Complete Analysis of Implementing Isogeny-based Cryptography using Huff Form of Elliptic Curves

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Abstract. In this paper, we present the complete analysis of Huff curves for implementing isogeny-based cryptography. In this regard, we first investigate the computational cost of the building-blocks when compression functions are used for Huff curves and presented an additional formula on Huff curves for implementing isogeny-based cryptography. From our implementation, the performance of Huff-SIDH and Montgomery-SIDH is almost the same, and Huff-CSIDH is 6.4% faster than Montgomery-CSIDH but 4% slower than Hybrid-CSIDH. The result of our work shows that the Huff curve can be quite practical for implementing isogeny-based cryptography but has some limitations.

Keywords: Isogeny, Post-quantum cryptography, Montgomery curves, square-root Vélu formula, Huff curves, SIDH, CSIDH

1 Introduction

Quantum-resistant cryptosystems based on isogenies were first proposed by Couveignes [13] and later rediscovered by Stolbunov [25], which are currently called as the CRS scheme. However, not only the quantum sub-exponential attack exists for the scheme [8], but the scheme was also inefficient for practical use. After the introduction of the Supersingular Isogeny Diffie-Hellman (SIDH) by De Feo and Jao [18], the isogeny-based cryptography gains back its attention. Due to the non-commutative structure of the endomorphism ring of supersingular curves, SIDH resists the attack proposed in [8]. Additionally, instead of relying on the discrete logarithm problems where the intractability assumption of the problem is broken by Shor's algorithm, the security relies on the problem of finding an isogeny between two given isogenous elliptic curves over a finite field, which is known to have quantum-exponential complexity. The Supersingular Isogeny Key Encapsulation (SIKE), a key encapsulation mechanism based on SIDH, was submitted as one of the candidates to the NIST PQC standardization project [1], and is currently an alternative candidate of Round 3.

Recently, the CRS scheme was revisited by De Feo, Kieffer, and Smith in [14], and independently by Castryck et al. in [7]. The advantage of the CRS scheme is that CCA-secure encryption can be constructed so that a non-interactive key exchange can be obtained. In [14], they modernized the parameter selection in the CRS schemes for better performance and presented an efficient way to compute the CRS group action. The CRS scheme was further optimized by Castryck et al. in [7]. In [7], they proposed CSIDH (Commutative SIDH), which solves the parameter selection problem of the CRS schemes by using supersingular elliptic curves defined over $\mathbb{F}_p$. Currently, the full key exchange of CSIDH at a 128-bit security level requires approximately 80ms, which is slower than SIDH. However, the vital aspect of CSIDH is that a relatively efficient digital signature scheme than SIDH can be constructed based on CSIDH [4]. CSI-FiSh [4] offers a practical digital signature scheme that requires 390ms to sign a message. For isogeny-based cryptography, this was a significant result, which facilitated the construction of various cryptographic primitives through elliptic curve isogenies. To summarize, SIDH and CSIDH have their own advantages, and their common disadvantage is that the performance is slower than any other quantum-resistant algorithm.

The implementation of isogeny-based cryptography involves complicating isogeny operations in addition to the standard elliptic curve arithmetic over a finite field. Regarding the isogeny operations, the degree of an isogeny used in the cryptosystem depends on the prime chosen for the scheme. The SIDH-based algorithms use the prime p of the form $p = \ell_A^{e_A}\ell_B^{e_B}f \pm 1$, where ℓ_A and ℓ_B are coprime to each other. The ℓ_A and ℓ_B corresponds to the degree of isogenies used in the scheme. Since the complexity of computing isogenies increases as the degree increases, isogenies of degree 3- and 4- were mostly considered for implementation. The CSIDH-based algorithms use the prime p of the form $p = 4\ell_1\ell_2\cdots\ell_n - 1$, where ℓ_i are odd-primes. Similarly, as ℓ_i are degrees of isogenies used in the scheme, demands for an efficient odd-degree isogeny formula have increased after the proposal of CSIDH. In [10], Costello and Hisil proposed an efficient way to compute arbitrary odd-degree isogenies on Montgomery curves. Classical ways for computing ℓ-isogeny requires $\tilde{O}(\ell)$ field operations. In [3], Bernstein et al. proposed the square-root Vélu formula which computes the ℓ-isogeny in $\tilde{O}(\sqrt{\ell})$ field operations. This ground-breaking work allows computing higher odd-degree isogenies efficiently, which suits well for implementing CSIDH and B-SIDH [9]. Regarding the elliptic curve arithmetic, it is important to select the form of elliptic curves that can provide efficient curve operations. The majority of isogeny-based cryptography implementations use Montgomery curves as it offers fast isogeny computation and curve arithmetic. The state-of-the-art implementation proposed in [11, 12] is also based on Montgomery curves.

Currently, there is ongoing research on whether other forms of the elliptic curve can yield efficient arithmetic or isogeny computation. The primary candidate is twisted Edwards curves, as it is birationally equivalent to Montgomery curves, and mapping a point on one curve to a point on the other curve is costless when projective coordinates are used. The first use of Edwards curves was by

Meyer et al. in [23], which used twisted Edwards curves for elliptic curve arithmetic and Montgomery curves for isogeny computation [23]. This was further optimized in [20], which used the Edwards curves for isogeny computation and Montgomery curves for the elliptic curve arithmetic. However, as stated in [5] and [20], using only Edwards curves for implementing SIDH-based algorithms is not as efficient as using only Montgomery curves.

The efficiency of using Edwards curves began to stand out when used for implementing CSIDH. Unlike SIDH-based algorithms, CSIDH-based algorithms use higher odd-degree isogenies. Montgomery curves offer efficient isogeny evaluation of arbitrary odd-degree isogenies [10]. However, it is hard to obtain an efficient formula for recovering the coefficient of the image curve on Montgomery curves. On the other hand, Edwards curves can provide an efficient formula for computing the coefficient of the image curve. Therefore, in [22], they implemented CSIDH by using Montgomery curves for isogeny evaluation and twisted Edwards curves for recovering the coefficient of the image curve. In [21], they proposed an optimized odd-degree isogeny formula by using the w-coordinate on Edwards curves. By adapting the formula in [21], Edwards-only CSIDH can be implemented, which is faster than Montgomery-CSIDH [7] or Hybrid-CSIDH [22]. The work of [21] shows that a certain form of an elliptic curve can lead to a better result for certain isogeny-based algorithms. Hence it is important to check the implementation results on various elliptic curves.

This work aims to provide an efficient method to exploit Huff curves in isogeny-based cryptography. Isogenies on Huff curves were first proposed in [24]. However, due to inefficient elliptic curve arithmetic and isogeny formula, it has not been studied until the work of [15], and in [17]. The proposed compression functions in [15,17] for the points on a Huff curve allow Montgomery-like elliptic curve arithmetic formulas, which result in faster isogeny computations. In this paper, based on the formula presented in [15] and [17], we examine the applicability of Huff curves for isogeny-based cryptography. The following list details the main contributions of this work.

- We examine the computational costs of the lower-level functions – differential addition, doubling, and isogeny computation – on Huff curves when a compression method in [15] and [17] are used. These are presented in Section 3. Also, we present 4-isogeny on Huff curves using the compression method proposed in [15], by applying a similar method to derive the 4-isogeny formula presented in [17]. The formulas for 4-isogeny are presented in the Appendix.
- We apply the square-root Vélu formula on Huff curves to compute higher-odd degree isogenies efficiently. To use the square-root Vélu formula proposed in [3], biquadratic polynomials must be redefined to express the relationship between points $P, Q, P-Q$, and $P+Q$ in w-coordinate. We derive biquadratic polynomials for Huff curves and demonstrate that the computational cost for evaluating the main polynomial h_S is the same as Montgomery curves. The definition of h_S and details of the formula is presented in Section 4.
- We also present the formula to recover the coefficient of the Huff curve, for SIDH-based cryptography. As using the compression method in [15] is

faster for recovering the coefficient, we used compression method in [15] to implement SIDH. For CSIDH-based cryptography, one has to examine where supersingular Huff curves exist for a chosen prime. We deduce that for a prime $p \equiv 7 \mod 8$, there exist supersingular Huff curves over $\mathbb{F}_p$, and $\mathbb{F}_p$ has no supersingular Huff curves when $p \equiv 3 \mod 8$.
- We present the implementation result of SIDH and CSIDH using Huff curves. Based on our experiment, for SIDH, the performance of Montgomery-SIDH and Huff-SIDH is almost the same. For CSIDH, Huff-CSIDH is 6.4% faster than Montgomery-CSIDH. When the square-root Vélu formula is used, then Montgomery-CSIDH is 4% faster than Huff-CSIDH. The details of the results are presented in Section 5.

This paper is organized as follows: In Section 2, we introduce two main isogeny-based key exchange algorithms and a form of Huff curves that will be used for the implementation. In Section 3, we demonstrate the computational cost of lower-level functions for implementing isogeny-based cryptography. In Section 4, we introduce the square-root Vélu formula and present the main polynomials for Huff curves to exploit the square-root Vélu formula. In Section 5, we present the implementation result of SIDH and CSIDH on Huff curves. We draw our conclusion in Section 6.

