# XORBoost: Tree Boosting in the Multiparty Computation Setting 

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#### Abstract

We present a novel protocol XORBoost for both training gradient boosted tree models and for using these models for inference in the multiparty computation (MPC) setting. Similarly to [2], our protocol is the first one supporting training for generically split datasets (vertical and horizontal splitting, or combination of those) while keeping all the information about the features and thresholds associated with the nodes private, thus, having only the depths and the number of the binary trees as public parameters of the model. By using optimization techniques reducing the number of oblivious permutation evaluations as well as the quicksort and real number arithmetic algorithms from the recent Manticore MPC framework [5], we obtain a scalable implementation operating under information-theoretic security model in the honest-butcurious setting with a trusted dealer. On a training dataset of 25,000 samples and 300 features in the 2-player setting, we are able to train 10 regression trees of depth 4 in less than 5 minutes per tree (using histograms of 128 bins).


## 1 Introduction

Gradient boosting is a machine learning technique for regression and classification problems that yields a prediction model in the form of an ensemble of weak prediction models, typically decision trees [14]. XGBoost [1], [6] is currently one of the most popular open-source libraries supporting gradient boosting for various programming environments and architectures.

Multiparty computation (MPC) is a method for cryptographic computing allowing several parties holding private data to evaluate a public function on their aggregate data while revealing only the output of the function and nothing else. Recent advances in the area make these protocols practical and suitable for real-world applications among which machine and statistical learning [3], [4], [5], [10], [12], [15], [16], [17], [21], [22], [24].

### 1.1 Prior work

Specific attempts have been made to adapt classical boosting methods to privacypreserving settings, both for training and inference [7], [13], [18], [19], [20].

In [7] and [13], frameworks based on federated learning and homomorphic encryption are proposed that allow for training a boosting model on vertically split datasets, that is, datasets where for each feature, all the data belongs to a single party ${ }^{1}$. While this is suitable for applications where external private features can be added to enhance the model performance, it does not cover horizontally split data, that is, cases where private datasets with the same features can be concatenated to build a larger dataset (a classical federated learning/edge computing scenario). This limitation is not present in our work, where any data partition type is supported.

A different method for training that does not require vertical splitting is proposed in [19]. The protocol based on federated learning, secret-sharing and homomorphic encryption reveals some information about the structure of every tree: for every internal node, the feature and threshold leading to the maximum reduction of the loss function are revealed. This potentially leaks sensitive information about the original data and that can be problematic for some use cases. Even more information about the underlying data and the trained model is revealed in [7].

Recent work on XGBoost inference based on fully-homomorphic encryption ( FHE ) is studied in [20].

A solution based on secure enclaves is proposed in [18]. As such, the security model is significantly different from the one considered here. Our approach is based on information theoretic security as opposed to hardware security. Even if measures are taken to obfuscate memory access and thus limit side channel attacks in [18], the approach remains vulnerable to attacks targeting the secure enclave.

Finally, in [2], a protocol is presented for training and evaluating classification trees on data split horizontally or vertically (or any mixture thereof). The authors implemented their protocol on the MP-SPDZ framework [15]. The authors of [2] observe that the output of their chosen tree learning algorithm (C4.5 trees) only depends on the relative order of the input, which allows to map the input data to the integer domain in an arbitrary, but order-preserving manner and to skillfully avoid costly privacy-preserving operations on fixed- or floating-point numbers.

The gradient boosted tree model outlined in here is based on XGBoost [6] and is inherently different from single classification and regression trees or random forests. Whereas random forest trains many trees independently on subsampled datasets, XGBoost constitutes an ensemble of learners by, at each step, adding to the ensemble the tree with the greatest loss reduction. Furthermore the protocol outlined in this paper leverages fixed-point arithmetic, which allows to compute prediction weights accurately to train regression trees instead of being limited to classification trees with categorical response variables.

[^0]
### 1.2 Our contributions

We present a generic algorithm to train and evaluate a gradient boosted tree model based on [6] in the MPC setting. As previously mentioned, our method supports any combination of horizontal and vertical splittings of the dataset.

In contrast to [19] and [7], the computing parties in our protocol learn only the shape of the model, i.e., the tree depth and the number of trees.

The feature indices and the threshold values associated to the non-leaf nodes as well as the weights (or prediction values) associated to the leaf nodes are all secret-shared. We also ensure that during training, no information about the first and second order statistics is revealed. Furthermore, during training and prediction, it is impossible to deduce the path taken by any sample in a decision tree. Similarly to [2], our protocol is built on generic primitives and can thus be implemented on any MPC framework supporting these primitives.

We used the Manticore MPC framework which provides access to Boolean arithmetic as well as arithmetic with real numbers represented using modular integers [5] or the prior floating-point numbers framework [4]. Manticore is operating in semi-honest security model with an offline trusted dealer, with full-threshold security across an arbitrary number of players. Furthermore all communication between the trusted dealer and the players, as well as all communication between the players during the online phase, is end-to-end encrypted. This makes it secure against malicious external adversaries.

By building on top of Manticore's representation of real numbers via modular integers, our protocol is secure in the information-theoretic setting.

Similarly to [2], Manticore allows us to privately compute the permutation sorting a given feature column. Applying these permutations is an expensive operation and we introduce in section 3.2 a novel algorithm that allows to reduce the number of permutation calls with respect to the naive algorithm.

The improvements presented in this paper allow to train a gradient boosted tree model of moderate depth within a reasonable time frame (see Section 8). These include the use of a very efficient sorting mechanism [5], the precomputation of generator vectors to apply inverse permutations (Section 3.2), the use of compressed instance vectors (Section 4.2) and storing permuted instance vectors to reduce the number of times the permutation function is called. Taken together, the number of permutation calls after an initial preprocessing phase is $2+D$ per tree, with $D$ the tree depth.

## 2 Background and Preliminaries

For a detailed review of XGBoost, we refer the reader to [6]. Consider a dataset $X$ of size $N \times k$ and a response variable $y$ (a vector of size $N$ ). We use $X^{(j)}$ to refer to column $j$ of $X$ and $X_{i}$ to refer to the $i$ th row of $X$. Thus, $X_{i}^{(j)}$ denotes the $i$ th element of the $j$ th column.

### 2.1 Binary decision trees

A binary decision tree of depth $D$ on the feature space $\mathbb{R}^{k}$ consists of $2^{D}-1$ inner nodes (also called non-leaf nodes or split nodes) and $2^{D}$ outer nodes (also called leaf nodes). Associated to each inner node is a feature index (in $\{1, \ldots, k\}$ ) and a threshold. Associated to each leaf node is a weight value. We thus represent a tree as

```
Tree = (TreeStructure, TreeWeights),
```

where

$$
\text { TreeStructure }=\left(\left(j_{1}, t_{1}\right), \ldots,\left(j_{2^{D}-1}, t_{2^{D}-1}\right)\right)
$$

and

$$
\text { TreeWeights }=\left(w_{1}, \ldots, w_{2^{D}}\right)
$$

We further refine the TreeStructure to consist of $D$ layers at depth $0,1, \ldots, D-$ 1 and of size $2^{0}, 2^{1}, \ldots, 2^{D-1}$ respectively. The evaluation of a tree on a sample $\mathbf{x}=\left(x_{1}, \ldots, x_{k}\right)$ is as the following recursive procedure:

```
eval(x, node) =
    if node is a leaf (w): return w
    if node is a inner ((j,t), nleft , nright ):
            if }\mp@subsup{x}{j}{}<t\mathrm{ return eval( }\mathbf{x},\mp@subsup{n}{\mathrm{ left }}{})\mathrm{ else eval( }\mathbf{x},\mp@subsup{n}{\mathrm{ right }}{})
```

And by extension eval( $\mathbf{x}$, Tree) means to call eval on $\mathbf{x}$ and the root node of Tree: In a prediction scenario when the Tree is fixed and the sample varies, we often abbreviate the notation as $\operatorname{Tree}(\mathbf{x})$ seen as a $\mathbb{R}^{k} \rightarrow \mathbb{R}$ piecewise constant function, and on training scenario where $\mathbf{x}$ is fixed and the Tree varies, we use eval $\mathbf{x}_{\mathbf{x}}$ (Tree), which continuously differentiable over the TreeWeights for each TreeStructure.

