

# A note on a Claim of Eldar & Hallgren: LLL already solves it

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**Abstract.** In a recent talk of Hallgren on a joint work with Eldar (Sept 21, 2021, Simons Institute), a polynomial-time quantum algorithm for solving BDD in a certain class of lattice was claimed. We show here that known classical (and even, deterministic) polynomial-time algorithms already achieve this result.

## 1 Context and Claims

The problem considered by Eldar and Hallgren [Hal21] can be read as a worst-case version of the LWE problem, with a secret dimension  $k = 1$ ,  $n$  samples, modulus  $q = c^n$  for some  $c > 1$ , and a subexponential approximation factor  $\alpha = 1/2^{\Theta(\sqrt{n})}$ .

More formally, let us start by defining the Bounded Distance Decoding.

**Definition 1.1 (Bounded Distance decoding (BDD)).** *The BDD problem in a lattice  $L \subset \mathbb{R}^n$  and radius with  $r > 0$  is the problem of, given  $\mathbf{t} = \mathbf{v} + \mathbf{e}$  for a lattice vector  $\mathbf{v} \in L$  and an error  $\mathbf{e} \in \mathbb{R}^n$  with  $\|\mathbf{e}\| < r$ , finding  $\mathbf{v}$ .*

For the solution to be unique, one requires  $r/\lambda_1(L) \leq 1/2$ . More generally, this ratio is referred to as the BDD approximation factor.

The family of lattices considered in [Hal21] are the  $q$ -ary lattices spanned by a single vector  $\mathbf{a} \in \mathbb{Z}^n$

$$L_{\mathbf{a}} = q\mathbb{Z}^n + \mathbf{a}\mathbb{Z}.$$

**Theorem 1.2 (Eldar & Hallgren [Hal21]).** *There exists a quantum polynomial-time algorithm that solves BDD in  $L_{\mathbf{a}}$  for any  $\mathbf{a} \in \mathbb{Z}^n$  and for any error up to radius  $\lambda_1(L_{\mathbf{a}}) \cdot 2^{-\Theta(\sqrt{n})}$ .*

In the average-case, this problem with these parameters is already known to be easy to solve, simply by ignoring all but  $O(\sqrt{n})$  many samples (geometrically, a projection onto certain canonical axes), applying the LLL reduction algorithm to the basis, and finally decoding with Babai nearest plane algorithm. During the panel discussion following the presentation, various expert discussed the plausibility of a provable classical algorithm achieving the same result via known randomization techniques.

While we share their optimism regarding the plausibility of such a classical rerandomisation, we will show that such randomization is not even needed !

Namely we will prove that the LLL [LLL82] and Babai [Bab86] algorithms already solves the problem in the full dimensional lattice, in the worst-case, and *deterministically*.

Proving so requires considering the  $q$ -ary structure of the lattice, and other guarantees of LLL than its approximation factor. Such reasoning are not new, and already played a role in lattice cryptanalysis [CL15, KF17]. More specifically, a key remark in our case is to note that the "perp lattice" (the dual lattice scaled up by  $q$ ) is an integer lattice with small determinant; the situation appears as the dual of [CL15].

We also provide constant in the exponent for more refined comparison. To this end, let us introduce  $\delta = \sqrt{4/3} + \epsilon$  for some arbitrary small  $\epsilon > 0$  as the constant appearing in László condition in LLL [LLL82]. The constant  $c > 1$  below is the constant such that  $q = c^n$ .

**Theorem 1.3 (This note).** *There exists a deterministic polynomial-time algorithm that solves BDD in  $L_{\mathbf{a}}$  for any  $\mathbf{a} \in \mathbb{Z}$  for any error up to radius  $\lambda_1(L_{\mathbf{a}}) \cdot (c\sqrt{\delta})^{-\sqrt{n}-O(1)}$ .*

## 2 Proof

The volume of this lattice is an integer comprised between  $q^{n-1}$  and  $q^n$ . A key remark to show that LLL already solves the problem is to also state the profile of a systematic basis, namely, a basis starting with  $n - 1$  distinct  $q$ -vectors  $(0, \dots, 0, q, 0, \dots, 0)$ . Denoting  $\mathbf{b}_i$  those basis vectors for  $i \in \{1, \dots, n\}$ ,  $\mathbf{b}_i^*$  the associated Gram-Schmidt vectors, and  $\ell'_i = \ln \|\mathbf{b}_i^*\|$ , we have:

$$\ell'_i = \ln q \text{ for } i < n, \quad \ell'_n \in [0, \ln q].$$

Let us now denote  $(\mathbf{c}_i)_i$  the basis obtained by applying LLL to *this*<sup>1</sup> basis  $(\mathbf{b}_i)_i$ , and denote  $(\mathbf{c}_i^*)_i$  the associated Gram-Schmidt vectors, and finally  $\ell_i = \|\mathbf{c}_i^*\|$ .

The first constraint on the profile is what is typically used to control LLL-reduced bases, together with the invariant  $\sum \ell_i = \sum \ell'_i$ .