2 Preliminaries

In this section, we provide the necessary background that will be used throughout the paper. First, we introduce two main streams in isogeny-based cryptography – SIDH and CSIDH. Lastly, we describe variants of Huff curves and their arithmetic.

2.1 Isogeny-based cryptosystems

We recall the SIDH and CSIDH key exchange protocol proposed in [18] and [7]. For more information, please refer to [18] and [7] for SIDH and CSIDH, respectively. The notations used in this section will continue to be used throughout the paper.

SIDH protocol Fix two coprime numbers ℓ_A and ℓ_B. Let p be a prime of the form $p = \ell_A^{e_A}\ell_B^{e_B}f \pm 1$ for some integer cofactor f, and e_A and e_B be positive integers such that $\ell_A^{e_A} \approx \ell_B^{e_B}$. Then construct a supersingular elliptic curve E over $\mathbb{F}_{p^2}$ of order $(\ell_A^{e_A}\ell_B^{e_B}f)^2$. We have full ℓ^e-torsion subgroup on E over $\mathbb{F}_{p^2}$ for $\ell \in \{\ell_A, \ell_B\}$ and $e \in \{e_A, e_B\}$. Choose basis $\{P_A, Q_A\}$ and $\{P_B, Q_B\}$ for the $\ell_A^{e_A}$- and $\ell_B^{e_B}$-torsion subgroups, respectively.

Suppose Alice and Bob want to exchange a secret key. Let $\{P_A, Q_A\}$ be the basis for Alice, and $\{P_B, Q_B\}$ be the basis for Bob. For key generation, Alice chooses random elements $m_A, n_A \in \mathbb{Z}/\ell_A^{e_A}\mathbb{Z}$, not both divisible by ℓ_A, and computes the subgroup $\langle R_A \rangle = \langle [m_A]P_A + [n_A]Q_A \rangle$. Then using Vélu's formula, Alice

computes a curve $E_A = E/\langle R_A \rangle$ and an isogeny $\phi_A : E \to E_A$ of degree $\ell_A^{e_A}$, where $ker\phi_A = \langle R_A \rangle$. Alice computes and sends $(E_A, \phi_A(P_B), \phi_A(Q_B))$ to Bob. Bob repeats the same operation as Alice so that Alice receives $(E_B, \phi_B(P_A), \phi_B(Q_A))$.

For the key establishment, Alice computes the subgroup $\langle R'_A \rangle = \langle [m_A]\phi_B(P_A) + [n_A]\phi_B(Q_A) \rangle$. By using Vélu's formula, Alice computes a curve $E_{AB} = E_B/\langle R'_A \rangle$. Bob repeats the same operation as Alice and computes a curve $E_{BA} = E_A/\langle R'_B \rangle$. The shared secret between Alice and Bob is the j-invariant of E_{AB}, *i.e.* $j(E_{AB}) = j(E_{BA})$.

CSIDH protocol CSIDH uses commutative group action on supersingular elliptic curves defined over a finite field $\mathbb{F}_p$. Let $\mathcal{O}$ be an imaginary quadratic order. Let $\mathcal{E}\ell\ell_p(\mathcal{O})$ denote the set of elliptic curves defined over $\mathbb{F}_p$ with the endomorphism ring $\mathcal{O}$. It is well-known that the class group $Cl(\mathcal{O})$ acts freely and transitively on $\mathcal{E}\ell\ell_p(\mathcal{O})$. We call the group action as CM-action and denote the action of an ideal class $[\mathfrak{a}] \in Cl(\mathcal{O})$ on an elliptic curve $E \in \mathcal{E}\ell\ell_p(\mathcal{O})$ by $[\mathfrak{a}]E$.

Let $p = 4\ell_1\ell_2 \cdots \ell_n - 1$ be a prime where $\ell_1, \cdots, \ell_n$ are small distinct odd primes. Let E be a supersingular elliptic curve over $\mathbb{F}_p$ such that $\mathrm{End}_p(E) = \mathbb{Z}[\pi]$, where $\mathrm{End}_p(E)$ is the endomorphism ring of E over $\mathbb{F}_p$. Note that $\mathrm{End}_p(E)$ is a commutative subring of the quaternion order $\mathrm{End}(E)$. Then the trace of Frobenius is zero, hence $E(\mathbb{F}_p) = p + 1$. Since $\pi^2 - 1 = 0 \mod \ell_i$, the ideal $\ell_i\mathcal{O}$ splits as $\ell_i\mathcal{O} = \mathfrak{l}_i\bar{\mathfrak{l}_i}$, where $\mathfrak{l}_i = (\ell_i, \pi - 1)$ and $\bar{\mathfrak{l}_i} = (\ell_i, \pi + 1)$. The group action $[\mathfrak{l}_i]E$ (resp. $[\bar{\mathfrak{l}_i}]E$) is computed via isogeny $\phi_{\mathfrak{l}_i}$ (resp. $\phi_{\bar{\mathfrak{l}_i}}$) over $\mathbb{F}_p$ (resp. $\mathbb{F}_p$) using Vélu's formulas.

Suppose Alice and Bob want to exchange a secret key. Alice chooses a vector $(e_1, \cdots, e_n) \in \mathbb{Z}^n$, where $e_i \in [-m, m]$, for a positive integer m. The vector represents an isogeny associated to the group action by the ideal class $[\mathfrak{a}] = [\mathfrak{l}_1^{e_1} \cdots \mathfrak{l}_n^{e_n}]$, where $\mathfrak{l}_i = (\ell_i, \pi - 1)$. Alice computes the public key $E_A := [\mathfrak{a}]E$ and sends E_A to Bob. Bob repeats the similar operation with his secret ideal $\mathfrak{b}$ and sends the public key $E_B := [\mathfrak{b}]E$ to Alice. Upon receiving Bob's public key, Alice computes $[\mathfrak{a}]E_B$ and Bob computes $[\mathfrak{b}]E_A$. Due to the commutativity, $[\mathfrak{a}]E_B$ and $[\mathfrak{b}]E_A$ are isomorphic to each other so that they can derive a shared secret value from the elliptic curves.

2.2 Huff curves and their arithmetic

Huff curves Huff models for elliptic curves was first introduced by Joye, Tibouchi, and Vergnaud in [19]. They proposed the group law and formula for computing Tate pairings on Huff form of elliptic curves. Let K be a finite field of characteristic not equal to 2. The Huff form of elliptic curve is given by the equation:

$$H_{a,b} = ax(y^2 - 1) = by(x^2 - 1)$$

where $a^2 \neq b^2$ and $a, b \neq 0$. The point $O = (0, 0)$ is the neutral element and $-(x, y) = (-x, -y)$. Also, every Huff curve has three points at infinity, which

are also points of order 2. The curve $H_{a,b}$ can also be simplified as

$$H_c = cx(y^2 - 1) = y(x^2 - 1)$$

where $c = a/b$, $c \neq \pm 1$. The general Huff curves which contains the Huff form of elliptic curves is introduced in [27]. General Huff curves are given by the equation

$$G_{a,b} = x(ay^2 - 1) = y(bx^2 - 1)$$

where $a \neq b$ and $a, b \neq 0$. Similar to the Huff curves, the point $O = (0, 0)$ is the neutral element and $-(x, y) = (-x, -y)$. The j-invariant of the curve $G_{a,b}$ is $j = \frac{2^8(a^2-ab+b^2)^3}{a^2b^2(a-b)^2}$, and the j-invariant of the curve $H_{a,b}$ is $j = \frac{2^8(a^4-a^2b^2+b^4)^3}{a^4b^4(a^2-b^2)^2}$.

Isomorphisms The Huff curve $H_{a,b}$ is isomorphic to a Weierstrass curve of the form

$$W_{A,B} : y^2 = x^3 + Ax^2 + Bx$$

where $A = (a^2 + b^2)$ and $B = a^2b^2$. The Huff curve $H_{a,b}$ is isomorphic to an Edwards curve of the form

$$E : x^2 + y^2 = 1 + \left(\frac{a-b}{a+b}\right)^2 x^2y^2$$

and corresponding Montgomery curve of the form

$$M_D : y^2 = x^3 + Dx^2 + x$$

where $D = (a^2 + b^2)/ab$.

Arithmetic on Huff curves For points $P = (x_p, y_p)$ and $Q = (x_q, y_q)$ on a Huff curve $H_{a,b}$, the addition of two points $P + Q = (x_r, y_r)$ is defined as below, and doubling can be performed with exactly the same formula.

$$x_r = \frac{(x_p + x_q)(1 + y_py_q)}{(1 + x_px_q)(1 - y_py_q)}$$
$$y_r = \frac{(y_p + y_q)(1 + x_px_q)}{(1 - x_px_q)(1 + y_py_q)}$$

The above formula is same for the curve H_c. For a general Huff curve $G_{a,b}$, the unified addition is performed as below:

$$x_r = \frac{(x_p + x_q)(ay_py_q + 1)}{(bx_px_q + 1)(ay_py_q - 1)}$$
$$y_r = \frac{(y_p + y_q)(bx_px_q + 1)}{(bx_px_q - 1)(ay_py_q + 1)}$$

where $P = (x_p, y_p)$ and $Q = (x_q, y_q)$ are points on $G_{a,b}$, and $P + Q = (x_r, y_r)$.

