On Hilbert-Poincaré series of affine semi-regular polynomial sequences and related Gröbner bases

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Abstract Gröbner bases are nowadays central tools for solving various problems in commutative algebra and algebraic geometry. A typical use of Gröbner bases is the multivariate polynomial system solving, which enables us to construct algebraic attacks against post-quantum cryptographic protocols. Therefore, the determination of the complexity of computing Gröbner bases is very important both in theory and in practice: One of the most important cases is the case where input polynomials compose an (overdetermined) affine semi-regular sequence. The first part of this paper aims to present a survey on the Gröbner basis computation and its complexity. In the second part, we shall give an explicit formula on the (truncated) Hilbert-Poincaré series associated to the homogenization of an affine semi-regular sequence. Based on the formula, we also study (reduced) Gröbner bases of the ideals generated by an affine semi-regular sequence and its homogenization. Some of our results are considered to give mathematically rigorous proofs of the correctness of methods for computing Gröbner bases of the ideal generated by an affine semi-regular sequence.

1 Introduction

Let K be a field, and $R = K[x_1, \ldots, x_n]$ the polynomial ring in n variables over K. For a polynomial f in R, let f^{top} denote its maximal total-degree part which is called the *top part* here, and let f^h denote its homogenization in $R' = R[y]$ by an extra variable y, see Subsection 3.1.1 below for details. We denote by $\langle F \rangle_R$ (or $\langle F \rangle$ simply) the ideal generated by a non-empty subset F of R. For a finitely generated

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graded R-(or R'-)module M, we also denote by HF_M and HS_M its Hilbert function and its Hilbert–Poincaré series, respectively. A *Gröbner basis* of an ideal I in R is defined as a special kind of generator set for I, and it gives a computational tool to determine many properties of the ideal. A typical application of computing Gröbner bases is solving the multivariate polynomial (MP) problem: Given m polynomials $f_1, \ldots, f_m$ in R, find $(a_1, \ldots, a_n) \in K^n$ such that $f_i(a_1, \ldots, a_n) = 0$ for all i with $1 \le i \le m$. A particular case where polynomials are all quadratic is called the MQ problem, and its hardness is applied to constructing public-key cryptosystems and digital signature schemes that are expected to be quantum resistant. Therefore, analyzing the complexity of computing Gröbner bases is one of the most important problems both in theory and in practice.

An algorithm for computing Gröbner bases was proposed first by Buchberger [6], and so far a number of its improvements such as the F_4 [18] and F_5 [19] algorithms have been proposed, see Subsection 3.1 below for a summary. In general, it is very difficult to determine the complexity of computing Gröbner bases, but in some cases, we can estimate it with several algebraic invariants such as the solving degree, the degree of regularity, the Castelnuovo–Mumford regularity, and the first and last fall degrees; we refer to [8] for the relations between these invariants.

The first part of this paper aims to survey the Gröbner basis computation, and to review its complexity in the case where input polynomials generate a zero-dimensional ideal. For this, in Section 2, we first recall foundations in commutative algebra such as Koszul complex, Hilbert-Poincaré series, and semi-regular sequence, which are useful ingredients to estimate the complexity of computing Gröbner bases. Then, we overview existing Gröbner basis algorithms in Subsection 3.1. Subsequently, it will be described in Subsection 3.2 how to estimate the complexity of computing the reduced Gröbner basis of a zero-dimensional ideal, with the notion of homogenization.

In the second part, we focus on *affine semi-regular* polynomial sequences, where a sequence $F = (f_1, \ldots, f_m) \in R^m$ of (not necessarily homogeneous) polynomials is said to be affine (cryptographic) semi-regular if $F^{\mathrm{top}} = (f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}})$ is (cryptographic) semi-regular, see Definitions 4, 7, and 8 for details. Note that homogeneous semi-regular sequences are conjectured by Pardue [30, Conjecture B] to be generic sequences of polynomials, and affine (cryptographic) semi-regular sequences are often appearing in the construction of multivariate public key cryptosystems and digital signature schemes. In Section 4 below, we relate the Hilbert-Poincaré series of $R'/\langle F^h\rangle$ with that of $R/\langle F^{\mathrm{top}}\rangle$. As a corollary, we obtain an explicit formula of the truncation at degree $D-1$ of the Hilbert-Poincaré series of $R'/\langle F^h\rangle$, where D is the degree of regularity for $\langle F^{\mathrm{top}}\rangle$. The following theorem summarizes these results:

Theorem 1 (Theorem 7, Corollaries 1 and 2) *With notation as above, assume that F is affine cryptographic semi-regular. Then $\mathrm{HF}_{R'/\langle F^h\rangle}(d) = \sum_{i=0}^{d} \mathrm{HF}_{R/\langle F^{\mathrm{top}}\rangle}(i)$ and $(\langle \mathrm{LM}(\langle F^h\rangle)\rangle_{R'})_d = (\langle \mathrm{LM}(\langle F^{\mathrm{top}}\rangle)\rangle_{R'})_d$ for each d with $d < D$. Hence, we also obtain $\mathrm{HS}_{R'/\langle F^h\rangle}(z) \equiv \prod_{i=1}^{m}(1-z^{d_i})/(1-z)^{n+1} \pmod{z^D}$, so that F^h is D-regular (see Definition 4 for the definition of the d-regularity).*

As an application of this theorem, we explore reduced Gröbner bases of $\langle F\rangle$, $\langle F^h\rangle$, and $\langle F^{\mathrm{top}}\rangle$ in Section 5 below, dividing the cases into the degree less than D or not. In particular, we rigorously prove some existing results, which are often used for analyzing the complexity of computing Gröbner bases, and moreover extend them.

2 Preliminaries

In this section, we recall definitions of Koszul complex, Hilbert–Poincaré series, and semi-regular polynomial sequences, and collect some known facts related to them. Throughout this section, let $R = K[X] = K[x_1, \ldots, x_n]$ be the polynomial ring of n variables $X = (x_1, \ldots, x_n)$ over a field K. As a notion, for a polynomial f in R, we denote its total degree by $\deg(f)$. As R is a graded ring with respect to total degree, for a polynomial f, its maximal total degree part, denoted by f^{top}, is defined as its graded component of $\deg(f)$, that is, the sum of all terms of f whose total degree equals to $\deg(f)$.

2.1 Koszul complex and its homology

Let $f_1, \ldots, f_m \in R$ be homogeneous polynomials of degrees $d_1, \ldots, d_m$, and put $d_{j_1\cdots j_i} := \sum_{k=1}^{i} d_{j_k}$. For each $0 \le i \le m$, we define a free R-module of rank $\binom{m}{i}$

$$K_i(f_1, \ldots, f_m) := \begin{cases} \displaystyle\bigoplus_{1\le j_1<\cdots<j_i\le m} R(-d_{j_1\cdots j_i})\mathbf{e}_{j_1\cdots j_i} & (i \ge 1) \\ R & (i = 0), \end{cases}$$

where $\mathbf{e}_{j_1\cdots j_i}$ is a standard basis. We also define a graded homomorphism

$$\varphi_i : K_i(f_1, \ldots, f_m) \longrightarrow K_{i-1}(f_1, \ldots, f_m) \tag{1}$$

of degree 0 by

$$\varphi_i(\mathbf{e}_{j_1\cdots j_i}) := \sum_{k=1}^{i} (-1)^{k-1} f_{j_k} \mathbf{e}_{j_1\cdots \hat{j}_k \cdots j_i}. \tag{2}$$

Here, $\hat{j}_k$ means to omit j_k. For example, we have $\mathbf{e}_{1\hat{2}3} = \mathbf{e}_{13}$. To simplify the notation, we set $K_i := K_i(f_1, \ldots, f_m)$. Then,

$$K_\bullet : 0 \to K_m \xrightarrow{\varphi_m} \cdots \xrightarrow{\varphi_3} K_2 \xrightarrow{\varphi_2} K_1 \xrightarrow{\varphi_1} K_0 \to 0 \tag{3}$$

is a complex, and we call it the *Koszul complex* of $(f_1, \ldots, f_m)$. The i-th homology group is given by

$$H_i(K_\bullet) = \mathrm{Ker}(\varphi_i)/\mathrm{Im}(\varphi_{i+1}). \tag{4}$$

In particular, we have
$$H_0(K_\bullet) = R/\langle f_1, \dots, f_m \rangle_R. \tag{5}$$
The kernel and the image of a graded homomorphism are both graded submodule in general, so that $\mathrm{Ker}(\varphi_i)$ and $\mathrm{Im}(\varphi_{i+1})$ are graded R-module, and so is the quotient module $H_i(K_\bullet)$. In the following, we denote by $H_i(K_\bullet)_d$ the degree-d homogeneous part of $H_i(K_\bullet)$.