If there are many samples, we write $\operatorname{Tree}(X) \in \mathbb{R}^{N}$, to mean Tree evaluated at each row of $X$. Given a tree ensemble $\left\{\operatorname{Tree}{ }^{(1)}, \ldots, \operatorname{Tree}^{(T)}\right\}$, one obtains predictions on $X$ as

$$
\begin{equation*}
\widehat{y}^{(T)}=\sum_{t=1}^{T} \operatorname{Tree}^{(t)}(X) \in \mathbb{R}^{N} \tag{1}
\end{equation*}
$$

In order to reduce overfitting during the boosting procedure, it is standard to use a weighting coefficient $\eta$ called the learning rate when updating $\widehat{y}$ with the prediction from a new tree in 1.

### 2.2 Objective function

Gradient tree boosting is an iterative process using a current prediction $\widehat{y}^{(T)}$ on T trees to greedily (see [14] for a definition) grow a new $\mathrm{T}+1$-th tree that most reduces a certain objective function. For a given function loss: $\mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$, such
as e.g. mean-squared error or logistic loss and fixed training set $X, y$, consider the function

$$
\begin{equation*}
\mathcal{L}\left(\operatorname{Tree}{ }^{(T+1)}\right)=\sum_{i=1}^{N} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}+\operatorname{eval}_{X_{i}}\left(\operatorname{Tre} \mathrm{e}^{(T+1)}\right)\right)+\operatorname{Reg}\left(\operatorname{Tree}{ }^{(T+1)}\right) \tag{2}
\end{equation*}
$$

The regularization function Reg is used to reduce overfitting by penalizing large parameter values (similarly to Ridge and Lasso regression models). We use L2regularization on the leaf weights

$$
\operatorname{Reg}\left(\operatorname{Tree}^{(t)}\right)=\gamma|L|+\frac{\lambda}{2} \sum_{\ell \in L} w_{\ell}^{2}
$$

where $\lambda$ and $\gamma$ are fixed hyperparameters, $L$ is the set of leaves of $\operatorname{Tree}{ }^{(t)}$ and $w_{\ell}$ is the leaf weight for leaf $\ell \in L$.

The overall goal is to find a tree that minimizes the objective function $\mathcal{L}$. Note that for a fixed TreeStructure the restriction of the objective function $\mathcal{L}$ on the TreeWeights space is differentiable, convex in the case of logistic loss, and even quadratic in the case of mean-square loss. The basic idea of the greedy XGBoost algorithm is to recursively take a Tree (initially a single leaf), replace one of its leaves by a fixed number of splits, and for each of these potential tree structures, retain the one that shows the maximal gain on $\mathcal{L}$, and repeat the process until the tree is a full tree of depth $D$. Under the assumption that $\mathcal{L}$ is equal or well-approximated by its second-order expansion at zero the gain on $\mathcal{L}$ is estimated as

$$
\text { score }(\operatorname{TreeStructure})=\frac{1}{2} \operatorname{grad}(\mathcal{L})^{t} \cdot \operatorname{Hessian}^{-1}(\mathcal{L}) \cdot \operatorname{grad}(\mathcal{L})
$$

with gradient and hessian over the TreeWeights space, all evaluated in zero. In the present case, the score formula simplifies into the following lemma, which corresponds to equation 6 of [6], and a proof is also given in Appendix A:

Lemma 1. Let $g_{i}=\partial_{b} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}\right)$ and $h_{i}=\partial_{b}^{2} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}\right)$ for $1 \leq i \leq$ $N$, we easily verify (see [6]) that:

$$
\begin{equation*}
\operatorname{score}(\text { TreeStructure })=\sum_{\text {leaves } L} \frac{G^{2}}{H+\lambda}, \tag{3}
\end{equation*}
$$

where $G=\sum_{i \in L} g_{i}$ and $H=\sum_{i \in L} h_{i}$.
Between each step, we update the tree structure by picking the split (feature and threshold value) maximizing the score: since we split only one node at each time, we only need to account for the contribution of the new left and right leaves in the above score, the rest of the leaves remain constant.

Consider now splitting a node $n$, i.e. attaching two children nodes $n_{\text {left }}$ and $n_{\text {right }}$ to $n$. The gain associated to this split is the difference in loss resulting from attaching these two children nodes:

$$
\begin{align*}
\operatorname{gain} & =\operatorname{score}\left(G_{n_{\text {left }}}, H_{n_{\text {left }}}, \lambda\right)+\operatorname{score}\left(G_{n_{\text {right }}}, H_{n_{\text {right }}}, \lambda\right)  \tag{4}\\
& -\operatorname{score}\left(G_{n}, H_{n}, \lambda\right)-\gamma,
\end{align*}
$$

(where $\operatorname{score}\left(G_{n}, H_{n}, \lambda\right)=\operatorname{score}\left(G_{n_{\text {left }}}+G_{n_{\text {right }}}, H_{n_{\text {left }}}+H_{n_{\text {right }}}, \lambda\right)$ ). Note that the gain needs to take into account $-\gamma$, since splitting a node results in an increment of the total number of nodes.

### 2.3 MPC represenation of a tree

The TreeStructure is represented as follows: for every internal node $n$, the feature index $j_{n}$ is secret and is encoded as an additively secret-shared elementary vector $\mathrm{e}_{n}$ of size $k$. As way of example, if we need to access the $j_{n}$ th column, we can multiply the original dataset $X$ with $\mathrm{e}_{n}$. This technique is used throughout to keep the tree structure private. The threshold $\mathbf{t}_{n}$ is secret and is additively secret-shared.
The TreeWeights are secret and are additively secret-shared.

## 3 Data Preprocessing Phase

Note that one of the challenges for splitting a node into a left child and a right child in the overall xgboost training procedure is the need to maximize the gain. A priori, this is done over one discrete variable (the feature) and one continuous variable (the threshold corresponding to that feature). To discretize the search for the threshold, we use histograms for the feature values. Computing these histograms in privacy-preserving manner is challenging - it requires to obliviously sort the feature vectors and extract the sorting permutations, as explained in Section 3.1.

The sorting permutations are required later on in the training procedure (see Algorithm 6). In Section 3.2 we introduce a novel algorithm, that allows to reduce the number of times these permutations need to be applied, leading to a significant gain in efficiency $\left(\mathcal{O}\left(\log _{2} B\right)\right.$ instead of $\mathcal{O}(B)$ per feature, where $B$ is the number of bins in the histogram).

### 3.1 Sorting feature vectors and building histograms

The Manticore framework enables for efficient sorting of numerical vectors. It does so by applying a dealer-generated, secret-shared uniformly random permutation to the target vector, on which a variant of quicksort is applied (c.f. Section 5.2 of [5]). Throughout the computation, the input vector remains secret-shared and nothing about the underlying data is revealed. As an output we receive the sorted vector and / or the secret-shared sorting permutation and their inverses.

While the Manticore algorithm supports oblivious sorting of vectors with repeated values, this method is not suited for categorical feature vectors, for which refer the reader to section 6 .

We now introduce some notation used throughout this paper.

- Given a permutation $\sigma$ on $N$ elements and a feature vector $X^{(j)}$ of size $N$, $\sigma$ acts on $X^{(j)}$ by permuting the indices:

$$
\sigma\left(X^{(j)}\right):=\left(X_{\sigma(i)}^{(j)}\right)_{i=1}^{N}
$$

We define the function apply_permutation, that takes as input a secretshared vector $v$ of size $N$, together with a list of $r$ permutations $\sigma_{1}, \ldots, \sigma_{r}$ and returns a $N \times r$ matrix, whose $j$-th column corresponds to $\sigma_{j}(v)$.

- We use $\pi_{j}$ as a label for the sorting-permutation of feature-column $X^{(j)}$, and we use $\pi_{j}^{-1}$ to denote the inverse permutation.
3.1.1 Bucket vectors As mentioned in the introduction of this section, we restrict the possible choices of the thresholds by relying on histograms. More specifically, given the number of bins $B$ in the histogram and a feature vector $X^{(j)}$, we only consider the $B-1$ values

$$
\boldsymbol{t}_{b}^{(j)}:=\pi_{j}\left(X^{(j)}\right)_{b\lfloor N / B\rfloor+1}, b=1, \ldots, B-1,
$$

as possible threshold candidates for that feature.
Recall that an inner node of a decision tree consists of a feature-index $j$ and a threshold-value $t$, and that the evaluation of such a node for sample $x$ is given by $x_{j}<t$.