3 w-coordinates on Huff curves

Recently, in [15] and independently in [17], they proposed a compression function on Huff curves, which allows faster elliptic curve arithmetic and isogeny computation. We shall express this compression function as w-coordinate and examine the computational cost of formulas on Huff curves when w-coordinate is used. For the simplicity of the explanation, we shall denote w-function for the compression method proposed in [15], and w_{inv}-function for the compression method proposed in [17].

3.1 Compression function w_{inv} on Huff curves

Huang et al. proposed a compression method for Huff curves and presented an isogeny formula on Huff curves [17]. For the simplicity of the formula, they used the Huff curve of the form H_c.

Let $P = (x, y)$ be a point on a Huff curve H_c. In [17], the defined the compression function w_{inv} as $w(P) = 1/xy$. Then $w(P) = w(-P)$ and $w(O) = \infty$. Using this function, doubling and differential addition formula are defined in [17]. Now, for $P_1, P_2 \in H_c$, let $w_1 = w(P_1)$ and $w_2 = w(P_2)$. Let $w_0 = w(2P_1), w_3 = w(P_1 + P_2)$, and $w_4 = w(P_1 - P_2)$. Then,

$$w_0 = \frac{(w_1^2 - 1)^2}{4w_1(w_1 + c)(w_1 + 1/c)}$$
$$w_3 w_4 = \frac{(w_1 w_2 - 1)^2}{(w_1 - w_2)^2},$$

For the rest of this subsection, we examine the computational cost of the doubling, differential additions, and odd-degree isogeny formula on Huff curves when w_{inv} is used for the compression. We consider WZ-coordinate as projective w-coordinate on Huff curves, where $w = W/Z$. The $\mathbf{M}$ and $\mathbf{S}$ refer to a field multiplication and squaring, respectively.

Doubling Let $P = (x, y)$ be a point on a Huff curve H_c. Let $\hat{c} = \hat{C}/\hat{D}$, where $\hat{c} = \frac{1}{4}\left(c + \frac{1}{c} - 2\right)$. Let $w = 1/xy$, and $w = W/Z$. For $w(P) = (W : Z)$ in projective w-coordinates, the doubling of P gives $w([2]P) = (W' : Z')$, where W' and Z' are defined as:

$$W' = \hat{D}(W - Z)^2(W + Z)^2$$
$$Z' = 4WZ(\hat{D}(W + Z)^2 + \hat{C} \cdot 4WZ)$$

The computational cost is $4\mathbf{M} + 2\mathbf{S}$, given $\hat{C}$ and $\hat{D}$.

Differential addition Let $P_1 = (W_1 : Z_1)$ and $P_2 = (W_2 : Z_2)$ be the points on H_c. Let $w_0 = w(P_1 - P_2)$ and $w_3 = w(P_1 + P_2)$. Let $w_0 = W_0/Z_0$ and $w_3 = W_3/Z_3$. Then,

$$W_3 = Z_0(W_1W_2 - Z_1Z_2)^2,$$
$$Z_3 = W_0(W_1Z_2 - Z_1W_2)^2.$$

The computational cost of differential addition and doubling on Huff curves is $6\mathbf{M}+4\mathbf{S}$ using affine curve coefficients and $8\mathbf{M}+4\mathbf{S}$ using projective coordinates and projective curve coefficients.

Odd-degree isogeny formula Using the compression function w_{inv} for the points on a Huff curve, the odd-degree isogeny formula is presented as the following theorem [17]:

Theorem 1 (Odd-degree isogeny on H_c using w_{inv}-function [17]).

Let P be a point on a Huff curve H_c of odd order $\ell = 2s + 1$. Let $\langle P\rangle = \{(0,0), \pm(\alpha_1, \beta_1), \cdots, \pm(\alpha_s, \beta_s)\}$, where $P = (\alpha_1, \beta_1)$. Let $w_i = 1/\alpha_i\beta_i$ for $1 \leq i \leq s$. and $w = w(Q)$, where $Q = (x, y) \in H_c$. Then for ℓ-isogeny ϕ from H_c to $H_{\hat{c}} = H_c/\langle P\rangle$ the evaluation of w, $w(\phi)$, is given by,

$$w(\phi) = w\prod_{i=1}^{s}\frac{(ww_i - 1)^2}{(w - w_i)^2} \tag{1}$$

where

$$\hat{c} = c\prod_{i=1}^{s}\frac{(1 + cw_i)^2}{(c + w_i)^2}$$

To transform the above formula by using projective coordinates and projective curve coeffcients, for $(\alpha_i, \beta_i) \in H_c$, let $(W_i : Z_i) = (w_i : 1)$ for $i = 1, ..., s$ where $w_i = 1/\alpha_i\beta_i$. For an additional input point $(W : Z)$ on the curve H_c, the output is expressed as $(W' : Z')$ where $(W' : Z') = \phi(W : Z)$. Then, the equation (1) can be rewritten as:

$$W' = W \cdot \prod_{i=1}^{s}(WW_i - ZZ_i)^2,$$
$$Z' = Z \cdot \prod_{i=1}^{s}(WZ_i - ZW_i)^2.$$

Similarly, for $\ell = 2s + 1$-isogeny, evaluation of an isogeny using w_{inv}-function costs $(4s)\mathbf{M}+2\mathbf{S}$. To compute the curve coefficients, let $c = C/D$ and $\hat{c} = \hat{C}/\hat{D}$. Then we have,

$$\hat{C} = C \cdot \prod_{i=1}^{s}(DZ_i + CW_i)^2$$
$$\hat{D} = D \cdot \prod_{i=1}^{s}(CZ_i + DW_i)^2$$

Therefore, recovering the curve coefficient costs $(4s)\mathbf{M}+2\mathbf{S}$.

Coefficient transformation When w_{inv}-function is used for elliptic curve arithmetic on Huff curves, instead of using the projective curve coefficients C and D, we use $(C-D)^2$ and $4CD$ for efficient computation. Hence, after obtaining the coefficient of the image curve $\hat{C}$ and $\hat{D}$, $(\hat{C}-\hat{D})^2$ and $4\hat{C}\hat{D}$ must be computed in order to proceed with elliptic curve arithmetic on the image curve. Intuitively, this requires $2\mathbf{S}$.

3.2 Compression function w on Huff curves

In [15], they proposed a compression method for Huff curves and presented an isogeny formula. As they proposed the formulas for elliptic curve arithmetic and isogenies for Huff curve of the form $H_{a,b}$, we shall present the formulas in this setting. However, for the implementation, we apply the compression function w on H_c, for simplicity. Note that this is equivalent to the case on $H_{a,b}$, with $b=1$, and the resulting formula is almost the reciprocal of the formula on H_c using w_{inv} as a compression method. For detailed formula on H_c using w, please refer to the Appendix.

For a point $P=(x,y)$ on a Huff curve $H_{a,b}$, define the compression function w as $w(P)=xy$. Then $w(P)=w(-P)$ and $w(O)=0$. Using this function, doubling and differential addition can be expressed as follows [15].

For $P_1, P_2 \in H_{a,b}$, let $w_1=w(P_1)$ and $w_2=w(P_2)$. Let $w_0=w(2P_1), w_3=w(P_1+P_2)$, and $w_4=w(P_1-P_2)$. Then,

$$w_0=\frac{4w_1(w_1^2+ew_1+1)}{(w_1^2-1)^2}, w_3w_4=\frac{(w_1-w_2)^2}{(w_1w_2-1)^2},$$

where $e=\frac{b}{a}+\frac{a}{b}$.

For the rest of this subsection, we examine the computational cost of the doubling, differential additions, and odd-degree isogeny formula on Huff curves, in the setting of isogeny-based cryptosystems.

Doubling Let $P=(x,y)$ be a point on a Huff curve $H_{a,b}$. Let $a=A/D$, $b=B/D$, $w=xy$, and $w=W/Z$. For $w(P)=(W:Z)$ in projective w-coordinates, the doubling of P gives $w([2]P)=(W':Z')$, where W' and Z' are defined as:

$$W'=4WZ(4AB(W-Z)^2+(A+B)^2(4WZ))$$
$$Z'=4AB((W-Z)^2(W+Z)^2)$$

The computational cost is $4\mathbf{M}+2\mathbf{S}$, when we assume that $(A+B)^2$ and $4AB$ are precomputed.

Differential addition Let $P_1 = (W_1 : Z_1)$ and $P_2 = (W_2 : Z_2)$ be the points on $H_{a,b}$. Let $w_0 = w(P_1 - P_2)$ and $w_3 = w(P_1 + P_2)$. Let $w_0 = W_0/Z_0$ and $w_3 = W_3/Z_3$. Then,

$$W_3 = Z_0(W_1Z_2 - W_2Z_1)^2,$$
$$Z_3 = W_0(W_1W_2 - Z_1Z_2)^2.$$

The computational cost of differential addition and doubling on Huff curves is $6\mathbf{M}+4\mathbf{S}$ using affine curve coefficients and $8\mathbf{M}+4\mathbf{S}$ using projective coordinates and projective curve coefficients.

