_{k_{r_j,j}}^i, D_{k_{r_j,j}}^i \in (F_2^t)^{r_j}$ from $A_{k_{m,j}}^q, B_{k_{m,j}}^q$, where $1 \leq j \leq z, 0 \leq i \leq q-1$. And there are two basic function series, of which code are j_1, j_2 respectively, are similar. Then we define the degree of similarity for the two basic function series is deg . Supposed that $G_j \in F_2^{r_j}, 1 \leq j \leq z; \gamma, \lambda, v, w \in F_2^{q \cdot t}, Out, In_1 \in F_2^d, In \in Sim_{deg}^{j_1, j_2}$, where $d = \sum_{i=0}^z kr_i$, $In = (E_{k_{r_1,1}}, \dots, E_{k_{r_z,z}})$. Let $\lambda_q^1 = \lambda, \gamma_q^1 = \gamma, V_q^1 = v, W_q^1 = w$. We have:

$$Cor_{In}^q \left(A_{k_{r_j,j}}^q, \begin{matrix} \gamma_q^1, \lambda_q^1, V_q^1, W_q^1 \\ B_{k_{r_j,j}}^q, G_j \end{matrix}, 1 \leq j \leq z \right) = L \prod_{i=0}^{q-1} Ma^i Q^T$$

where $L = (1, 1, \dots, 1) \in F_2^{d-deg}, Q \in F_2^{d-deg}$, of which the sole nonzero component satisfies $Q[F(In)] = 1$, and $Md^i \in R^{(d-deg) \times (d-deg)}$ satisfying $Ma^i[out, in] = Cor_{F^{-1}(out), F^{-1}(in)}^1 \left(T_{k_{r_j,j}}^i, \begin{matrix} \gamma_i^2, \lambda_i^2, V_i^2, W_i^2 \\ D_{k_{r_j,j}}^i, G_j \end{matrix}, 1 \leq j \leq z \right), 0 \leq i \leq q-1, 0 \leq out, in \leq 2^{d-deg} - 1.$

Proof. According to the **remark 5**, for $1 \leq i \leq q$, we have :

$$Cor_{Out, In}^i \left(A_{k_{r_j,j}}^i, \begin{matrix} \gamma_i^1, \lambda_i^1, V_i^1, W_i^1 \\ B_{k_{r_j,j}}^i, G_j \end{matrix}, 1 \leq j \leq z \right) = 0, \text{ when } in \notin Sim_{deg}^{j_1, j_2} \text{ and } out \in Sim_{deg}^{j_1, j_2}.$$

$$Cor_{Out, In}^i \left(A_{k_{r_j,j}}^i, \begin{matrix} \gamma_i^1, \lambda_i^1, V_i^1, W_i^1 \\ B_{k_{r_j,j}}^i, G_j \end{matrix}, 1 \leq j \leq z \right) = 0, \text{ when } in \in Sim_{deg}^{j_1, j_2} \text{ and } out \notin Sim_{deg}^{j_1, j_2}.$$

Thus, for $1 \leq i \leq q$, when $In \in Sim_{deg}^{j_1, j_2}$,

$$Cor_{In}^i \left(A_{k_{r_j,j}}^i, \begin{matrix} \gamma_i^1, \lambda_i^1, V_i^1, W_i^1 \\ B_{k_{r_j,j}}^i, G_j \end{matrix}, 1 \leq j \leq z \right) = \sum_{Out \in Sim_{deg}^{j_1, j_2}} Cor_{Out, In}^i \left(A_{k_{r_j,j}}^i, \begin{matrix} \gamma_i^1, \lambda_i^1, V_i^1, W_{i+1}^1 \\ B_{k_{r_j,j}}^i, G_j \end{matrix}, 1 \leq j \leq z \right)$$

And for $1 \leq i \leq q-1$, when $In \in Sim_{deg}^{j_1, j_2}$,

$$\begin{aligned} & Cor_{Out, In}^{i+1} \left(\begin{array}{c} \gamma_{i+1}^1, \lambda_{i+1}^1, V_{i+1}^1, W_{i+1}^1, \\ A_{k_{r_j, j}}^{i+1}, B_{k_{r_j, j}}^{i+1}, G_j, 1 \leq j \leq z \end{array} \right) \\ = & \sum_{Mi \in Sim_{deg}^{j_1, j_2}} Cor_{Out, Mi}^1 \left(\begin{array}{c} \gamma_i^2, \lambda_i^2, V_i^2, W_i^2, \\ T_{k_{r_j, j}}^i, D_{k_{r_j, j}}^i, G_j, 1 \leq j \leq z \end{array} \right) \times Cor_{Mi, In}^i \left(\begin{array}{c} \gamma_i^1, \lambda_i^1, V_i^1, W_i^1, \\ A_{k_{r_j, j}}^i, B_{k_{r_j, j}}^i, G_j, 1 \leq j \leq z \end{array} \right) \end{aligned}$$

Next, use the bijection F to transform $Sim_{deg}^{j_1, j_2}$ into F_2^{d-deg} and do the similar way like the **theorem 6**. Then, the theorem holds. \square