As such, if we assume that all elements in feature column $X^{(j)}$ are unique, we find that we have

$$
\begin{equation*}
\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)=X^{(j)}<\boldsymbol{t}_{b}^{(j)}, \forall b=1, \ldots, B-1 \tag{5}
\end{equation*}
$$

where we define $\mathrm{BV}_{b}$ as the $b$-th bucket vector, such that

$$
\left(\mathrm{BV}_{b}\right)_{i}:=\left\{\begin{array}{l}
1 \text { if } i \leq b\lfloor N / B\rfloor  \tag{6}\\
0 \text { otherwise, }
\end{array} \quad \forall i=1, \ldots N\right.
$$

We call $\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)$ the selector vector of the $b$-th bucket for feature $j$, because it selects the first $b$ bins of the histogram of feature vector $X^{(j)}$.

During training, it is convenient to keep a record of the path taken by the samples in the training data. For this, we require all $(B-1) k$ selector vectors $\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)$ (c.f. section 4.2.1).

### 3.2 Bucket vectors and permutations

Naively, one would compute the selector vectors $\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)$ for feature $j$ and buckets $b=1, \ldots, B-1$, with $B-1$ calls to apply_permutation. Yet, it is possible to do this with only $\log _{2} B$ ones. This optimization is a major contribution of the paper and leads to practical speedups.

For this, we first explain how to construct the bucket vectors $\mathrm{BV}_{1}, \ldots, \mathrm{BV}_{B-1}$ from a set of publicly known generating vectors via Algorithm 1. We then leverage the commutative property between algorithm 1 and the function apply_permutation, to achieve an efficient generation of selector vectors in algorithm 2.
3.2.1 Generating bucket vectors For simplicity of the exposition and without loss of generality, we assume that $B$ is a power of two and divisor of $N$, thus each bin of the histogram holds $N / B$ samples.

First we define a helper function: for any $j \in\left[1, \log _{2} B\right]$, and any $b \in[0, B-1]$, we define $[b]_{j}$ as the $j$-th significant bit of the binary expansion of $b$. Specifically, we use 1-indexing for the bits and have

$$
b=\sum_{j=1}^{\log _{2} B}[b]_{j} 2^{j-1}, \forall b=0, \ldots, B-1
$$

We now define a binary matrix $C^{\prime}$ of dimension $B \times \log _{2} B$, such that row $i$ corresponds to the binary expansion of $B-i$, with the least-significant bit on the left, and the most significant bit on the right, e.g. we have

$$
\begin{aligned}
C_{i, j}^{\prime} & :=[B-i]_{j}, j=1, \ldots, \log _{2} B, \forall i=1, \ldots, B \\
& =\neg[i-1]_{j}, j=1, \ldots, \log _{2} B, \forall i=1, \ldots, B
\end{aligned}
$$

Further we construct the $N \times \log _{2}(B)$ binary matrix $C$ from $C^{\prime}$, by repeating each row $N / B$ times. Its columns are the generating vectors in algorithm 1:

$$
\begin{align*}
C_{i}^{(j)} & :=C_{1+\lfloor(i-1) B / N\rfloor, j}^{\prime}, & & j=1, \ldots, \log _{2} B, i \in[1, N] .  \tag{7}\\
& =[B-1-\lfloor(i-1) B / N\rfloor]_{j}, & & j=1, \ldots, \log _{2} B, i \in[1, N] .  \tag{8}\\
& =\neg[\lfloor(i-1) B / N\rfloor]_{j}, & & j=1, \ldots, \log _{2} B, i \in[1, N] . \tag{9}
\end{align*}
$$

For example, if $N=32$ and $B=8$ we have

$$
C^{t}=\left(\begin{array}{l}
111100001111000011110000111110000 \\
11111111000000001111111100000000 \\
111111111111111100000000000000000
\end{array}\right)
$$

```
Algorithm 1 bucket_vector
Input: This algorithm generates a bucket vector and takes as input
    - the public bucket index \(b \in[1, B-1]\)
    - The generating vectors \(\left\{C^{(1)}, \ldots, C^{\left(\log _{2} B\right)}\right\}\) as defined in equation 7
Output: \(b\) th bucket vector \(\mathrm{BV}_{b}\)
    res \(=0_{N}(\) the zero vector of size \(N)\)
    for \(j=1, \ldots, \log _{2} B\) do
        res \(=[b]_{j} ?\) res \(\vee C^{(j)}:\) res \(\wedge C^{(j)}\)
    end for
    return res
```

Lemma 2. Algorithm 1 yields the bth bucket vector $\mathrm{BV}_{b}$.
Proof. Let $b \in[1, B-1]$ an input for algorithm 1 and let res $_{j}$ be the state of variable res in the $j$-th iteration of the for-loop on line 3 of algorithm 1 , with $\operatorname{res}_{0}=0_{N}$ the initial value. Now, let $i \in[1, N]$ and let $b_{i} \in[0, B-1], r_{i} \in$ $[0, N / B-1]$, such that

$$
i-1=b_{i} \cdot N / B+r_{i}
$$

By definition of BV (c.f. equation 6), it suffices to prove that we have

$$
\begin{equation*}
\left(\operatorname{res}_{\log _{2} B}\right)_{i}=1 \Longleftrightarrow b_{i}<b \tag{10}
\end{equation*}
$$

Recall that by definition (c.f. equation 7 ), for any $j \in\left[1, \log _{2} B\right]$, we have

$$
\begin{equation*}
C_{i}^{(j)}=\neg\left[b_{i}\right]_{j} . \tag{11}
\end{equation*}
$$

Substituting (11) in line 3 of the algorithm yields

$$
\left(\operatorname{res}_{j}\right)_{i}=\left\{\begin{array}{l}
\left(\operatorname{res}_{j-1}\right)_{i} \text { OR } \neg\left[b_{i}\right]_{j} \text { if }[b]_{j}=1  \tag{12}\\
\left(\operatorname{res}_{j-1}\right)_{i} \text { AND } \neg\left[b_{i}\right]_{j} \text { else. }
\end{array}\right.
$$

Note that whenever $\left[b_{i}\right]_{j}=[b]_{j}$, we can substitute in equation 12 :

$$
\left(\operatorname{res}_{j}\right)_{i}=\left(\operatorname{res}_{j-1}\right)_{i} .
$$

Hence, if we define $j_{m}$ as the maximal index such that $\left[b_{i}\right]_{j_{m}} \neq[b]_{j_{m}}$, with the convention of $j_{m}=0$, if $b_{i}=b$, we find that

$$
\begin{equation*}
\left(\operatorname{res}_{\log _{2} B}\right)_{i}=\left(\operatorname{res}_{j_{m}}\right)_{i} \tag{13}
\end{equation*}
$$

Finally, we can observe that equation 10 holds for each of the three possible cases and conclude the proof:
case $b_{i}>b$ : we find $1=\left[b_{i}\right]_{j_{m}}>[b]_{j_{m}}=0$ and thus

$$
\left(\operatorname{res}_{\log _{2} B}\right)_{i} \stackrel{(13)}{=}\left(\operatorname{res}_{j_{m}}\right)_{i} \stackrel{(12)}{=}\left(\operatorname{res}_{j_{m}-1}\right)_{i} \text { AND } 0=0
$$

case $b_{i}=b$ : we have

$$
\left(\mathrm{res}_{\log _{2} B}\right)_{i} \stackrel{(13)}{=}\left(\operatorname{res}_{0}\right)_{i}=\left(0_{N}\right)_{i}=0
$$

case $b_{i}<b$ : we find $0=\left[b_{i}\right]_{j_{m}}<[b]_{j_{m}}=1$ and thus

$$
\left(\operatorname{res}_{\log _{2} B}\right)_{i} \stackrel{(13)}{=}\left(\operatorname{res}_{j_{m}}\right)_{i} \stackrel{(12)}{=}\left(\operatorname{res}_{j_{m}-1}\right)_{i} \text { OR } 1=1
$$

3.2.2 Constructing Selector Vectors Since algorithm 1 only requires AND and OR operations on the generating vectors $C^{(1)}, \ldots, C^{\left(\log _{2} B\right)}$, we find that the function apply_permutation and algorithm 1 are commuting, i.e. more specifically, for each sorting permutation $\pi_{j}$, we have that the selector vector $\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)$ is given by

$$
\begin{aligned}
\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right) & =\pi_{j}^{-1}\left(\text { bucket_vector }\left(b,\left\{C^{(1)}, \ldots, C^{\left(\log _{2} B\right)}\right\}\right)\right) \\
& =\text { bucket_vector }\left(b,\left\{\pi_{j}^{-1}\left(C^{(1)}\right), \ldots, \pi_{j}^{-1}\left(C^{\left(\log _{2} B\right)}\right)\right\}\right)
\end{aligned}
$$