Odd-degree isogeny formula In [15], they proposed The odd-degree isogeny formula on Huff curves by composing the isomorphism between Huff and general Huff curves, and odd-degree isogeny on general Huff curves.

Theorem 2 (Odd-degree isogeny on $H_{a,b}$ using w-function [15]).
Let P be a point on a Huff curve $H_{a,b}$ of odd order $\ell = 2s + 1$. Let $\langle P\rangle = \{(0, 0), \pm(\alpha_1, \beta_1), \cdots, \pm(\alpha_s, \beta_s)\}$, where $P = (\alpha_1, \beta_1)$. Let $w_i = \alpha_i\beta_i$ for $1 \le i \le s$. and $w = w(Q)$, where $Q = (x, y) \in H_{a,b}$. Then for ℓ-isogeny ϕ from $H_{a,b}$ to $H_{a',b'} = H_{a,b}/\langle P\rangle$ the evaluation of w, $w(\phi)$, is given by,

$$w(\phi) = w\prod_{i=1}^{s}\frac{(w - w_i)^2}{(ww_i - 1)^2} \tag{2}$$

where

$$a' = \frac{a\prod_{i=1}^{s}(bw_i + a)}{\prod_{i=1}^{s} w_i(aw_i + b)} \quad \text{and} \quad b' = \frac{b\prod_{i=1}^{s}(aw_i + b)}{\prod_{i=1}^{s} w_i(bw_i + a)}$$

Now to projectivize the formula, for $(\alpha_i, \beta_i) \in H_{a,b}$, let $(W_i : Z_i) = (w_i : 1)$ for $i = 1, ..., s$ where $w_i = \alpha_i\beta_i$. For an additional input point $(W : Z)$ on the curve $H_{a,b}$, the output is expressed as $(W' : Z')$ where $(W' : Z') = \phi(W : Z)$. Then, the equation (2) can be rewritten as:

$$W' = W \cdot \prod_{i=1}^{s}(WZ_i - ZW_i)^2,$$
$$Z' = Z \cdot \prod_{i=1}^{s}(WW_i - ZZ_i)^2,$$

where the computing $(WZ_i - ZW_i)$ and $(WW_i - ZZ_i)$ requires $2\mathbf{M}$. Hence for $\ell = 2s + 1$-isogeny, evaluation of an isogeny costs $(4s)\mathbf{M}+2\mathbf{S}$. To compute the curve coefficients, let $a = A/D$ and $b = B/D$. Then we have,

$$A' = A \cdot \prod_{i=1}^{s} Z_i(BW_i + AZ_i)^2,$$
$$B' = B \cdot \prod_{i=1}^{s} Z_i(AW_i + BZ_i)^2,$$
$$D' = D \cdot \prod_{i=1}^{s} W_i(AW_i + BZ_i)(BW_i + AZ_i).$$

Note that only A' and B' are required when implementing isogeny-based cryptography on Huff curves. Moreover, since we use only the ratio of A and B, the term Z_i can be omitted when computing A' and B'. Therefore, recovering the curve coefficient costs $(4s)\mathbf{M}+2\mathbf{S}$.

Coefficient transformation Note that when w-function is used for elliptic curve arithmetic on Huff curves, instead of using the projective curve coefficients A, B, and D, we use $(A+B)^2$ and $4AB$ for efficient computation. Hence, after obtaining the coefficient of the image curve A' and B', $(A'+B')^2$ and $4A'B'$ must be computed in order to proceed with the elliptic curve arithmetic on the image curve. Intuitively, this requires $2\mathbf{S}$.

Remark 1. It is obvious that when w-function is used for H_c, the formula for doubling, differential addition, and odd-degree isogenies is almost the reciprocal of the case when w_{inv}-function is used for H_c. Let $P_1 = (x_1, y_1)$ and $P_2 = (x_2, y_2)$ be the points on H_c, and $w(P) = xy$ for $P = (x, y) \in H_c$. Let $w_1 = w(P_1)$, $w_2 = w(P_2)$, $w_0 = w(2P_1), w_3 = w(P_1 + P_2)$, and $w_4 = w(P_1 - P_2)$. Then,

$$w_0 = \frac{4w_1(w_1 + c)(w_1 + 1/c)}{(w_1^2 - 1)^2}$$
$$w_3 w_4 = \frac{(w_1 - w_2)^2}{(w_1 w_2 - 1)^2},$$

which is almost the reciprocal of the case when w_{inv}-function is used. Hence every computational cost for elliptic curve arithmetic and isogeny is the same when H_c with w is used and when H_c with w_{inv} is used.

4 Square-root Vélu formula for Huff curves

Recently, Bernstein et al. proposed an efficient algorithm that computes ℓ-isogeny in $\tilde{O}(\sqrt{\ell})$ field operations [3]. The conventional Vélu formula computes ℓ-isogeny in $\tilde{O}(\ell)$ field operations. The high-level view of the Vélu formula can be considered as evaluation of polynomials over K whose roots are values of a function

from a cyclic group to K. Let G be a cyclic group with generator P. Then for a finite subset S of $\mathbb{Z}$, define a polynomial

$$h_S(X) = \prod_{s \in S} (X - f([s]P)) \tag{3}$$

where $[s]P$ denotes the sum of s copies of P. In isogeny-based setting, let $E(K)$ be an elliptic curve, $P \in E(K)$. Then $G = \langle P \rangle$ is a kernel of an ℓ-isogeny $\phi : E \to E'$, and $f([s]P)$ can be considered as the x-coordinate of P.

Let M_a be a Montgomery curve, $P \in M_a$ be a point of prime order $\ell \neq 2$. The isogeny $\phi : M_a \to M_{a'}$ with kernel $\langle P \rangle$ is given by the equation below, expressed in terms of equation (3):

$$\phi(X) = \frac{X^\ell \cdot h_S(1/X)^2}{h_S(X)^2}$$

where $a' = 2(1+d)/(1-d)$ for $d = ((a-2)/(a+2))^\ell \cdot (h_S(1)/h_S(-1))^8$, and $S = \{1, 3, \ldots, \ell - 2\}$. Now, $\phi(X)$ can be evaluated in $\tilde{O}(\sqrt{\ell})$ field operations if h_S is evaluated in $\tilde{O}(\sqrt{\ell})$ field operations.

The key for evaluating h_S in $\tilde{O}(\sqrt{\ell})$ field operations is to decompose the set S into smaller set I and J, having size similar to $\sqrt{S}$, satisfying certain conditions. In [3], I and J are chosen so that most of the elements in S is represented as elements of $(I + J) \cup (I - J)$. For details on the conditions of the set and algorithms, please refer to [3]. Hence, the problem of evaluating a polynomial whose roots are $[s]P$ for $s \in P$ is transformed to the problem of evaluating a polynomial, whose roots are $[i]P$ and $[j]P$ for $i \in I$ and $j \in J$, respectively. Then, by computing the resultant of polynomials relating to the set I and J, we can obtain the evaluation of h_S. To do this, we need to find the relations between the x-coordinate of $[i]P, [j]P, [i+j]P$, and $[i-j]P$ for $i \in I$ and $j \in J$. Below, Lemma 1 states the existence of biquadratic polynomials of an elliptic curve E that shows the relationship between points $P, Q, P+Q$, and $P-Q$ for $P, Q \in E$.

Lemma 1 (Biquadratic relations on x-coordinates [3]). *Let q be a prime power. Let $E(\mathbb{F}_q)$ be an elliptic curve. There exist biquadratic polynomials F_0, F_1, and F_2 in $\mathbb{F}_q[X_1, X_2]$ such that*

$$(X - x(P+Q))(X - x(P-Q)) = X^2 + \frac{F_1(x(P), x(Q))}{F_0(x(P), x(Q))} X + \frac{F_2(x(P), x(Q))}{F_0(x(P), x(Q))}$$

for all $P, Q \in E$ such that $O \notin \{P, Q, P+Q, P-Q\}$. The $x(P)$ denotes the x-coordinate of a point P.

If E is defined by affine Montgomery equation $By^2 = x^3 + Ax^2 + x$, then the polynomials F_0, F_1, and F_2 are defined as follows [3].

$$\begin{aligned}
F_0(X_1, X_2) &= (X_1 - X_2)^2 \\
F_1(X_1, X_2) &= -2((X_1 X_2 + 1)(X_1 + X_2) + 2A X_1 X_2) \\
F_2(X_1, X_2) &= (X_1 X_2 - 1)^2
\end{aligned}$$

To use the square-root formula, we define the following biquadratic polynomials specifically for Huff curves of the form H_c. Similarly, the relationship between the w-coordinates of points $P, Q, P+Q$, and $P-Q$ on Huff curves can be written as follows:

$$(W - w(P+Q))(W - w(P-Q)) = W^2 + \frac{G_1(w(P), w(Q))}{G_0(w(P), w(Q))}W + \frac{G_2(w(P), w(Q))}{G_0(w(P), w(Q))}$$

For the curve H_c using the w-function, then the polynomials G_0, G_1, and G_2 are defined as follows:

$$\begin{aligned} G_0(W_1, W_2) &= (W_1W_2 - 1)^2 \\ G_1(W_1, W_2) &= -2((W_1W_2 + 1)(W_1 + W_2) + 2\hat{C}W_1W_2 + 4W_1W_2) \\ G_2(W_1, W_2) &= (W_1 - W_2)^2 \end{aligned}$$

where $\hat{C} = c + \frac{1}{c} - 2$. When w_{inv}-function is used for compression, then the polynomials G_0, G_1, and G_2 are defined as follows:

$$\begin{aligned} G_0(W_1, W_2) &= (W_1 - W_2)^2 \\ G_1(W_1, W_2) &= -2((W_1W_2 + 1)(W_1 + W_2) + 2\hat{C}W_1W_2 + 4W_1W_2) \\ G_2(W_1, W_2) &= (W_1W_2 - 1)^2 \end{aligned}$$

where $\hat{C} = c + \frac{1}{c} - 2$. Using this biquadratic polynomials, the square-root formula for Huff curves directly follows [3], and the computational cost for evaluating biquadratic polynomials for Huff curve is the same as Montgomery curves. The following proposition states the isogeny formula on Huff curves H_c using w- and w_{inv}-function, expressed in terms of equation (3).