Note that $\mathrm{Ker}(\varphi_1) = \mathrm{syz}(f_1, \dots, f_m)$ (the right hand side is the module of syzygies), and that $\mathrm{Im}(\varphi_2) \subset K_1 = \bigoplus_{j=1}^m R(-d_j)\mathbf{e}_j$ is generated by

$$\{\mathbf{t}_{i,j} := f_i \mathbf{e}_j - f_j \mathbf{e}_i : 1 \leq i < j \leq m\}. \tag{6}$$

Hence, putting

$$\mathrm{tsyz}(f_1, \dots, f_m) := \langle \mathbf{t}_{i,j} : 1 \leq i < j \leq m \rangle_R, \tag{7}$$

we have
$$H_1(K_\bullet) = \mathrm{syz}(f_1, \dots, f_m)/\mathrm{tsyz}(f_1, \dots, f_m). \tag{8}$$

Definition 1 (Trivial syzygies) With notation as above, we call each generator $\mathbf{t}_{i,j}$ (or each element of $\mathrm{tsyz}(f_1, \dots, f_m)$) a *trivial syzygy* for $(f_1, \dots, f_m)$. We also call $\mathrm{tsyz}(f_1, \dots, f_m)$ the *module of trivial syzygies*.

We also note that $H_m(K_\bullet) = 0$, since φ_m is clearly injective by definition.

Remark 1 When $K = \mathbb{F}_q$, a vector of the form $f_i^{q-1}\mathbf{e}_i$ is also referred to as a trivial syzygy, in the context of Ding-Schmidt's definition for *first fall degree* [15] (see [7, Section 4.2] or [28, Section 3.2] for reviews). More concretely, putting $B := R/\langle x_1^q, \dots, x_n^q \rangle_R$ and $\overline{f}_i := f_i \bmod \langle x_1^q, \dots, x_n^q \rangle$, we define the Koszul complex of $(\overline{f}_1, \dots, \overline{f}_m) \in B^m$ similarly to that of $(f_1, \dots, f_m) \in R^m$, and denote it by $\overline{K}_\bullet = \overline{K}_\bullet(\overline{f}_1, \dots, \overline{f}_m)$. Then, the vectors $\overline{f}_i\mathbf{e}_j - \overline{f}_j\mathbf{e}_i$ and $\overline{f}_i^{q-1}\mathbf{e}_i$ in B^m for $1 \leq i < j \leq m$ are syzygies for $(\overline{f}_1, \dots, \overline{f}_m)$. Each $\overline{f}_i\mathbf{e}_j - \overline{f}_j\mathbf{e}_i$ is called a Koszul syzygy, and the Koszul syzygies together with $\overline{f}_i^{q-1}\mathbf{e}_i$'s are referred to as trivial syzygies for $(\overline{f}_1, \dots, \overline{f}_m)$. The *first fall degree* $d_{\mathrm{ff}}(f_1, \dots, f_m)$ is equal to the minimal integer d with $\mathrm{syz}(\overline{f}_1, \dots, \overline{f}_m)_d \supsetneq \mathrm{tsyz}^+(\overline{f}_1, \dots, \overline{f}_m)_d$ in $(B_{d-d_i})^m$, where $\mathrm{tsyz}^+(\overline{f}_1, \dots, \overline{f}_m)$ denotes the submodule in B^m generated by the trivial syzygies for $(\overline{f}_1, \dots, \overline{f}_m)$.

Note that, for each i, a homomorphism $H_i(K_\bullet) \to H_i(\overline{K}_\bullet)$ is canonically induced by taking modulo $\langle x_1^q, \dots, x_n^q \rangle_R$. In particular, we have the following composite K-linear map:

$$\eta_d : H_1(K_\bullet)_d \to H_1(\overline{K}_\bullet)_d \to \mathrm{syz}(\overline{f}_1, \dots, \overline{f}_m)_d/\mathrm{tsyz}^+(\overline{f}_1, \dots, \overline{f}_m)_d.$$

for each d. Putting $d = d_{\mathrm{ff}}(f_1, \dots, f_m)$ and letting D to be the minimal integer with $H_1(K_\bullet)_D \neq 0$, it is straightforward to verify the following:

- If $q > D$, then η_D is injective, and $\mathrm{syz}(\overline{f}_1, \ldots, \overline{f}_m)_D \supsetneq \mathrm{tsyz}^{+}(\overline{f}_1, \ldots, \overline{f}_m)_D$, whence $D \geq d$.
- If $q > d$, then η_d is surjective. In this case, $H_1(K)_d \neq 0$, so that $D \leq d$.

See [28, Lemmas 4.2 and 4.3] for a proof. Therefore, we have $d = D$ for sufficiently large any q.