Hence, if we define

$$
C_{j, m}:=\pi_{j}^{-1}\left(C^{(m)}\right), j=1, \ldots, k, m=1, \ldots, \log _{2}(B)
$$

We can re-construct all $(B-1) k$ selection vectors via Algorithm 2 and only require $\log _{2}(B) k$ calls to apply_permutation.

```
Algorithm 2 selector_vector
Input: This algorithm takes as input:
    - Public bucket index \(b \in[1, B-1]\)
    - Secret-shared generating vectors \(\left\{C_{j, 1}, \ldots, C_{j, \log _{2} B}\right\}\) for feature \(j\)
Output: Secret-shared selection-vector \(\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)\) for the first \(b\) bins of feature \(j\)
    res \(=0_{N}(\) the zero vector of size \(N)\)
    for \(m=1, \ldots, \log _{2} B\) do
            res \(:=[b]_{m}\) ? res \(\vee C_{j, m}: \operatorname{res} \wedge C_{j, m}\)
    end for
    return res
```


## 4 Description of the XORBoost Training Algorithm

The input to the training algorithm is the feature matrix $X$ as well as the response vector $y$ (both secret shared among the players).

To train an ensemble of a fixed number $T$ of binary decision trees, we proceed as follows: assuming that we have already trained the first $t-1$ trees, to grow
the $t$ th tree to a given depth $D$, we iterate by layers starting from layer zero, that is, the root node. In each iteration, we 'split' each leaf into a left and a right child via a splitting criterion, that is, a pair of a feature index and a threshold value, maximizing the gain.

For efficiency reasons, we have made the possible threshold values discrete by introducing the histogram/buckets in the data preprocessing phase, i.e., obliviously sorting each feature as described in Section 3.1 and building the histograms of Section 3.1.1.

To compute the optimal splitting criterion in plaintext, we simply iterate over all possible splitting criteria (all feature indices and all bucket indices) and select the optimal one. This yields a $(B-1) \times k$ gain matrix (of all possible gain values).

To adapt this simple optimization procedure to the MPC setting, we compute the optimal splitting criteria obliviously by first computing the gain matrix obliviously and then computing (also obliviously) the (secret shared) feature selector e and (secret shared) threshold selector $\mathbf{t}$ corresponding to the feature index $j$ and the bucket index $b$ of the optimal splitting criterion. Subsequently, pick the right generator vectors $C_{j, m}$ for $\mathrm{BV}_{b}$ and compute (again obliviously) the preimage $\pi_{j}^{-1}\left(\mathrm{BV}_{b}\right)$ of the bucket vector. This allows us to decide which samples would go to the left child and to the right child, effectively splitting the node. We now go into more detail for each of these steps.

```
Algorithm 3 xorboost_train
Input: The training algorithm takes the following input:
    - X - feature matrix of size \(N \times k\) (secret shared)
    - \(y\) - response vector of size \(N\) (secret shared)
    - \(T\) - the size of the ensemble
    - \(D\) - the depth of each binary decision tree
Output: A tree ensemble ensemble \(=\left\{\operatorname{Tree}^{(1)}, \ldots\right.\), Tree \(\left.^{(T)}\right\}\). Each tree Tree \({ }^{(t)}\) con-
    sists of the following data:
        \(-\mathbf{t}_{n}^{(t)}\) - secret shared threshold selector for each non-leaf node \(n\)
        - \(\mathbf{e}_{n}^{(t)}\) - secret shared feature selector for each non-leaf node \(n\)
        - \(w_{\ell}^{(t)}\) - secret shared weight for each leaf node \(\ell\)
        - \(\operatorname{Tr} e^{(t)}(X)\) - secret shared vector of predictions on the training data
    \(\widehat{y}^{0}:=\) initialize \((X, y)\)
    ensemble \(=\{ \}\)
    for \(t=1, \ldots, T\) do
        \(g^{(t-1)}:=\left(\partial_{b} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(t-1)}\right)\right)_{i \in[1, N]}\)
        \(h^{(t-1)}:=\left(\partial_{b}^{2} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(t-1))}\right)_{i \in[1, N]}\right.\)
        \(\operatorname{Tree}{ }^{(t)}:=\) grow_tree \(\left(g^{(t-1)}, h^{(t-1)}\right)\)
        Add \(\operatorname{Tree}{ }^{(t)}\) to ensemble
    end for
    return ensemble
```


### 4.1 Computing initial predictions

The computation of the first and second order statistics vectors $g$ and $h$ of the loss function only depends on the response variable $y$ and the current estimate $\widehat{y}^{(t)}$. Since the tree ensemble is initially empty, we must provide an initial estimate $\widehat{y}^{0}$ in order to grow the first tree. There are several possibilities for the initialization of $\widehat{y}^{0}$ :

- The zero vector.
- The constant vector with value $\alpha$ minimizing $\sum_{i=1}^{N} \operatorname{loss}\left(y_{i}, \alpha\right)$. For instance, for $L 2$ loss, this corresponds to $\alpha=\frac{1}{N} \sum_{i=1}^{N} y_{i}$, and for logistic loss this corresponds to $\alpha=\sigma^{-1}\left(\frac{1}{N} \sum_{i=1}^{N} y_{i}\right)$, where $\sigma$ is the sigmoid function $\sigma(x)=$ $\frac{1}{1+e^{-x}}$.
- Leveraging previous work on ridge regression (respectively logistic regression) [5], we can use its prediction to initialize the boosted trees model in the case of $L 2$ loss (respectively logistic loss). The aim here is to bootstrap the gradient boosting procedure by starting with a better initial value for $\widehat{y}^{0}$ and reducing the number of trees required to obtain a model with good predictive power.
In all cases, we assume that we have defined a function initialize $(X, y)$ that will compute the initial predictions.


### 4.2 Oblivious permutations and computing gain matrices

We now explain how to efficiently apply oblivious permutations in order to compute gains and weights. Recall from Section 3.1 that we have obliviously sorted the feature columns of $X$ by a set $\Pi=\left(\pi_{1}, \ldots, \pi_{k}\right)$ of $k$ secret-shared sorting permutations. The computation of the weights and the gain matrices includes two types of secret-shared vectors: instance vectors and the already introduced bucket vectors from Section 3.2.
4.2.1 Efficiently computing instance vectors. The main idea lies behind manipulating secret-shares of the so-called instance vectors. Associated to each node $n \in \mathcal{N}$ is an instance vector $\mathrm{IV}_{n}$, that is, the binary vector of size $N$ indicating which samples of the dataset $X$ go through this node. To parallelize computations and reduce complexity, we also store the permuted instance matrix $\Pi\left(\mathrm{IV}_{n}\right)$ - this is an $N \times k$ matrix.

When growing the tree, we will be working per layer. For a given depth $d$, let $L_{d}$ be the $d$ th layer, that is, the set of the $2^{d}$ nodes at depth $d$ ( $L_{0}$ consists of the Root only, $L_{1}$ consists of the two children of the Root, etc.). Observe the following two basic properties:

1. $\sum_{n \in L_{d}} \mathrm{IV}_{n}=1_{N}$,
2. $\mathrm{IV}_{n}=\mathrm{IV}_{n_{\text {left }}}+I V_{n_{\text {right }}}$.