Proposition 1 (Square-root Vélu formula on Huff curves). *Let H_c be an elliptic curve over $\mathbb{F}_q$ in Huff form, and let P be a point of prime order $\ell \neq 2$ in H_c. Let w be a compression function for point on H_c. For $Q \in H_c$ let $w(Q) = W$. Then the evaluation of $w(\phi(Q))$ where $\phi : H_c \to H_{c'}$, a quotient isogeny with kernel $\langle P \rangle$, is given as:*

$$w(\phi(W)) = \frac{W^\ell h_S(W)^2}{h_S(1/W)^2}$$

where $S = \{1, 3, \ldots, \ell - 2\}$ and $c' = c^\ell \cdot h_S(-c)^2 / h_S(-1/c)^2$.

Now, let w_{inv} be a compression function for point on H_c. For $Q \in H_c$ let $w_{inv}(Q) = W$. Then the evaluation of $w_{inv}(\phi(Q))$ where $\phi : H_c \to H_{c'}$, a quotient isogeny with kernel $\langle P \rangle$, is given as:

$$w_{inv}(\phi(W)) = \frac{W^\ell h_S(1/W)^2}{h_S(W)^2}$$

where $S = \{1, 3, \ldots, \ell - 2\}$ and $c' = c^\ell \cdot h_S(-1/c)^2 / h_S(-c)^2$.

Remark 2. While preparing this paper, we notice the recent work by Wroński, introducing the application of the square-root Vélu formula on Huff curves [26]. In [26], a new compression functions on Huff are introduce to be suitable for application.

Recovering the curve coefficient From the biquadratic polynomials, we were able to derive the formula for recovering the coefficient of the Huff curve from the w-coordinate of the points P, Q and $P - Q$ on a Huff curve.

When implementing SIDH-based cryptography, $P_A - Q_A$ and $P_B - Q_B$ are also considered as a public key for faster kernel computation using the Montgomery ladder. Hence $\phi_A(P_B - Q_B)$ and $\phi_B(P_A - Q_A)$ are also computed and exchanged to compute the shared secret key efficiently. This can be thought of as an increase in the public key size. But using the fact that the coefficient a of the Montgomery curve M_a relates to the x-coordinates of P, Q, and $P - Q$ for $P, Q \in M_a$, sending the coefficient of the image curve is omitted [12]. Therefore, $(\phi_A(P_B), \phi_A(Q_B), \phi_A(P_B - Q_B))$ and $(\phi_B(P_A), \phi_B(Q_A), \phi_B(P_A - Q_B))$ are exchanged during the protocol, and upon the receipt of the public key, the coefficient is recovered using the relationship, which costs $4\mathbf{M}+1\mathbf{S}+1\mathbf{I}$. The $\mathbf{I}$ denotes the field inversion.

For Huff curves, similar relationship can be obtained. Let H_c be a Huff curve using w as a compression function. For P, Q, and $P - Q$ in H_c, let $w(P) = w_p, w(Q) = w_q$, and $w(P - Q) = w_{pq}$. Then the following holds:

$$w(P+Q) + w(P-Q) = \frac{2((w_p w_q + 1)(w_p + w_q) + 2\hat{c} w_p w_q + 4 w_p w_q)}{(w_p w_q - 1)^2}$$
$$\frac{(w_p - w_q)^2}{w_{pq}(w_p w_q - 1)^2} + w_{pq} = \frac{2((w_p w_q + 1)(w_p + w_q) + 2\hat{c} w_p w_q + 4 w_p w_q)}{(w_p w_q - 1)^2}$$

so that

$$\begin{aligned}\hat{c} &= \frac{(w_p - w_q)^2 + w_{pq}^2(w_p w_q - 1)^2 - 2w_{pq}((w_p w_q + 1)(w_p + w_q) + 4 w_p w_q)}{4 w_{pq} w_p w_q} \\ &= \frac{((w_p - w_q) - (w_{pq}(w_p w_q - 1)))^2 - 4 w_{pq}(w_p + w_p w_q^2 + 2 w_p w_q)}{4 w_{pq} w_p w_q} \end{aligned} \tag{4}$$

where $\hat{c} = c + \frac{1}{c} - 2$. The computational cost is $3\mathbf{M}+1\mathbf{S}+1\mathbf{I}$. Similar relationship can be obtain for w_{inv}-function, which is as follows.

$$w_{inv}(P+Q) + w_{inv}(P-Q) = \frac{2((w_p w_q + 1)(w_p + w_q) + 2\hat{c} w_p w_q + 4 w_p w_q)}{(w_p - w_q)^2}$$
$$\frac{(w_p w_q - 1)^2}{w_{pq}(w_p - w_q)^2} + w_{pq} = \frac{2((w_p w_q + 1)(w_p + w_q) + 2\hat{c} w_p w_q + 4 w_p w_q)}{(w_p - w_q)^2}$$

so that

$$\begin{aligned}\hat{c} &= \frac{(w_p w_q - 1)^2 + w_{pq}^2(w_p - w_q)^2 - (2w_{pq}((w_p w_q + 1)(w_p + w_q) + 4w_p w_q)}{4w_{pq} w_p w_q} \\ &= \frac{((w_p w_q - 1) - (w_{pq}(w_p - w_q)))^2 - 4w_{pq}(w_p + w_p w_q^2 + 2w_p w_q)}{4w_{pq} w_p w_q} \end{aligned} \tag{5}$$

where $\hat{c} = c + \frac{1}{c} - 2$. The computational cost is $5\mathbf{M}+1\mathbf{S}+1\mathbf{I}$.

Summarizing the section, Table 1 denotes the computational cost of the building blocks of isogeny-based cryptography on Montgomery curves and on Huff curves. The middle rule in Table 1 divides the functions into two groups – the upper half is the functions that are commonly used in SIDH and CSIDH-based cryptography, and the lower half is the functions that are explicitly used in SIDH-based cryptography.

In Table 1, `DBLADD` refers to the differential addition and doubling in projective coordinates, and `DBL` refers to the doubling. ℓ-`isog eval` refers to the evaluation of an ℓ-isogeny and ℓ-`isog coeff` refers to the computation of the coefficient of the ℓ-isogenous image curve, where $\ell = 2s + 1$. `CoeffTrans` refers to the cost of transforming the coefficient for efficient elliptic curve arithmetic, which only occurs on Huff curves. The `TPL` refers to tripling of a point, and `3-isogeny` (resp. `4-isogeny`) is the combined computational cost of isogeny evaluation and coefficient computation. Lastly, `get_coeff` refers to recovering of the curve coefficient using points P, Q and $P - Q$ on an elliptic curve.

Also, `Mont` refers to Montgomery curve, and `Hybrid` refers to the hybrid method proposed by [22], where Montgomery curve is used for elliptic curve arithmetic and isogeny evaluation, and Edwards curves are used for computing the coefficient of the image curve. As the hybrid method is only used for a comparison between curves in CSIDH-based cryptography, the computational cost of the lower half of the table is omitted. The function $w(\ell)$ refers to $w(\ell) = (h - 1)\mathbf{M} + (t - 1)\mathbf{S}$. In $w(\ell)$, h denotes the hamming weight of ℓ and t is the bit length of ℓ.

As shown in Table 1, except for the ℓ-`isogeny coeff`, the computational cost of the lower-level functions is the same for Montgomery curves and Huff curves. Also, as the compression function w and w_{inv} are reciprocals of each other, the formula of the lower-level functions are almost reciprocals of each other so that w-function and w_{inv}-function induce the same computational cost. Hence, when implementing CSIDH-based cryptography on Huff curves, the compression function is free of one's choice. On the other hand, for SIDH-based cryptography, w-function is preferred as `get_coeff` is slightly efficient than w_{inv}-function.

Remark 3. On Huff curves, the `CoeffTrans` can be omitted when division polynomial is used to represent the curve coefficient in terms of the kernel points. This can be easily done for 3- and 4- isogenies. For general higher degree isogenies, as representing the coefficient of the curve using kernel point is difficult, extra `CoeffTrans` operation is required to proceed with the elliptic curve arithmetic further.