2.2 Hilbert–Poincaré series and semi-regular sequences

Definition 2 (Hilbert–Poincaré series) For a finitely generated graded R-module M, we define the *Hilbert function* HF_M of M, given by

$$\mathrm{HF}_M(d) = \dim_K M_d \tag{9}$$

for each $d \in \mathbb{Z}_{\geq 0}$. The *Hilbert–Poincaré series* HS_M of M is defined as the formal power series

$$\mathrm{HS}_M(z) = \sum_{d=0}^{\infty} \mathrm{HF}_M(d) z^d \in \mathbb{Z}[\![z]\!]. \tag{10}$$

Theorem 2 (cf. [4, Chapter 10]) *Let I be a homogeneous ideal of R generated by a set $G \subset R$ of homogeneous elements of degree not greater than d. Let $\prec$ be a graded ordering on the monomials in R. Then, the following are equivalent:*

1. $\langle \mathrm{LT}(G) \rangle_{\leq d} = \langle \mathrm{LT}(I) \rangle_{\leq d}$.
2. *Every $f \in I_{\leq d}$ is reduced to zero modulo G.*
3. *For every pair of $f, g \in G$ with $\deg(\mathrm{LCM}(\mathrm{LM}(f), \mathrm{LM}(g))) \leq d$, the S-polynomial $S(f, g)$ is reduced to zero modulo G.*

In this case, G is called a d-Gröbner basis of I with respect to $\prec$.

We also review the notion of semi-regular sequence defined by Pardue [30].

Definition 3 (Semi-regular sequences, [30, Definition 1]) Let I be a homogeneous ideal of R. A degree-d homogeneous element $f \in R$ is said to be *semi-regular* on I if the multiplication map $(R/I)_{t-d} \longrightarrow (R/I)_d$; $g \mapsto gf$ is injective or surjective, for every t with $t \geq d$. A sequence $(f_1, \ldots, f_m) \in R^m$ of homogeneous polynomials is said to be *semi-regular* on I if f_i is semi-regular on $I + \langle f_1, \ldots, f_{i-1} \rangle_R$, for every i with $1 \leq i \leq m$.

Throughout the rest of this subsection, let $f_1, \ldots, f_m \in R$ be homogeneous elements of degree $d_1, \ldots, d_m$ respectively, and put $I = \langle f_1, \ldots, f_m \rangle_R$. Furthermore, put $I^{(0)} := \{0\}$, $A^{(0)} := R/I^{(0)} = R$. For each i with $1 \leq i \leq m$, we also set $I^{(i)} := \langle f_1, \ldots, f_i \rangle_R$, $A^{(i)} := R/I^{(i)}$. The degree-d homogeneous part $A^{(i)}_d$ of each $A^{(i)}$ is given by $A^{(i)}_d = R_d / I^{(i)}_d$, where $I^{(i)}_d = I^{(i)} \cap R_d$. We denote by ψ_{f_i} the multiplication map

$$A^{(i-1)} \longrightarrow A^{(i-1)} \ ; \ g \mapsto g f_i, \tag{11}$$

which is a graded homomorphism of degree d_i. For every $t \geq d_i$, the restriction map

$$\psi_{f_i}|_{A^{(i-1)}_{t-d_i}} : A^{(i-1)}_{t-d_i} \longrightarrow A^{(i-1)}_t$$

is a K-linear map. On the other hand, as for the surjective homomorphism

$$\phi_{i-1} : A^{(i-1)} \longrightarrow A^{(i)} \ ; \ f + I^{(i-1)} \mapsto f + I^{(i)}, \tag{12}$$

it is straightforward to see that for each t with $0 \leq t \leq d_i - 1$, the restriction map

$$\phi_{i-1}|_{A^{(i-1)}_t} : A^{(i-1)}_t \longrightarrow A^{(i)}_t$$

is an isomorphism of K-linear spaces, whence

$$\dim_K A^{(i-1)}_t = \dim_K A^{(i)}_t \quad (0 \leq t \leq d_i - 1). \tag{13}$$

Lemma 1 *With notation as above, for each $1 \leq i \leq m$ and for each $t \geq d_i$, we have the following equalities:*

$$\dim_K(A^{(i)}_t) = \dim_K(A^{(i-1)}_t) - \dim_K \mathrm{Im}\left(A^{(i-1)}_{t-d_i} \xrightarrow{\times f_i} A^{(i-1)}_t\right), \tag{14}$$

$$\dim_K \mathrm{Im}\left(A^{(i-1)}_{t-d_i} \xrightarrow{\times f_i} A^{(i-1)}_t\right) = \dim_K(A^{(i-1)}_{t-d_i}) - \dim_K(0 : f_i)_{t-d_i}, \tag{15}$$

where we set $(0 : f_i) = \{g \in A^{(i-1)} : g f_i = 0\}$. Hence,

- *The multiplication map $A^{(i-1)}_{t-d_i} \xrightarrow{\times f_i} A^{(i-1)}_t$ is injective if and only if*

$$\dim_K(A^{(i)}_t) = \dim_K(A^{(i-1)}_t) - \dim_K(A^{(i-1)}_{t-d_i}). \tag{16}$$

  In this case, one has $\dim_K(A^{(i-1)}_{t-d_i}) \leq \dim_K(A^{(i-1)}_t)$.
- *The multiplication map $A^{(i-1)}_{t-d_i} \xrightarrow{\times f_i} A^{(i-1)}_t$ is surjective if and only if*

$$\dim_K(A^{(i)}_t) = 0. \tag{17}$$

  In this case, one has $\dim_K(A^{(i-1)}_{t-d_i}) \geq \dim_K(A^{(i-1)}_t)$.