Since applying the function apply_permutation to a vector of size $N$ is costly, we use an optimization technique to reduce the number of calls to this function from $O\left(2^{D}\right)$ to $O(D)$ (i.e., linear on the depth of the tree). More precisely, we define the compressed instance vectors IVC

$$
\operatorname{IVC}_{d+1}:=\bigoplus_{n \in L_{d}} \operatorname{IV}_{n_{\text {left }}}
$$

where $\bigoplus$ is the logical XOR operator. Using Property 1. above, we see that the $\oplus$ operation in the definition is equivalent to $\vee$. This allows us to compute the oblivious permutation matrix $\Pi\left(\mathrm{IV}_{n_{\text {left }}}\right)$ for the left child of a given node by simply applying the AND operator between the permutation matrix for that node with the compressed instance vector for that level, i.e., by combining the identity

$$
\begin{equation*}
\mathrm{IV}_{n_{\text {left }}}=I V_{n} \wedge \mathrm{IVC}_{d+1} \tag{14}
\end{equation*}
$$

with the commutativity of permutations and logical operations. Thus,

$$
\begin{equation*}
\Pi\left(\mathrm{IV}_{n_{\text {left }}}\right)=\Pi\left(\mathrm{IV}_{n}\right) \wedge \Pi\left(\mathrm{IVC}_{d+1}\right) \tag{15}
\end{equation*}
$$

This explains the interest in IVC: instead of applying $\Pi$ to each $\mathrm{IV}_{n_{\text {left }}}$, we can apply $\Pi$ to $\operatorname{IVC}_{d+1}$ and compute $\Pi\left(\mathrm{IV}_{n_{\text {left }}}\right)$ using a significantly cheaper $\wedge$ operation.
4.2.2 Computing gain matrices. Recall that computing gain (eqs. (4)) requires computing the quantities $G_{n}$ and $H_{n}$ for various nodes $n$ and their left and right children. These can conveniently be written in terms of instance vectors and sorting permutations as follows:

$$
\begin{equation*}
G_{n}=\pi(g)^{t} \cdot \pi\left(\mathrm{IV}_{n}\right) \quad \text { and } \quad H_{n}=\pi(h)^{t} \cdot \pi\left(\mathrm{IV}_{n}\right) \tag{16}
\end{equation*}
$$

where $\pi$ is any permutation on $N$ letters (in particular, any of the sorting permutation). Note that this expression does not depend on the choice of $\pi$.

We are now ready to compute the gain function to be used in the tree growing algorithm. Since our goal is to bucket each feature value into one of the $B$ possible buckets, we only want to compute the gains corresponding to $B-1$ splittings at the bucket thresholds (thus, discretizing the possible splits, as explained in Section 3.1.1). As we do this for any of the $k$ input features, we can conveniently package these gains into a $(B-1) \times k$ matrix gains that can be computed via (16) as

$$
\begin{aligned}
\operatorname{gains}_{n}(i, j) & =\operatorname{score}\left(\mathrm{G}_{n_{\text {left }}}(i, j), \mathrm{H}_{n_{\text {left }}}(i, j), \lambda\right)+ \\
& +\operatorname{score}\left(\mathrm{G}_{n_{\text {right }}}(i, j), \mathrm{H}_{n_{\text {right }}}(i, j), \lambda\right)- \\
& -\operatorname{score}\left(\mathrm{G}_{n_{\text {left }}}(i, j)+\mathrm{G}_{n_{\text {right }}}(i, j), \mathrm{H}_{n_{\text {left }}}(i, j)+\mathrm{H}_{n_{\text {left }}}(i, j), \lambda\right)-\gamma,
\end{aligned}
$$

for $i=1, \ldots, B-1$ and $j=1, \ldots, k$, where

$$
\mathrm{G}_{n_{\mathrm{left}}}=\left(\Pi(g) \odot \Pi\left(\mathrm{I} \mathrm{~V}_{n}\right)\right)^{t} \cdot \mathrm{BM}, \quad \mathrm{G}_{n_{\mathrm{right}}}=\left(\Pi(g) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \neg \mathrm{BM}
$$

and

$$
\mathrm{H}_{n_{\text {left }}}=\left(\Pi(h) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \mathrm{BM}, \quad \mathrm{H}_{n_{\mathrm{right}}}=\left(\Pi(h) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \neg \mathrm{BM}
$$

where $\odot$ denotes the Hadamard product (i.e., elementwise product) and BM is the concatenation of the bucket vectors: $\mathrm{BM}=\left[\mathrm{BV}_{1}|\ldots| \mathrm{BV}_{B-1}\right]$ (a $N$-by- $(B-1)$ matrix). See Section 2 for the definition of the score function and an explanation of the gain terminology.

```
Algorithm 4 gain_matrix
Input: This algorithm takes as input
    - \(N \times k\) matrices \(\Pi(g)\) and \(\Pi(h)\) for the first- and second-order statistics vectors
    - Permutation matrix \(\Pi\left(\mathrm{IV}_{n}\right)\) for node \(n\)
Output: The \((B-1) \times k\) gain matrix gains \({ }_{n}\) for node \(n\)
    \(\mathrm{G}_{n_{\text {left }}}=\left(\Pi(g) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \mathrm{BM}\)
    \(\mathrm{G}_{n_{\text {right }}}=\left(\Pi(g) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \neg \mathrm{BM}\)
    \(\mathrm{H}_{n_{\text {left }}}=\left(\Pi(h) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \mathrm{BM}\)
    \(\mathrm{H}_{n_{\text {right }}}=\left(\Pi(h) \odot \Pi\left(\mathrm{IV}_{n}\right)\right)^{t} \cdot \neg \mathrm{BM}\)
    return \(\operatorname{score}\left(\mathrm{G}_{n_{\text {left }}}, \mathrm{H}_{n_{\text {left }}}, \lambda\right)+\operatorname{score}\left(\mathrm{G}_{n_{\text {right }}}, \mathrm{H}_{n_{\text {right }}}, \lambda\right)-\operatorname{score}\left(\mathrm{G}_{n_{\text {left }}}+\right.\)
    \(\left.\mathrm{G}_{n_{\text {right }}}, \mathrm{H}_{n_{\text {left }}}+\mathrm{H}_{n_{\text {right }}}, \lambda\right)-\gamma\)
```

Algorithm 4 summarizes the above computation. Note that in Step 5, we have overloaded notation so that $\operatorname{score}\left(\mathrm{G}_{n_{\text {left }}}, \mathrm{H}_{n_{\text {left }}}, \lambda\right)$ is now a $(B-1) \times k$ matrix.

Remark 1. By tracking a 0-1 matrix $X_{\text {nan }}$ of missing values in the initial dataset, we can adapt the algorithm to handle missing values. Let

$$
\Pi\left(X_{n a n}\right)=\left[\pi_{1}\left(X^{(1)}\right), \ldots, \pi_{k}\left(X^{(k)}\right)\right] .
$$

We can consider $\Pi\left(\mathrm{IV}_{n, n a n}\right)=\Pi\left(\mathrm{IV}_{n}\right)$ XOR $\Pi\left(X_{\text {nan }}\right)$. We then obtain two gain matrices, one where $G_{n_{\text {left }}}$ and $\mathrm{H}_{n_{\text {left }}}$ are computed using $\Pi\left(\mathrm{IV}_{n, \text { nan }}\right)$ instead of $\Pi\left(\mathrm{IV}_{n}\right)$ and another where it is $\mathrm{G}_{n_{\text {right }}}$ and $\mathrm{H}_{n_{\text {right }}}$ that use $\Pi\left(\mathrm{IV}_{n, n a n}\right)$. We can then proceed to take the argmax over the concatenation of these two matrices.

### 4.3 Tree growing

The algorithm will split the layers until the desired depth $D$ and compute the leaf weights to obtain the predictions $\operatorname{Tree}(X)$. The threshold selectors $\mathbf{t}_{n}$ and feature selectors $e_{n}$ are computed during the layer splitting while the weights $w=\left(w_{\ell}\right)_{\ell \in L_{D}}$ are computed once the tree has reached its full depth. We denote by $\mathbf{T}=\left\{\mathbf{t}_{n}: n \in\left[1,2^{D}-1\right]\right\}$ the set of thresholds of Tree and by $\mathbf{E}=\left\{\mathrm{e}_{n}: n \in\right.$
$\left.\left[1,2^{D}-1\right]\right\}$ the set of feature selectors. The data $\{\mathbf{T}, \mathbf{E}\}$ constitutes the tree structure and is secret-shared.

```
Algorithm 5 grow_tree
Input: As input, this algorithm takes
    - First and second order statistic vectors \(g\) and \(h\)
    - Depth \(D\) of the tree to grow
Output: A tree \(\{\mathbf{T}, \mathbf{E}, w\}\)
    \(\Pi(g) \leftarrow\) apply_permutation \((g)\)
2: \(\Pi(h) \leftarrow\) apply_permutation \((h)\)
3: \(\mathrm{IV}_{\text {Root }} \leftarrow 1_{N \times 1}\)
4: \(\Pi\left(\mathrm{IV}_{\text {Root }}\right) \leftarrow 1_{N \times k}\)
5: \(\mathbf{T}, \mathbf{E},\left\{\mathrm{IV}_{\ell}: \ell \in L_{D}\right\} \leftarrow \operatorname{split\_ layer}\left(1, \mathbf{T}=\emptyset, \mathbf{E}=\emptyset, \Pi(g), \Pi(h),\left\{\mathrm{IV}_{\text {Root }}\right\},\left\{\Pi\left(\mathrm{IV}_{\text {Root }}\right)\right\}, D\right)\)
6: for \(\ell \in L_{D}\) do
            \(w_{\ell} \leftarrow\) privatedivide \(\left(-g^{T} \cdot \mathrm{IV}_{\ell}, h^{T} \cdot \mathrm{IV}_{\ell}+\lambda\right)\)
    end for
    \(\operatorname{Tree}(X) \leftarrow \sum_{\ell \in L_{D}} w_{\ell} \cdot \operatorname{IV}_{\ell}\)
    return \(\mathbf{T}, \mathbf{E}, w=\left(w_{\ell}\right)_{\ell \in L_{D}}, \operatorname{Tree}(X)\)
```