Table 1: Computational cost of building-blocks of isogeny-based cryptography on Huff curves and Montgomery curves

	`Mont` [7, 10]	`Hybrid` [22]	w	w_{inv}
`DBLADD`	$6\mathbf{M} + 4\mathbf{S}$	$6\mathbf{M} + 4\mathbf{S}$	$6\mathbf{M} + 4\mathbf{S}$	$6\mathbf{M} + 4\mathbf{S}$
`DBL`	$4\mathbf{M} + 2\mathbf{S}$	$4\mathbf{M} + 2\mathbf{S}$	$4\mathbf{M} + 2\mathbf{S}$	$4\mathbf{M} + 2\mathbf{S}$
ℓ-`isog eval`	$4s\mathbf{M} + 2\mathbf{S}$	$4s\mathbf{M} + 2\mathbf{S}$	$4s\mathbf{M} + 2\mathbf{S}$	$4s\mathbf{M} + 2\mathbf{S}$
ℓ-`isog coeff`	$(6s-2)\mathbf{M} + 3\mathbf{S}$	$(2s)\mathbf{M} + 6\mathbf{S} + 2w(\ell)$	$4s\mathbf{M} + 2\mathbf{S}$	$4s\mathbf{M} + 2\mathbf{S}$
`CoeffTrans`	-	-	$2\mathbf{S}$	$2\mathbf{S}$
`TPL`	$7\mathbf{M} + 5\mathbf{S}$	-	$7\mathbf{M} + 5\mathbf{S}$	$7\mathbf{M} + 5\mathbf{S}$
`3-isogeny`	$6\mathbf{M} + 5\mathbf{S}$	-	$6\mathbf{M} + 5\mathbf{S}$	$6\mathbf{M} + 5\mathbf{S}$
`4-isogeny`	$6\mathbf{M} + 6\mathbf{S}$	-	$6\mathbf{M} + 6\mathbf{S}$	$6\mathbf{M} + 6\mathbf{S}$
j-invariant	$3\mathbf{M} + 4\mathbf{S} + 1\mathbf{I}$	-	$3\mathbf{M} + 4\mathbf{S} + 1\mathbf{I}$	$3\mathbf{M} + 4\mathbf{S} + 1\mathbf{I}$
`get_coeff`	$4\mathbf{M} + 1\mathbf{S} + 1\mathbf{I}$	-	$5\mathbf{M} + 1\mathbf{S} + 1\mathbf{I}$	$3\mathbf{M} + 1\mathbf{S} + 1\mathbf{I}$

5 Implementation

In this section, we provide the performance result of isogeny-based cryptography. First, we present the implementation result of SIDH entirely on Huff curves. Then, we present the implementation result of CSIDH.

To evaluate the performance, the algorithms are implemented in the C language. For SIDH, we modified the field arithmetic implemented in SIDH library version 3.1 [11]. For CSIDH, we use the field arithmetic implemented in [7]. All the cycle counts were obtained on one core of an Intel Core i7-6700 at 3.40 GHz, running Ubuntu 16.04 LTS. For the compilation, we used clang version 6.0.0 with an optimization level -O3.

5.1 Implementation of SIDH

We first present the parameter settings for SIDH implementation. Then we present the implementation result with analysis. For implementing SIDH, we used the Huff curve of the form H_c with w as compression function, as w-function is more efficient than w_{inv} for recovering the coefficient of the curve after the first round of the protocol.

Parameter Settings The prime used in SIDH-based cryptography is of the form $p = \ell_A^{e_A}\ell_B^{e_B}f \pm 1$. In this section, we present two implementations on Huff curves when $\{\ell_A, \ell_B\} = \{2, 3\}$ and $\{\ell_A, \ell_B\} = \{3, 5\}$. The former is the general choice of ℓ_A and ℓ_B for implementing SIDH-based cryptography. As an extra coefficient transformation is required for Huff curves for higher degree isogenies, the latter is to examine the performance change caused by this.

For $\{\ell_A, \ell_B\} = \{2, 3\}$, we used the 751-bit prime proposed in [2], which is as follows:

$$p_{751} = 2^{372} \cdot 3^{239} - 1$$

For $\{\ell_A, \ell_B\} = \{3, 5\}$, we used the 621-bit prime of the form:

$$p_{621} = 2^{67} \cdot 3^{175} \cdot 5^{119} - 1,$$

and $3^{175} \approx 2^{277.368}$ and $5^{119} \approx 2^{276.309}$.

Over finite field $\mathbb{F}_{p_a^2} = \mathbb{F}_{p_a}(i)$ for $i^2 = -1$ and $a \in \{751, 621\}$, we used the supersingular Montgomery curve of the form as the base curve:

$$M : y^2 = x^3 + 6x^2 + x,$$

which is isomorphic to a Huff curve of the form:

$$H_{c_a} : c_a x(y^2 - 1) = y(x^2 - 1).$$

For $a = 751$, then $c_{751} = 3 + \sqrt{8} \in \mathbb{F}_{p_{751}^2}$ and for $a = 621$, then $c_{621} = 3 + \sqrt{8} \in \mathbb{F}_{p_{621}^2}$.

For p_{621}, the generator points for the Huff curve are the points P_A, Q_A and P_B, Q_B such that $P_A, Q_A \in E[3^{175}]$ and both points have exact order 3^{175}, $P_B, Q_B \in E[5^{119}]$ and both points have an exact order 5^{119}. To select such a point, we first search for the points on the following Weierstrass curve:

$$W : y^2 = x^3 + (c_{621}^2 + 1)x^2 + c_{621}^2 x,$$

which is isomorphic to the Huff curve $H_{c_{621}}$. When (P_A, Q_A) and (P_B, Q_B) are found, we compute the Weil paring $e(P_A, Q_A) \in E[3^{175}]$ and $e(P_B, Q_B) \in E[5^{119}]$ to check that the result has order 3^{175} and 5^{119}, respectively. When the points are found, we transform the points on W to points on $H_{c_{621}}$, and express in w-coordinate. The generator points on Montgomery curves are found in a similar manner.

Also, when implementing 5-isogeny, we used the formula from [10] for isogeny evaluation. For recovering the coefficient of the image curve, we used the 2-torsion method described in [10]. The reason is that using the 2-torsion method, the cost for recovering the coefficient of the image curve is $8\mathbf{M}+4\mathbf{S}$, while using the projectivized formula of [10] presented in [7], the cost is $10\mathbf{M}+3\mathbf{S}$.

Implementation Results Table 2 presents the implementation results of SIDH on Montgomery curves and Huff curves. Using the prime p_{621} and p_{751}, the performance of SIDH is compared between Montgomery curves and Huff curves. The implementation using p_{751} uses 3- and 5-isogeny formula, which is presented in Section 3.2. The prime p_{621} uses 3- and 4-isogeny, and the corresponding formula on H_c using w-function is in the Appendix.

As denoted in Table 2, for p_{751}, the performance of the Montgomery-SIDH and Huff-SIDH are almost the same. This is obvious as the computational cost for

Table 2: Performance results of SIDH implementation. The results were rounded to the nearest 10^3 clock cycles.

	Montgomery Curve		Huff Curve	
	p_{621}	p_{751}	p_{621}	p_{751}
Isogeny degree used	3,5	2,3	3,5	2,3
Alice's Keygen	134,497	221,498	134,529	221,449
Bob's Keygen	149,234	248,712	148,495	249,294
Alice's Shared Key	114,053	182,015	114,109	182,089
Bob's Shared Key	133,132	211,899	132,761	211,952
Total	530,916	864,124	529,894	864,784
Security (Classical)	138	186	138	186

the formulas for implementing isogeny-based cryptography is almost the same. For p_{621}, although the Huff curve requires to transform the curve coefficient on Bob's side, the performance of the Montgomery-SIDH and Huff-SIDH is almost the same. We shall analyze the results in detail by dividing them into key generation and shared key computation phases.

Public key generation During this phase, it is natural that there is no difference when comparing the computational cost of the two curves for Alice's side. For Bob's side on Huff curves, after calculating the coefficient of the image curve, extra coefficient transformation is required for efficient quintupling. Hence, computing the coefficient on Huff curves costs $8\mathbf{M}+4\mathbf{S}$ for a total. For Montgomery curves, 2-torsion is used for recovering the coefficient of the image curve. Hence, after evaluating isogeny at 2-torsion point for a Montgomery curve, recovering the curve coefficient is required, and total also costs $8\mathbf{M}+4\mathbf{S}$. Therefore, the computation of 5-isogeny on both curves is almost the same.

Computing the shared key The difference in the computation between two curves occurs when calculating the curve coefficient upon the receipt of $(\phi_i(P_j), \phi_i(Q_j), \phi_i(P_j - Q_j))$ for $(i, j) \in \{(A, B), (B, A)\}$. Now, note that upon the receipt of $(\phi_i(P_j), \phi_i(Q_j), \phi_i(P_j - Q_j))$ for $(i, j) \in \{(A, B), (B, A)\}$, the Huff coefficient $\hat{c} = c + 1/c + 2$ of H_c is recovered using equation (4), not c itself. For Alice, as $\hat{c}$ is directly used for tripling and isogeny computation, the performance on Huff curves and Montgomery curves is almost the same. However, on Bob's side in Huff curves, recovering $\hat{c}$ is not enough – $\hat{c}$ is used for quintupling, but we need the actual c to compute the coefficient of the isogenous curve. On the other hand, for Montgomery curves, Bob uses the extra 2-torsion point on the base curve to compute the image curve's coefficient. To reduce the key size, when computing the shared key on Bob's side, we compute the 2-torsion, given the coefficient of the Montgomery curve. Hence both curves require to solve quadratic equation over $\mathbb{F}_{p^2}$, which requires 1 field multiplication and 1 square-root computation for both curves. Hence, the performance of Montgomery-SIDH and Huff-SIDH is almost the same.