**Fact 2.1 (László Condition)** *For all  $i \in \{1, \dots, n - 1\}$ ,*

$$\ell_{i+1} \geq \ell_i - \ln \delta. \tag{1}$$

But there is much more that can be said about the LLL algorithm. In particular, one can also show that partial volumes do not increase during the algorithm (this is even a key fact to prove termination of LLL).

**Fact 2.2 (Partial Volumes Decreases)** *For all  $i \in \{1, \dots, n\}$ :*

$$\sum_{j=1}^i \ell_j \leq \sum_{j=1}^i \ell'_j. \tag{2}$$

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<sup>1</sup> as opposed to taking an arbitrary LLL-reduced basis

At last, one can invoke duality to upper-bound the aggregated Gram-Schmidt length of the last vectors. Indeed, the dual lattice is contained in  $\frac{1}{q}\mathbb{Z}^n$ , so any partial basis of the dual of dimension  $d$  has volume at least  $1/q^d$ . By duality (see [Mic17, Lecture 3 (Duality), Section 3] or [DD18, Lecture 5, Section 3]), this implies the following.

**Fact 2.3 (Left-over volume for  $q$ -ary lattices)** For all  $i \in \{1, \dots, n\}$ :

$$\sum_{j=i+1}^n \ell_j \leq (n-i) \ln q. \quad (3)$$

The rest of our proof is a game of inequalities towards the following.

**Proposition 2.4.** With the notations above, and  $\lambda_1$  the minimal distance of our lattice, it holds that

$$\min_i \ell_i \geq \ln \lambda_1 - \sqrt{n} \cdot \ln(c\sqrt{\delta}) + O(1).$$

The final claim that Babai will properly solve BDD in the worst-case directly follows.

*Proof.* For  $i \leq \sqrt{n}$ , the bound  $\ell_i \geq \ln \lambda_1 - \sqrt{n} \ln \delta$  directly follows from Lovasz condition, noting that  $\ell_1 \geq \ln \lambda_1$  because  $\mathbf{c}_0^* = \mathbf{c}_0$  is a non-zero lattice vector.

It remains to prove the statement for  $i \geq \lceil \sqrt{n} \rceil$ . First, note that the volume invariant gives:

$$(n-1) \ln q \leq \sum_{j=1}^n \ell_j \leq n \ln q \quad (4)$$

Because of the  $q$ -ary structure of the input basis, this in fact generalizes to any partial volumes:

$$(i-1) \ln q \leq \sum_{j=1}^i \ell_j \leq i \ln q, \quad \text{for all } i \in \{1, \dots, n\} \quad (5)$$

The upper bounds follows from Fact 2.2 and the upper bound  $\ell'_i \leq \ln q$  on the initial basis. The lower bound follows from subtracting the inequality of Fact 2.3 to the volume invariant.

Let  $d = \lceil \sqrt{n} \rceil$ . Taking the difference of (5) at  $i \geq d$  and  $i-d+1$  we get:

$$\sum_{j=i-d+1}^i \ell_j \geq (d-1) \ln q. \quad (6)$$

By Lovasz condition (2.1) we further have that  $\ell_j \leq \ell_i + (i-j) \ln \delta$ , which implies

$$d\ell_i + \sum_{j=0}^{d-1} j \ln \delta \geq (d-1) \ln q. \quad (7)$$

which rewrites as:

$$d\ell_i + d(d-1)\frac{\ln \delta}{2} \geq (d-1) \ln q. \quad (8)$$

Using that  $q = c^n$ ,  $\lambda_1 \leq q$ , and  $d = \lceil \sqrt{n} \rceil$  we get

$$\min_{d \leq i \leq n} \ell_i \geq \frac{d-1}{d} \ln q - \frac{1}{2}(d-1) \ln \delta \quad (9)$$

$$\geq \ln(q) - \frac{n}{d} \ln c - \frac{1}{2}d \ln \delta \quad (10)$$

$$\geq \ln \lambda_1 - \sqrt{n} \ln c - \frac{1}{2} \lceil \sqrt{n} \rceil \ln \delta \quad (11)$$

$$\geq \ln \lambda_1 - \sqrt{n} \cdot \ln(c\sqrt{\delta}) + O(1). \quad (12)$$

□

To conclude the proof of our main theorem, it remains to invoke the correctness condition for solving BDD with Babai's algorithm.

**Fact 2.5 (Correctness of Babai's Algorithm [Bab86])** *Given a basis  $(\mathbf{c}_i)_i$  of a lattice  $L$ , with associated Gram-Schmidt basis  $(\mathbf{c}_i^*)_i$ , Babai's Nearest Plane algorithms solves BDD up to a radius  $r = \min \|\mathbf{c}_i^*\|/2$ .*

**Acknowledgments.** The authors are grateful to Zvika Brakerski, Nico Dötting, Chris Peikert, Noah Stephens-Davidowitz, Damien Stehlé and Vinod Vaikuntanathan for helpful discussions.

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