Remark 2. For each split, one can track the value of the gain achieved. By verifying that the gain is positive using a private compare function, one can prune the tree of leaves which hold little explanatory power. This means computing the weight for every internal node and updating the weight of a leaf if its ancestor has negative gain. This can be determined using a private compare function.

### 4.4 Splitting a layer

We present a detailed description of the algorithm that, given the $d$ th layer of a tree, adds a new layer to a tree by splitting each node at depth $d$ into a left and right child node at depth $d+1$. Note that all the output data is secret-shared.

```
Algorithm 6 split_layer at depth \(d\)
Input: \(d\) - current depth, \(\mathbf{T}, \mathbf{E}, \Pi(g), \Pi(h),\left\{\mathrm{IV}_{n}: n \in L_{d}\right\},\left\{\Pi\left(\mathrm{IV}_{n}\right): n \in L_{d}\right\}, D-\)
    final depth
Output: T, E \(,\left\{\mathrm{IV}_{n}: n \in L_{d+1}\right\},\left\{\Pi\left(\mathrm{IV}_{n}\right): n \in L_{d+1}\right\}\)
    for \(n \in L_{d}\) do
        gains \(_{n} \leftarrow\) gain_matrix \(\left(\Pi(g), \Pi(h), \Pi\left(\mathrm{IV}_{n}\right)\right) \quad \triangleright\) Matrix of size \((B-1) \times k\)
        \(\mathbf{t}_{n}, \mathbf{e}_{n}, b, j \leftarrow \arg \max \left(\right.\) gains \(\left._{n}\right)\)
        Selector \(\leftarrow\) selector_vector \(\left(\left\{C_{m, j}: 1 \leq m \leq \log _{2} B\right\}, b\right)\)
        \(\mathrm{IV}_{n_{\text {left }}} \leftarrow\) Selector \(\wedge\) IV \(_{n}\)
        \(\mathrm{IV}_{n_{\text {right }}} \leftarrow \neg\) Selector \(\wedge\) IV \(_{n}\)
    end for
    \(\mathbf{T} \leftarrow \mathbf{T} \cup\left\{\mathbf{t}_{n}: n \in L_{d}\right\}, \mathbf{E} \leftarrow \mathbf{E} \cup\left\{\mathbf{e}_{n}: n \in L_{d}\right\}\)
    \(\left\{\mathrm{IV}_{n}: n \in L_{d+1}\right\} \leftarrow\left\{\mathrm{IV}_{n_{\text {left }}}: n \in L_{d}\right\} \cup\left\{\operatorname{IV}_{n_{\text {right }}}: n \in L_{d}\right\}\)
    \(\operatorname{IVC} \leftarrow \bigoplus_{n \in L_{d}}\) IV \(_{n_{\text {left }}}\)
    \(\Pi(\) IVC \() \leftarrow\) apply_permutation(IVC)
    for \(n \in L_{d}\) do
        \(\Pi\left(\mathrm{IV}_{n_{\text {left }}}\right) \leftarrow \Pi(\mathrm{IVC}) \wedge \Pi\left(\mathrm{IV}_{n}\right)\)
        \(\Pi\left(\mathrm{IV}_{n_{\mathrm{right}}}\right) \leftarrow \neg \Pi(\mathrm{IVC}) \wedge \Pi\left(\mathrm{IV}_{n}\right)\)
    end for
    \(\left\{\Pi\left(\mathrm{IV}_{n}\right): n \in L_{d+1}\right\} \leftarrow\left\{\Pi\left(\mathrm{IV}_{n_{\text {left }}}\right): n \in L_{d}\right\} \cup\left\{\Pi\left(\mathrm{IV}_{n_{\text {right }}}\right): n \in L_{d}\right\}\)
    if \(d=D-1\) then
        return \(\mathbf{T}, \mathbf{E},\left\{\mathrm{IV}_{n}: n \in L_{D}\right\}\)
    else
        return split_layer \(\left(d+1, \Pi(g), \Pi(h),\left\{\mathrm{IV}_{n}: n \in L_{d+1}\right\},\left\{\Pi(\mathrm{IV})_{n}: n \in L_{d+1}\right\}, D\right)\)
    end if
```

Note that the argmax only returns the secret-shared vectors $\mathbf{t}$ and e from which we could compute $b$ and $j$. This is purely for exposition since we can obtain the right generator vectors $C_{j, m}$ using e and matrix multiplication.

## 5 Prediction

Our inference protocol respects the same privacy requirements as our training protocol: only the shape of the model is public knowledge (tree depth, number of trees) everything else remains secret (threshold values, feature selectors, leaf- and prediction-values). The data used to evaluate the model remains secret. Similar to [2] we achieve this by evaluating the predicate of each node in the tree and thus hiding the path taken. Previous work on private and secure decision tree inference was done in [9] and [11] for passive and active security in two player settings.

In order to improve efficiency, the communication intensive comparisons and multiplications required for securely evaluating a tree are batched in Algorithm 7. We give a brief explanation of the protocol for a batch of size one, i.e. a single tree.

Recall that each non-leaf node consists of a feature selector and a threshold value. To evaluate a non-leaf node with a sample $x$, the feature selector is used
to extract the feature of interest from $x$, which is then compared against the threshold value. This can be achieved securely via a multiplication of the secretshared feature selector and the secret-shared sample $x$, followed by an oblivious comparison with the threshold value, yielding a secret-shared Boolean that indicates, if the left or the right subtree is to be evaluated next. Our iterative algorithm 7 evaluates one layer at a time - starting from the bottom layer and going up until reaching the root of the tree, each time obliviously selecting between secret-shared leaf-values and reducing the number of leaf candidates by a factor of two. Note how multiplications and comparisons on lines one and two of algorithm 7 have been batched to reduce the number of communication rounds.

```
Algorithm 7 xorboost_predict_batch
Input: The prediction function for a batch of trees takes as input:
    - secret-shared data matrix \(X^{\prime}\) of dimensions \(N^{\prime} \times k\)
    - secret-shared tree ensemble of \(T\) trees, each tree Tree \({ }^{(t)}\) of depth \(D\) and con-
        sisting of
            - threshold values \(\left\{\mathbf{t}_{n}^{t}: n \in\left[1,2^{D}-1\right]\right\}\)
            - feature selectors \(\left\{\mathrm{e}_{n}^{t}: n \in\left[1,2^{D}-1\right]\right\}\) : vectors of length \(k\) each
            - prediction weights \(w^{t}=\left(w_{\ell}\right)_{\ell \in\left[1,2^{D}\right]}\) (leaf-values)
Output: Predictions Tree \({ }^{(t)}\left(X^{\prime}\right)\) for \(t \in[1, T]\)
```