5.2 Implementation of CSIDH

For CSIDH-based cryptography, the compression function to use is free of one's choice as the computational cost of the building blocks for CSIDH is the same for w and w_{inv}. In this paper, we used w_{inv} for implementing CSIDH. To implement CSIDH-based cryptography, we first need to check whether a supersingular curve exists over a given prime field. In this section, we examine the existence of a supersingular Huff curve over $\mathbb{F}_p$ for a prime p and present the base curve H_c for the implementation. Then we present the implementation result of CSIDH using Huff curves.

Prime field and base curve In order to implement the CSIDH-based cryptosystem, we need to search for a supersingular Huff curve $H_{a,b}$ over a prime field p of the form $p = f \cdot \prod \ell_i - 1$, where ℓ_is are small distinct primes. Below is the theorem proving that a supersingular Huff curve over $\mathbb{F}_p$ exists when $p \equiv 7 \mod 8$. If $p \equiv 3 \mod 8$, there is no supersingular Huff curve over $\mathbb{F}_p$.

Theorem 3. *There exists a supersingular Huff curve of the form $H_{a,b}$ over $\mathbb{F}_p$ when $p \equiv 7 \mod 8$.*

Proof. In the CSIDH setting, for every supersingular elliptic curve over $\mathbb{F}_p$, there exist a corresponding supersingular Montgomery curve over $\mathbb{F}_p$. Hence it suffices to show that for a given supersingular Montgomery curve, there exists an isomorphic Huff curve over $\mathbb{F}_p$. Now, Huff curve $H_{a,b}$ is isomorphic to a Montgomery curve of the form:

$$M : y^2 = x^3 + \frac{a^2 + b^2}{ab}x^2 + x \tag{6}$$

Then, M is supersingular if and only if $H_{a,b}$ is supersingular. Let $(a^2 + b^2)/ab = A$. If we find a supersingular Montgomery curve $y^2 = x^3 + Ax^2 + x$ over $\mathbb{F}_p$, then by using the equation (6), we can find the corresponding supersingular Huff curve over $\mathbb{F}_p$. Solving the equation we have,

$$a = \frac{Ab \pm \sqrt{(Ab)^2 - 4b^2}}{2} \tag{7}$$

From the above equation, $H_{a,b}$ is defined over $\mathbb{F}_p$, if and only if $Ab^2 - 4b^2$ is a square in $\mathbb{F}_p$, i.e. $A^2 - 4$ is a square in $\mathbb{F}_p$.

Now, suppose $p \equiv 7 \mod 8$ and let M be a supersingluar curve having a 2-torsion point on $\mathbb{F}_p$ except for $(0, 0)$. Then the 2-torsion subgroup of M satisfy $| M[2] | = 4$. In this case, the supersingular curve M lies on the surface so that $\mathrm{End}_{\mathbb{F}_p}(M) \cong \mathbb{Z}[(1 + \sqrt{-p})/2]$ [6]. Then $A^2 - 4$ is a square in $\mathbb{F}_p$, so that the corresponding $H_{a,b}$ exists over $\mathbb{F}_p$. On the other hand, if $p \equiv 3 \mod 8$, then M lies on the floor so that $A^2 - 4$ is not a square in $\mathbb{F}_p$, so that there is no supersingular Huff curve $H_{a,b}$ over $\mathbb{F}_p$.

The original implementation of CSIDH uses the prime of the form $p \equiv 3 \mod 8$. However, from Theorem 3. we use the 511-bit prime presented in [16], which works over $\mathbb{F}_p$ where

$$p = 2^4 \cdot 3^3 \cdot 5 \cdot 7 \cdot 11^2 \cdot 13 \cdots 373 - 1.$$

In this field, we choose a supersingular Huff curve of the form as the base curve:

$$H_c : cx(y^2 - 1) = y(x^2 - 1)$$

where $c = 3 - \sqrt{8} \in \mathbb{F}_p$.

Remark 4. For a prime p such that $p \equiv 3 \mod 8$, there exist a supersingular general Huff curve over $\mathbb{F}_p$. However, as the computational cost of elliptic curve arithmetic and isogeny evaluations is slower than the Huff curve, we omit this case.

Selecting a random point over $\mathbb{F}_p$ When implementing CSIDH, one has to select a random point on a curve over $\mathbb{F}_p$ of a certain order to compute an isogeny using Velu's formula. For a Montgomery curve, first, a random element in $\mathbb{F}_p$ is selected, and we consider as an x-coordinate of a given Montgomery curve. Then, by using the curve equation $y^2 = x^3 + Ax^2 + x$, $r = x^3 + Ax^2 + x$ is computed and checked whether r is a square or nonsquare in $\mathbb{F}_p$. The computational cost for checking whether a random point on a Montgomery curve is in $\mathbb{F}_p$ or $\mathbb{F}_{p^2} \backslash \mathbb{F}_p$ costs $2\mathbf{M} + 1\mathbf{S}$ (we omit the computational costs for computing the Legendre symbol).

The following method checks whether the point $(x, y) \in H_c$ is on $\mathbb{F}_p$ or $\mathbb{F}_{p^2} \backslash \mathbb{F}_p$. Since $w = 1/xy$ for a point $(x, y) \in H_c$, $y = 1/wx$. Now, from the curve equation, the following holds:

$$\begin{aligned} cx(y^2 - 1) &= y(x^2 - 1) \\ \frac{c}{w}y - cx &= \frac{c}{w}x - y \\ \left(\frac{c}{w} + 1\right) y &= \left(\frac{1}{w} + c\right) x \\ x^2 &= (w + c)/w(1 + cw) \end{aligned}$$

Thus $x \in \mathbb{F}_p$ if $(cw + 1)(cw + w^2)$ is square in $\mathbb{F}_p$. The computational cost for checking whether a random point is in $\mathbb{F}_p$ or $\mathbb{F}_{p^2} \backslash \mathbb{F}_p$ costs $2\mathbf{M} + 1\mathbf{S}$.

Implementation results For the implementation, we used the prime field $\mathbb{F}_p$, as presented in Section 6. In order to compare the performance with Montgomery curves, we use the following supersingular curve over $\mathbb{F}_p$ as a base curve.

$$M : y^2 = x^3 + x$$

The original implementation of Montgomery-CSIDH in [7] does not use the optimization method when evaluating isogenies. Hence, we modified the implementation for a fair comparison with the Huff-CSIDH. The difference in the performance between the algorithms lies purely in the computation of the coefficient of the image curve and coefficient transformation for Huff curves. Table 3 presents the performance of the group action on Montgomery curves and Huff curves. In Table 3, Hybrid-CSIDH is a method proposed in [22], which implements CSIDH using Montgomery curves but uses Edwards curves for evaluating the coefficient of the image curve.

Table 3: Performance results of group action in using traditional Vélu formula.

	Montgomery-CSIDH	Hybrid-CSIDH	Huff-CSIDH
Group action	110,558,288	99,264,050	103,469,978

As shown in Table 3, Huff-CSIDH is 6.4% faster than Montgomery-CSIDH. This is because although an extra coefficient transformation is required when Huff curves are used for the implementation, recovering the Montgomery curve's coefficient is costly than on a Huff curve for odd-degree isogenies. Also, Hybrid-CSIDH is 10.2% faster than Montgomery-CSIDH and 4% faster than Huff-CSIDH.

Additionally, for CSIDH, it is important to optimize the odd-degree isogeny formula as isogeny computation contributes to the overall CSIDH performance. Hence we present the CSIDH implementation using the square-root Vélu formula in [3].

Table 4: Performance results of group action in CSIDH using the square-root Vélu formula.

	`sqrt-Mont`	`sqrt-Huff`
Group action	99,658,531	103,396,849

Table 4 presents the performance comparison of Montgomery-CSIDH and Huff-CSIDH, when the square-root Vélu formula is used for computing odd-degree isogenies. In Table 4, `sqrt-Mont` and `sqrt-Huff` refers to CSIDH implementation when the square-root Vélu formula is used for Montgomery curves and Huff curves, respectively. As the effect of square-root Vélu formula becomes more conspicuous when isogeny of degree larger than 113 is used, the effect is not immediate for the current parameter setting. However, as the square-root Vélu formula exploits Edwards curves for recovering the coefficient of Montgomery curve, we can see that the performance of `sqrt-Mont` is similar to Hybrid-

CSIDH in Table 3. The `sqrt-Huff` is 3.6% slower than `sqrt-Mont`. This is because when computing the coefficient of image curve, `sqrt-Huff` computes $h_S(-c)$ and $h_S(-1/c)$, while `sqrt-Mont` computes $h_S(1)$ and $h_S(-1)$, which is more efficient.