Proof. Let i and t be integers such that $1 \leq i \leq m$ and $t \geq d_i$. Since we have $(0 : f_i)_{t-d_i} = \{g \in A^{(i-1)}_{t-d_i} : g f_i = 0\}$, the sequence

$$0 \longrightarrow (0 : f_i)_{t-d_i} \longrightarrow A^{(i-1)}_{t-d_i} \xrightarrow{\times f_i} A^{(i-1)}_t \longrightarrow A^{(i)}_t \longrightarrow 0 \tag{18}$$

of K-linear maps is exact, where $(0 : f_i)_{t-d_i} \rightarrow A^{(i-1)}_{t-d_i}$ is an inclusion map. The exactness of this sequence implies the desired equalities (14) and (15). □

The semi-regularity is characterized by equivalent conditions in Proposition 1 below. In particular, the fourth condition enables us to compute the Hilbert–Poincaré series of each $A^{(i)}$ easily.

Proposition 1 (cf. [30, Proposition 1]) *With notation as above, the following are equivalent:*

1. *$(f_1, \ldots, f_m) \in R^m$ is semi-regular.*
2. *For each $1 \le i \le m$ and for each $t \ge d_i$, we have (16) or (17), namely*

$$\dim_K(A_t^{(i)}) = \max\{0, \dim_K(A_t^{(i-1)}) - \dim_K(A_{t-d_i}^{(i-1)})\}.$$

3. *For each i with $1 \le i \le m$, we have*

$$\mathrm{HS}_{A^{(i)}}(z) = [\mathrm{HS}_{A^{(i-1)}}(z)(1 - z^{d_i})]. \tag{19}$$

4. *For each i with $1 \le i \le m$, we have*

$$\mathrm{HS}_{A^{(i)}}(z) = \left[\frac{\prod_{j=1}^{i}(1 - z^{d_j})}{(1 - z)^n} \right]. \tag{20}$$

When K is an infinite field, Pardue also conjectured in [30, Conjecture B] that generic polynomial sequences are semi-regular.

2.3 Cryptographic semi-regular sequences

We here review the notion of *cryptographic semi-regular* sequences, which are defined by a condition weaker than one for semi-regular sequences. The notion of cryptographic semi-regular sequences is introduced first by Bardet et al. (e.g., [2], [3]) motivated to analyzing the complexity of computing Gröbner bases. Diem [13] also formulated cryptographic semi-regular sequences, in terms of commutative and homological algebra. The terminology "cryptographic" was named by Bigdeli et al. in their recent work [5], in order to distinguish such a sequence from a semi-regular sequence defined by Pardue (see Definition 3 in the previous subsection).

Definition 4 ([2, Definition 3]; see also [13, Definition 1]) Let $f_1, \ldots, f_m \in R$ be homogeneous polynomials of positive degrees $d_1, \ldots, d_m$ respectively, and put $I = \langle f_1, \ldots, f_m \rangle_R$. The notations $I^{(i)}$ and $A^{(i)}$ are also the same as in the previous subsection. For each integer d with $d \ge \max\{d_i : 1 \le i \le m\}$, we call a sequence $(f_1, \ldots, f_m) \in R^m$ of homogeneous polynomials *d-regular* if it satisfies the following condition:

- For each i with $1 \le i \le m$, if a homogeneous polynomial $g \in R$ satisfies $g f_i \in \langle f_1, \ldots, f_{i-1} \rangle_R$ and $\deg(g f_i) < d$, then we have $g \in \langle f_1, \ldots, f_{i-1} \rangle_R$. In other word, the multiplication map $A_{t-d_i}^{(i-1)} \longrightarrow A_t^{(i-1)}$; $g \mapsto g f_i$ is injective for every t with $d_i \le t < d$.