    Extract the feature-values and compare them to the threshold values:
    \(X^{\prime(t, n)}:=X^{\prime} \cdot \mathrm{e}_{n}^{t}\) for all \(n \in\left[1,2^{D}-1\right]\) and \(t \in[1, T]\)
    \(\beta^{t, n}:=X^{\prime(t, n)}<\mathbf{t}_{n}^{t}\), for all \(n \in\left[1,2^{D}-1\right]\) and \(t \in[1, T]\)
    Re-arrange the leaf values and obliviously extract the correct candidate:
    \(w^{t, 2^{D-1}+\ell}=w_{\ell}^{t} \cdot \mathbb{1}_{N^{\prime} \times 1}, \ell=1, \ldots, 2^{D}\), for all \(t \in[1, T]\)
    for \(d=D, \ldots, 1\) do
            \(w^{t, n}=w^{t, 2 n+1}+\beta^{t, n} \cdot\left(w^{t, 2 n}-w^{t, 2 n+1}\right)\) for \(n=1, \ldots, 2^{d-1}\) and \(t \in[1, T]\)
    end for
    return \(\left\{w^{t, 1}: t \in[1, T]\right\}\)
    Evaluating a gradient boosted tree model follows naturally: The tree-ensemble is split into batches of reasonable sizes, which are evaluated independently. The predictions for each batch are then multiplied with the learning parameter $\eta$ (c.f. section 2) and summed together with the initial prediction (c.f. section 4.1) as outlined in algorithm 8.

```
Algorithm 8 xorboost_predict
Input: The prediction function for the entire model takes as input:
    - secret-shared data \(N^{\prime} \times k\) matrix \(X^{\prime}\),
    - Learning rate parameter \(\eta\)
    - secret-shared initial prediction value \(\hat{y}^{0}\)
    - secret-shared tree ensemble of \(T\) trees, each tree Tree \({ }^{(t)}\) of depth \(D\) and con-
                sisting of
            - threshold values \(\left\{\mathbf{t}_{n}^{t}: n \in\left[1,2^{D}-1\right]\right\}\)
            - feature selectors \(\left\{\mathrm{e}_{n}^{t}: n \in\left[1,2^{D}-1\right]\right\}\) : vectors of length \(k\) each
            - prediction weights \(w^{t}=\left(w_{\ell}\right)_{\ell \in\left[1,2^{D}\right]}\) (leaf-values)
Output: secret-shared prediction vector \(\hat{y}:=\hat{y}^{0}+\sum_{t=1}^{T} \eta \operatorname{Tree}{ }^{(t)}\left(X^{\prime}\right)\)
    \(w^{1}, \ldots, w^{T}:=\) xorboost_predict_batch \(\left(X^{\prime}, \operatorname{Tree}^{(1)}, \ldots\right.\), Tree \(\left.^{(T)}\right)\)
    \(\hat{y}:=\hat{y}^{0}+\eta \sum_{t=1}^{T} w^{t}\)
    return \(\hat{y}\)
```


## 6 Categorical feature

Many datasets include categorical features which are informative and should be handled adequately by the machine learning model. While the original xgboost approach does not support categorical features, there has been development in this area. For a survey of different ways to do this, see [23]. We develop here an approach that is compatible with the work presented so far. Suppose given a feature column $X^{(j)}$ with categorical values drawn from a set $A$. Tree splits for this feature are induced by writing $A$ as a disjoint union, $A=A_{1} \bigcup A_{2}$ and we are interested in finding the split that produces the greatest reduction in loss. By the main result of [8], the following procedure will find the best split if the loss function is either $L 2$ loss or logistic loss. It only considers $|A|-1$ possible splits.

- For each categorical value $a \in A$, compute $y_{a}=\frac{\sum_{i, X_{i}^{(j)}=a} y_{i}}{\sum_{i, X_{i}^{(j)}=a} 1}$ the average response variable over all samples for $a$.
- Order the categorical values according to the $y_{a}$ values, yielding a total order $<$ (breaking ties arbitrarily) on the set $A$.
- For every $a \in A$ consider the split $S_{a}$ with $A_{1}=\left\{a^{\prime} \mid a^{\prime} \leq a\right\}$ and $A_{2}=$ $\left\{a^{\prime} \mid a^{\prime}>a\right\}$.

Let $\pi_{A}$ be the sorting permutation for the values $y_{a}$. Given the original one hot encoded matrix $M$, we apply $\pi_{A}$ to its columns to obtain $M_{A}$, the matrix where the columns are sorted according to the values of $y_{a}$. In the context of categorical variables, the bucket vectors' role is played by indicator vectors for the splits $S_{a}$. These indicator vectors for the splits $S_{a}$ can be computed as row sums of the sorted matrix $M_{A}$. If we let $a$ be the $b$ th categorical value according to the total order $<$, we have:

$$
\mathrm{BV}_{b}=\sum_{j \leq b} M_{A}^{(j)}=\left\{\begin{array}{l}
1 \text { if } X^{(j)}=a^{\prime}, a^{\prime} \leq a \\
0 \text { otherwise, }
\end{array} \quad \forall i=1, \ldots N\right.
$$

## 7 Oblivious Trees

Oblivious trees use the same split for every node across a given layer, see for example [23]. This means that the splits are no longer adaptive: they do not depend on the instance sets at a given node. Concretely, the tree growing step would proceed as follows, assuming the first and second order statistic vectors have already been computed:

- Compute the gain matrix. Let $\alpha_{1} \geq \ldots \geq \alpha_{D}$ be the $D$ entries with maximum gain.
- For each $\alpha_{d}$, compute $\mathbf{t}_{d}$ and $\mathrm{e}_{d}$.
- Compute the leaf instance vectors and leaf weights.

In terms of complexity, the main difference is that a single gain matrix is computed for the whole tree as opposed to one gain matrix per internal node in the original approach. This also means reducing the number of calls to the (expensive) argmax function. The tradeoff is in the inference step: since the splits are chosen all at once without taking into account previous splits, the reduction in loss at each tree split stands to be lower than in the original fully adaptive version of XGBoost. Oblivious trees can easily be adapted into the framework developed so far in order to provide a tradeoff between inferential power and computational complexity.

## 8 Benchmarks

All benchmarks have been done on a single n1-standard-8 (8 vCPUs, 30GB of RAM, SSD drive, Intel Xeon CPU Skylake 2.00 GHz ) Google Compute Engine virtual machine for 2 players. As such, all reported times do not take into account network transfer time.

### 8.1 Parameter Scaling

In figure 1, we show the impact of varying the dataset size for the utilization of network, memory and time. In figure 2 , we show the impact of varying the number of buckets as well as the tree depth on the same resources.

The benchmarks presented in Table 1 and Figure 3 are for a dataset of $N=20 K$ samples and $k=300$ features, 2 players, and models with 5 trees. Tree depth and number of buckets vary according to the table rows.

Referring to Table 1 we see the strong dependency on the depth and the number of buckets. The total processing time remains reasonable to grow 5


Fig. 1: Network Size, RAM usage and Wall time for depth 4, 64 buckets, 10 trees, $N \in\{5 K, 25 K, 100 K\}$ samples and $k \in\{10,300\}$ features.
trees. Figure 3 highlights that total wall-time hovers around a few minutes per tree for depth 4 and total time grows more or less linearly with the number of trees, as expected.

If we compare our results to [7], they are similar. We obtain a running time of 22 minutes for 5 trees, dataset dimensions of 20 K by 300 , a tree depth of 4


Fig. 2: Network size, Memory and Wall time for a dataset of dimension $20 K \times 300$ and a model with 5 trees.
and 64 buckets. They obtain a running time of around 5 minutes for a single tree with $30 K$ samples, 25 features and a tree depth of 4 . Delivering similar performance while not revealing any information is a significant improvement.

Although our framework leverages a bucketing strategy to handle larger datasets, Figure 4 presents a comparison with the main benchmarks of [2] where

|  |  | RAM |  | Communication |  | Wall time |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Depth | Number of | Offline(TD) | Online(pp) | Triples(pp) | Network(pp) | Offline(TD) | Online(pp) |
|  | buckets | GB |  |  |  | minutes |  |
| 3 | 16 | 1.2 | 2.24 | 17.4 | 14.7 | 5.61 | 6.23 |
| 4 |  | 1.41 | 2.70 | 31.8 | 23.5 | 9.98 | 10.6 |
| 5 |  | 1.95 | 3.64 | 60.0 | 40.5 | 18.2 | 19.9 |
| 3 | 32 | 1.27 | 2.44 | 17.8 | 15.1 | 5.78 | 6.58 |
| 4 |  | 1.51 | 2.95 | 32.4 | 24.2 | 10.4 | 11.3 |
| 5 |  | 2.43 | 3.98 | 61.3 | 41.8 | 18.4 | 20.8 |
| 3 | 64 | 1.35 | 2.64 | 18.4 | 15.6 | 6.05 | 7.18 |
| 4 |  | 1.94 | 3.26 | 33.6 | 25.3 | 10.7 | 12.8 |
| 5 |  | 3.38 | 4.42 | 63.7 | 44.0 | 18.7 | 23.4 |
| 3 | 128 | 2.10 | 2.93 | 19.4 | 16.6 | 6.41 | 8.41 |
| 4 |  | 3.81 | 3.89 | 35.8 | 27.3 | 11.2 | 14.4 |
| 5 |  | 6.53 | 6.53 | 68.3 | 48.0 | 20.4 | 28.7 |
| 3 | 256 | 5.50 | 5.51 | 21.3 | 18.3 | 6.93 | 10.1 |
| 4 |  | 10.9 | 11.0 | 40.1 | 31.0 | 12.5 | 19.6 |
| 5 |  | NA | NA | NA | NA | NA | NA |

Table 1: Benchmarks for $20 K$ samples, 300 features, 2 players and 5 trees.
no bucketing is possible. We use a tree depth of 1 and 2 features in order to use the same parameters. For 8192 samples, end to end execution time is around 12 seconds and communication is 50 MB for XORBoost compared to 35 seconds and 3.5 GB for Table 1 in [2] for the passive security setting.