6 Conclusion

In this paper, we present the complete analysis of Huff curves' usage for implementing isogeny-based cryptography. First, we analyzed the computational cost of the lower-level functions when the compression method is used for Huff curves. Then, we proposed additional functions on Huff curves to implement isogeny-based cryptography. We presented implementation results of the two main isogeny-based algorithms – SIDH and CSIDH – on Huff curves.

For SIDH, as the computational cost for the lower-level function is the same on Montgomery curves and Huff curves, the performance of Montgomery-SIDH and Huff-SIDH is almost the same. For CSIDH, we present the birational polynomials on Huff curves in order to exploit the square-root Vélu formula. We implemented CSIDH using the classical Vélu formula and the square-root Vélu formula and compare it with Montgomery curves. The `sqrt-Mont` is 4% faster than `sqrt-Huff`.

Base on our analysis, the Huff curve can be quite practical for implementing isogeny-based cryptography but has some limitations. The first is that as the points of order 2 are all at infinity on Huff curves, it is hard to construct a 2-isogeny formula using w-coordinate so that only e_A with an even number can be used to implement SIDH with Huff curves. The second is that a supersingular Huff curve exists on $\mathbb{F}_p$, where $p \equiv 7 \mod 8$. This result is contrary to the case where supersingular Montgomery curve exists on $\mathbb{F}_p$ for both $p \equiv 3 \mod 8$ and $p \equiv 7 \mod 8$.

Acknowledgement

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Appendix A 4-isogenies on Huff Curves

Although [15] omits the 4-isogeny formula on Huff curves using w-function, by adapting the idea from [17], we additionally present the 4-isogeny formula on Huff curves using w-function. In this section, we shall briefly state the formula for implementing SIDH using 4-isogeny.

Theorem 4 (2-isogenies for $H_{a,b}$ using w-function). *Let $\phi : H_{a,b} \to H_{a',b'}$ be a 2-isogeny with kernel $\{(0,0),(a:b:0)\}$. Let $w = xy$ for $(x,y) \in H_{a,b}$. Then the evaluation of w under ϕ is given by*

$$\phi(w) = \frac{(a^2 - b^2)w}{(bw + a)(aw + b)} \tag{8}$$

where

$$a' = \sqrt{-\left(\sqrt{\frac{1}{b^2}} + \sqrt{\frac{1}{a^2}}\right)^2}, \quad b' = \sqrt{-\left(\sqrt{\frac{1}{b^2}} - \sqrt{\frac{1}{a^2}}\right)^2}$$

To derive equation (8), we adapt the method used in [17]. That is, ϕ is first derived from the below composition:

$$H_{a,b} \xrightarrow{\iota} G_{a,b} \xrightarrow{\psi} G_{\hat{a},\hat{b}} \xrightarrow{\iota^{-1}} H_{a',b'}$$

where ι denotes the transformation from a Huff curve to a general Huff curve, ψ is a 2-isogeny on general Huff curve from [24]. Then $\phi = \iota^{-1} \circ \psi \circ \iota$. Similarly, we can derive Huff 2-isogenies for the curve of the form H_c.

Theorem 5 (2-isogenies for H_c using w-function). *Let $\phi : H_c \to H_{c'}$ be a 2-isogeny with kernel $\{(0,0),(c:1:0)\}$. Let $w = xy$ for $(x,y) \in H_c$. Then the evaluation of w under ϕ is given by*

$$\phi(w) = \frac{w(c^2 - 1)}{(w + c)(cw + 1)} \tag{9}$$

where $c' = \mid (c+1) \mid / \mid (1-c) \mid$.

As shown in equation (9), 2-isogeny on H_c using w-function is identical to setting $b = 1$ in equation (8). Also, equation (9) is just reciprocal of the 2-isogeny on H_c using w_{inv} function defined in [17]. Lastly, we state the 4-isogeny on Huff curve using w-function, directly derived from the idea presented in [17].

Theorem 6 (4-isogenies for H_c using w-function). *Let $\phi : H_c \to H_{c'}$ be a 4-isogeny with kernel P such that $w(P) = w_4$ and P has order 4 in H_c. Let $w(Q) = w = xy$ for a point $Q = (x, y) \in H_c$. Then the evaluation of w under ϕ is given by*

$$\phi(w) = \frac{w(w - w_4)^2(ww_4^2 + w - 2w_4)}{(2ww_4 - w_4^2 - 1)(ww_4 - 1)^2}$$

where $c' = (1 + \sqrt{1 - w_4^4})/(1 - \sqrt{1 - w_4^4})$.

Appendix B Formulas for implementing SIDH-based cryptography

In this section, we present the doubling, tripling, 3-isogeny, and 4-isogeny formula on a Huff curve of the form H_c, using w as a compression function. For corresponding formulas on H_c using w function, please refer to [17].

Doubling Let $P = (x, y)$ be a point on a Huff curve H_c. Let $c = C/D$ and $\hat{c} = \hat{C}/\hat{D}$, where $\hat{c} = \frac{1}{4}\left(c + \frac{1}{c} - 2\right)$. For $w(P) = (W : Z)$ in projective w-coordinates, the doubling of P gives $w([2]P) = (W' : Z')$, where W' and Z' are defined as:

$$W' = 4WZ(\hat{D}(W + Z)^2 + \hat{C} \cdot 4WZ)$$
$$Z' = \hat{D}(W - Z)^2(W + Z)^2$$

The computational cost is $4\mathbf{M} + 2\mathbf{S}$, given $\hat{C}$ and $\hat{D}$. Therefore, instead of using the projective curve coefficient $(C : D)$, it is efficient to use $(\hat{C} : \hat{D}) = ((C-D)^2 : 4CD)$ for implementation.

Tripling Let $P = (x, y)$ be a point on a Huff curve H_c. Let $c = C/D$ and $\hat{c} = \hat{C}/\hat{D}$, where $\hat{c} = \frac{1}{4}\left(c + \frac{1}{c} - 2\right)$. For $w(P) = (W : Z)$ in projective w-coordinates, the tripling of P gives $w([3]P) = (W' : Z')$, where W' and Z' are defined as:

$$W' = W(\hat{D}W^4 - 6\hat{D}W^2Z^2 - 16\hat{C}WZ^3 - 8\hat{D}WZ^3 - 3\hat{D}Z^4)^2$$
$$Z' = Z(3\hat{D}W^4 + 16\hat{C}W^3Z + 8\hat{D}W^3Z + 6\hat{D}W^2Z^2 - \hat{D}Z^4)^2$$

The tripling formula is the same to the case when w_{inv} is used. The computational cost is $7\mathbf{M} + 5\mathbf{S}$, given $\hat{C}$ and $\hat{D}$. For tripling it efficient to keep the projective curve coefficient as $((C - D)^2 : (C + D)^2)$.

3-isogeny Let $P = (x_3, y_3)$ be a 3-torsion point on a Huff curve H_c, $w(P) = (W_3 : Z_3)$. Let $\phi : H_c \to H_{c'}$ be a 3-isogeny generated by a kernel $\langle P \rangle$, such that $H_{c'} = H_c/\langle P \rangle$. Let $Q = (W : Z)$ be another point on H_c. Then the image $w(\phi(Q)) = (W' : Z')$ is computed as:

$$W' = W(WZ_3 - ZW_3)^2$$
$$Z' = Z(WW_3 - ZZ_3)^2$$

and

$$\hat{C} = (W_3 - Z_3)(W_3 + 3Z_3)^3$$
$$\hat{D} = (W_3 + Z_3)(W_3 - 3Z_3)^3$$

where $c' = C'/D'$ and $\hat{C} = (C' + D')^2$ and $\hat{D} = (C' - D')^2$, to continue with the tripling efficiently. The computational cost for 3-isogeny evaluation is $4\mathbf{M} + 2\mathbf{S}$ and the computational cost for computing the coefficient of the image curve is $2\mathbf{M} + 3\mathbf{S}$.

4-isogeny Let $P = (x_4, y_4)$ be a 4-torsion point on a Huff curve H_c, $w(P) = (W_4 : Z_4)$. Let $\phi : H_c \to H_{c'}$ be a 4-isogeny generated by a kernel $\langle P \rangle$, such that $H_{c'} = H_c/\langle P \rangle$. Let $Q = (W : Z)$ be another point on H_c. Then the image $w(\phi(Q)) = (W' : Z')$ is computed as:

$$W' = W(2W_4Z_4Z - W(W_4^2 + Z_4^2))(Z_4W - W_4Z)^2$$
$$Z' = Z(2W_4Z_4W - Z(W_4^2 + Z_4^2))(WW_4 - ZZ_4)^2$$

and

$$\hat{C} = 4Z_4^4 - 4W_4^4$$
$$\hat{D} = 4W_4^4$$

where $\hat{c} = \hat{C}/\hat{D} = \frac{1}{4}\left(c' + \frac{1}{c'} - 2\right)$, to continue with the doubling efficiently. The computational cost for 4-isogeny evaluation is $6\mathbf{M} + 2\mathbf{S}$ and the computational cost for computing the coefficient of the image curve is $4\mathbf{S}$.