Diem [13] determined the (truncated) Hilbert series of d-regular sequences as in the following proposition:

Theorem 3 (cf. [13, Theorem 1]) *With the same notation as in Definition 4, the following are equivalent for each d with $d \geq \max\{d_i : 1 \leq i \leq m\}$:*

1. *$(f_1, \ldots, f_m) \in R^m$ is d-regular. Namely, for each (i, t) with $1 \leq i \leq m$ and $d_i \leq t < d$, the equality (16) holds.*
2. *We have*
$$\mathrm{HS}_{A^{(m)}}(z) \equiv \frac{\prod_{j=1}^{m}(1 - z^{d_j})}{(1 - z)^n} \pmod{z^d}. \tag{21}$$
3. *$H_1(K_\bullet(f_1, \ldots, f_m))_{\leq d-1} = 0$.*

Proposition 2 ([13, Proposition 2 (a)]) *With the same notation as in Definition 4, let D and i be natural numbers. Assume that $H_i(K(f_1, \ldots, f_m))_{\leq D} = 0$. Then, for each j with $1 \leq j < m$, we have $H_i(K(f_1, \ldots, f_j))_{\leq D} = 0$.*

Definition 5 A finitely generated graded R-module M is said to be *Artinian* if there exists a sufficiently large $D \in \mathbb{Z}$ such that $M_d = 0$ for all $d \geq D$.

Definition 6 ([2, Definition 4], [3, Definition 5]) For a homogeneous ideal I of R, we define its *degree of regularity* $d_{\mathrm{reg}}(I)$ as follows: If the finitely generated graded R-module R/I is Artinian, we set $d_{\mathrm{reg}}(I) := \min\{d : R_d = I_d\}$, and otherwise we set $d_{\mathrm{reg}}(I) := \infty$.

As for an upper-bound on the degree of regularity, we refer to [22, Theorem 21].

Remark 2 In Definition 6, since R/I is Noetherian, it is Artinian if and only if it is of finite length. In this case, the degree of regularity $d_{\mathrm{reg}}(I)$ is equal to the *Castelnuovo-Mumford regularity* $\mathrm{reg}(I)$ of I (see e.g., [16, §20.5] for the definition), whence $d_{\mathrm{reg}}(I) = \mathrm{reg}(I) = \mathrm{reg}(R/I) + 1$.

Definition 7 ([2, Definition 5], [3, Definition 5]; see also [14, Section 2]) A sequence $(f_1, \ldots, f_m) \in R^m$ of homogeneous polynomials is said to be *cryptographic semi-regular* if it is $d_{\mathrm{reg}}(I)$-regular, where we set $I = \langle f_1, \ldots, f_m \rangle_R$.

The cryptographic semi-regularity is characterized by equivalent conditions in Proposition 3 below. In particular, the second condition enables us to compute the Hilbert–Poincaré series of $A^{(i)}$ easily.

Proposition 3 ([13, Proposition 1 (d)]; see also [3, Proposition 6]) *With the same notation as in Definition 4, we put $D = d_{\mathrm{reg}}(I)$. Then, the following are equivalent:*

1. *$(f_1, \ldots, f_m) \in R^m$ is cryptographic semi-regular.*
2. *We have*
$$\mathrm{HS}_{R/I}(z) = \left[\frac{\prod_{j=1}^{m}(1 - z^{d_j})}{(1 - z)^n} \right]. \tag{22}$$
3. *$H_1(K_\bullet(f_1, \ldots, f_m))_{\leq D-1} = 0$.*

Remark 3 By the definition of *degree of regularity*, if $(f_1, \ldots, f_m) \in R^m$ is cryptographic semi-regular, $d_{\mathrm{reg}}(I)$ coincides with $\deg(\mathrm{HS}_{R/I}(z)) + 1$, where we set $I = \langle f_1, \ldots, f_m \rangle$.

In 1985, Fröberg already conjectured in [21] that, when K is an infinite field, a sequence of homogeneous polynomials $f_1, \ldots, f_m \in R$ of degrees $d_1, \ldots, d_m$ generates an ideal I with the Hilbert series of the form (22), namely $(f_1, \ldots, f_m)$ is cryptographic semi-regular. It can be proved (cf. [30]) that Fröberg's conjecture is equivalent to Pardue's one [30, Conjecture B] introduced in Subsection 2.2. We also note that Moreno-Socías conjecture [27] is stronger than the above two conjectures, see [30, Theorem 2] for a proof.

It follows from the fourth condition of Proposition 1 together with the second condition of Proposition 3 that the semi-regularity implies the cryptographic semi-regularity.

Definition 8 (Affine semi-regular sequences) A sequence $F = (f_1, \ldots, f_m) \in R^m$ of not necessarily homogeneous polynomials $f_1, \ldots, f_m$ is said to be semi-regular (resp. cryptographic semi-regular) if $F^{\mathrm{top}} = (f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}})$ is semi-regular (resp. cryptographic semi-regular). In this case, we call F an *affine semi-regular (resp. affine cryptographic semi-regular)* sequence.