### 8.2 Division and argmax.

For each split node we need to compute a division and an argmax on a matrix of size $(B-1) \times k$ (which from a performance point of view is equivalent to computing a division/argmax on a vector of size $(B-1) \cdot k)$. Table 2 shows execution times for a 2-player division for different sizes of input vectors. Input vectors are 60 -bit wide and the output has a relative precision of 50 bits compared to the plaintext division (performed on double-precision floating-point).

Table 3 shows execution times for a 2-player argmax for different sizes of input vectors.

|  | RAM |  | Communication |  | Wall Time |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Input size | Offline(TD) | Online(pp) | Triples(pp) | Network(pp) | Offline(TD) |  |
|  | Online(pp) |  |  |  |  |  |
|  | MB |  |  | seconds |  |  |
| 512 | 8.8 | 9.4 | 0.380 | 0.290 | 0.17 | 1.65 |
| 2048 | 13.8 | 14.1 | 1.411 | 1.099 | 0.21 | 1.71 |
| 16384 | 49.8 | 46.8 | 10.662 | 8.615 | 0.29 | 1.79 |
| 65536 | 176.6 | 110.3 | 42.381 | 34.383 | 0.72 | 2.35 |

Table 2: 2-player private division execution times and communication sizes


Fig. 3: Growth of RAM usage during the online phase and of total execution time for 64 buckets and a tree depth of 4 .


Fig. 4: Resource utilization without bucketing, comparison with (AEV) [2]

|  | RAM |  | Communication |  | Wall Time |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Input size | Offline(TD) | Online(pp) | Triples(pp) | Network(pp) | Offline(TD) | Online(pp) |
|  | MB |  |  |  | seconds |  |
| 512 | 8.5 | 9.7 | 0.057 | 0.031 | 0.13 | 1.68 |
| 2048 | 10.1 | 11.1 | 0.148 | 0.147 | 0.15 | 1.71 |
| 16384 | 16.2 | 17.9 | 0.758 | 1.032 | 0.29 | 1.78 |
| 65536 | 28.8 | 34.5 | 2.801 | 4.046 | 0.42 | 1.94 |

Table 3: 2-player private argmax execution times and communication sizes

### 8.3 Comparison with plaintext algorithms

We have also ascertained that minimal predictive power is lost with respect to plaintext implementations. Since there is no unique minimum loss model, imple-
mentation decisions such as how the bucketing is performed result in different models even when comparing only plaintext models. We compared the L2 loss of the predictions made by XORboost and by several well-known plaintext implementations on the training dataset: scikit-learn (with and without bucketing), xgboost, lightgbm. We ran a simulation with 50 different datasets generated at random. If we let minLoss $s_{i}$ be the loss value achieved by the best model and maxLoss $i_{i}$ be the loss achieved by the worst model for the $i$ th dataset, then maxLoss is on average $12 \%$ higher than minLoss. On average, XORboost is $6 \%$ higher than minLoss. We conclude from this that XORBoost behaves similarly to other gradient boosting implementations with respect to predictive power.

As an example, for a particular dataset, XORBoost achieved a loss of 58.1 (lower is better), while sklearn without histograms achieved 51.5 , sklearn with histograms achieved 56.1, lightgbm achieved 60.8 and xgboost achieved 61.9.

## A Appendix: Gradient and Hessian of $\mathcal{L}$

Proof. (Lemma 1) In a training context, the dataset $X$ and $y$ are constant, and since we are doing a gradient descent to train the weights of Tree ${ }^{(T+1)}$, so all previous trees (structure and weights) Tree ${ }^{(1)}, \ldots$, Tree $^{T}$ are fixed, as well as the structure of the current tree. The only free variables that remain are the tree weights: $\left(w_{1}, \ldots, w_{L}\right)$ associated to the leaves of Tree ${ }^{(T+1)}$.

For a sample $i \in[1, N]$ and a leaf $j \in[1, L]$, let $\delta_{i \in n_{j}}$ be the Kronecker symbol of the partition induced by the structure of Tree ${ }^{(T+1)}$ :

$$
\delta_{i \in n_{j}}=\left\{\begin{array}{l}
1 \text { iff. } \operatorname{Tree}\left(x_{i}\right) \text { ends in leaf } n_{j} \\
0 \text { otherwise }
\end{array}\right.
$$

For all sample $i \in[1, N]$, the evaluation function rewrites as:

$$
\operatorname{eval}_{x_{i}}\left(\operatorname{Tree}{ }^{(T+1)}\right)=\sum_{j=1}^{2^{d}} w_{j} \delta_{i \in n_{j}}
$$

in particular, this implies:

$$
\frac{\partial \mathrm{eval}}{\partial w_{j}}\left(w_{1}, \ldots, w_{2^{d}}\right)=\delta_{i \in n_{j}} \text { is constant. }
$$

Applying it to the loss function $\mathcal{L}$ of $\mathrm{Eq}(2)$, and since $y_{i}, \widehat{y}_{i}{ }^{(T)}$ are all constent, we deduce:

$$
\left.\left.\begin{array}{rl}
\frac{\partial \mathcal{L}}{\partial w_{j}} & =\sum_{i=1}^{N} \partial_{b} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}+\operatorname{eval}_{x_{i}}(\operatorname{Tree}\right. \\
\\
& \left.=\sum_{i=1}^{N} \partial_{b} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}\right)\right) \cdot \frac{\partial \operatorname{eval}_{x_{i}}}{\partial w_{j}}(\operatorname{Tree} \\
x_{i} & (T+1)
\end{array} \operatorname{Tree}^{(T+1)}\right)\right) \cdot \frac{\partial \operatorname{Reg}}{\partial w_{j}} . \delta_{i \in n_{j}}+\lambda w_{j} .
$$

And thus, the second derivative across $w_{i}$ and $w_{k}$, we get:

$$
\frac{\partial^{2} \mathcal{L}}{\partial w_{j} \partial w_{k}}=\sum_{i=1}^{N} \partial_{b}^{2} \operatorname{loss}\left(y_{i}, \widehat{y}_{i}^{(T)}+\operatorname{eval}_{x_{i}}\left(\operatorname{Tree}{ }^{(T+1)}\right)\right) \cdot \delta_{i \in n_{j}} \delta_{i \in n_{k}}+\lambda \delta_{j, k}
$$

All second derivatives across 2 different variables are zero, so the hessian of $\mathcal{L}$ is a pure Diagonal. Applied to the zero weights (i.e. eval $x_{x_{i}}\left(\operatorname{Tree}^{(T+1)}=0\right)$, the gradient and hessian are:

$$
\begin{aligned}
\frac{\partial \mathcal{L}}{\partial w_{j}}(0, \ldots, 0) & =\sum_{i=1}^{N} g_{i} \delta_{i \in n_{j}}=G \\
\frac{\partial^{2} \mathcal{L}}{\partial w_{j}^{2}}(0, \ldots, 0) & =\sum_{i=1}^{N} h_{i} \delta_{i \in n_{j}}+\lambda=H+\lambda
\end{aligned}
$$

which concludes the proof of Lemma 1.

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[^0]:    ${ }^{1}$ The implementation of [7] has been extended to allow for horizontally split datasets as well. However, we have not found any accompanying papers Homo SecureBoost [25].