Remark 4 For an affine crytographic semi-regular sequence $F = (f_1, \ldots, f_m) \in R^m$ with $K = \mathbb{F}_q$, it follows from Proposition 3 that $d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle) \leq d_{\mathrm{ff}}(f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}})$ for $q \gg 0$, where $d_{\mathrm{ff}}(f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}})$ is the first fall degree defined in Remark 1.

3 Quick review on the computation of Gröbner basis

In this section, we first review previous studies on the computation of Gröbner bases for polynomial ideals.

3.1 Overview of existing Gröbner basis algorithms

Since Buchberger [6] discovered the notion of Gröbner bases and a fundamental algorithm for computing them, many efforts have been done for improving the efficiency of the computation based on Buchberger's algorithm. In his algorithm, S-polynomials play an important role for the Gröbner basis construction and give a famous termination criterion called Buchberger's criterion, that is, for a given ideal I of a polynomial ring over a field, its finite generating subset G is a Gröbner basis of I with respect to a monomial ordering if and only if the S-polynomial $S(g, g')$ for any distinct pair $g, g' \in G$ is reduced to 0 modulo G. For details on Buchberger's algorithm and monomial orders, see e.g., [4].

In the below, we list effective improvements for algorithms which are, at the same time, very useful to analyze the complexity of the Gröbner basis computation. Here we note that the choice of a monomial order is also very crucial for the efficiency of the Gröber basis computation, but we here do not discuss about its choice. (In general, the degree reverse lexicographical (DRL) ordering[1] is considered as the most efficient ordering for the computation.)

(1) **Related to S-polynomial:**

(1-1) **Strategy for selecting S-polynomial:** It is considered to be very effective to apply so called the normal strategy, where we choose a pair (g, g') for which the least common multiple (LCM) of the leading monomials $\mathrm{LM}(g)$ and $\mathrm{LM}(g')$ with respect to the fixed ordering $\prec$ as smaller as possible. (See [4, Chapter 5.5].) The strategy is very suited for a homogeneous ideal with a *degree compatible* ordering such as DRL, as we can utilize the graded structure of a homogeneous ideal. Also, the *sugar strategy* is designed for a non-homogeneous ideal to make the computational behavior very close to that for the homogenized ideal. See Subsection 3.1.1 below for some details on homogenization. (See also [11, Chapter 2.10].)
(1-2) **Avoiding unnecessary S-polynomial computation:** In Buchberger's algorithm, we add a polynomial to a generating set G which is computed from an S-polynomial by possible reduction of elements in G. Since the cost of the construction of S-polynomials and their reduction dominate the whole computation, S-polynomials which are reduced to 0 are very harmful for the efficiency. Thus, it is highly desired to avoid such unnecessary S-polynomials as many as possible.
(A) **Based on simple rules:** At earlier stage, there are easily computable rules, called Buchberger's criterion and Gebauer-Möller's one. Those are using the relation of the LMs of a pair and those of a triple, see [4, Chapter 5.5]. Then, in 2002, Faugére [19] introduced the notion of *signature* and proposed his F_5 algorithm based on a general rule among signatures. We call algorithms using signatures including variants of F_5 *signature-based algorithms (SBA)*. See a survey [17] and the detail in Subsection 3.1.2 below.
(B) **Using invariants of polynomial ideal:** For a homogeneous ideal I of a polynomial ring R, when its Hilbert function $\mathrm{HF}_{R/I}(z)$ is known before the computation, we can utilize the value $\mathrm{HF}_{R/I}(d)$ for each $d \in \mathbb{N}$ (cf. [35]). Because, by the value $\mathrm{HF}_{R/I}(d)$, we can check whether we can stop the computation of S-polynomials of degree d or not. We call an algorithm using Hilbert functions a *Hilbert driven* (Buchberger's) algorithm. See [11, Chapter 10.2] or [12, Section 3.5].

(2) **Efficient computation of S-polynomial reduction:** Since the computation of S-polynomial reduction is a dominant step in the whole Gröbner basis computation, its efficiency heavily affects the total efficiency. As the reduction for a polynomial by elements of G is sequentially applied, we can transform the whole reduction to a Gaussian elimination of a matrix. This approach was suggested

[1] This ordering is also called the graded reverse lexicographical (grevlex) ordering.

in form of Macaulay matrices by Lazard [25] and the first efficient algorithm was given by Faugére [18], which is called the F_4 algorithm. Of course, we can combine the F_4 and F_5 algorithms effectively, which is called the *matrix-F_5 algorithm*.
(3) **Solving coefficient growth:** For a polynomial ideal over the rational number field $\mathbb{Q}$, the computation may be suffered by certain growth of coefficients in polynomials appearing during the Gröbner basis computation. To resolve this problem, several modular methods were proposed. As a typical one, we can use Chinese remainder algorithm (CRA), where we first compute the reduced Gröbner bases G_p over several finite fields $\mathbb{F}_p$ and then recover the reduced Gröbner basis from G_p's by CRA. See [29] for details about choosing primes p.

Remark 5 For several public key cryptosystems based on polynomial ideals over finite fields or the elliptic curve discrete logarithm problem, estimating the cost of finding zeros of polynomial ideals is important to analyze the security of those systems, where the computation of their Gröbner bases is a fundamental tool. In this situation, the F_5 algorithm and matrix-F_5 algorithm as its efficient variant with an efficient DRL ordering are considered, as not only those can attain efficient computation but also they are suited for estimating the computational complexity.

In the below, we introduce the notion of *homogenization* and an algorithm for Gröbner basis computation based on signature (F_5 or its variants) which will be used for our study in Section 5 below.

3.1.1 Homogenization of polynomials and monomial orders

We begin with recalling the notion of homogenization. (See [23, Chapter 4] for details.) Let K be a field, $X = \{x_1, \dots, x_n\}$ a set of variables, and $\mathcal{T}$ the set of all monomials in X. [2]

(1) For a non-homogeneous polynomial $f = \sum_{t \in \mathcal{T}} c_t t$ in $K[X]$ with $c_t \in K$, its *homogenization* f^h is defined, by introducing a new variable y, as

$$f^h = \sum_{t \in \mathcal{T}} c_t t y^{\deg(f) - \deg(t)}.$$

Thus f^h is a homogeneous polynomial in $X \cup \{y\}$ over K with total degree $d = \deg(f)$. Also for a set F of polynomials, its *homogenization* F^h is defined as $F^h = \{f^h \mid f \in F\}$. We also write X^h for $X \cup \{y\}$.
(2) Conversely, for a homogeneous polynomial h in $K[X \cup \{y\}]$, its *dehomogenization* h^{deh} is defined by substituting y with 1, that is, $h^{\mathrm{deh}} = h(X, 1)$. (It is also denoted by $h|_{y=1}$.) For a set H of homogeneous polynomials in $K[X \cup \{y\}]$, its *dehomogenization* H^{deh} (or $H|_{y=1}$) is defined as $H^{\mathrm{deh}} = \{h^{\mathrm{deh}} \mid h \in H\}$.

[2] As the symbolc m is used for the size of a generating set, we use $\mathcal{T}$ instead of $\mathcal{M}$.

(3) For an ideal I of $K[X]$, its homogenization I^h, as an ideal, is defined as $\langle I^h \rangle_{K[X\cup\{y\}]}$.
(4) For a homogeneous ideal J in $K[X \cup \{y\}]$, its dehomogenization J^{deh}, as a set, is an ideal of $K[X]$. We note that if a homogeneous ideal J is generated by H, then $J^{\mathrm{deh}} = \langle H^{\mathrm{deh}} \rangle_{K[X]}$ and for an ideal I of $K[X]$, we have $(I^h)^{\mathrm{deh}} = I$.
(5) For a monomial (term) order $\prec$ on the set of *monomials* $\mathcal{T}$ in X, its *homogenization* $\prec_h$ on the set of *monomials* $\mathcal{T}^h$ in $X \cup \{y\}$ is defined as follows: For two monomials $X^\alpha y^a$, $X^\beta y^b$ in $\mathcal{T}^h$, $X^\alpha y^a \prec_h X^\beta y^b$ if and only if one of the following holds:

 (i) $a + |\alpha| < b + |\beta|$, or
 (ii) $a + |\alpha| = b + |\beta|$ and $X^\alpha \prec X^\beta$,

 where $\alpha = (\alpha_1, \dots, \alpha_n) \in \mathbb{Z}^n_{\geq 0}$ and $|\alpha| = \alpha_1 + \cdots + \alpha_n$, and where X^α denotes $x_1^{\alpha_1} \cdots x_n^{\alpha_n}$. Here, for a monomial $X^\alpha y^a$, we call X^α and y^a *the* X*-part* and the y-part, respectively. If a monomial order $\prec$ is *degree-compatible*, that is, it first compares the total degrees, the restriction $\prec_h |_{\mathcal{T}}$ of $\prec_h$ on $\mathcal{T}$ coincides with $\prec$.

It is well-known that for a Gröbner basis H of $\langle F^h \rangle$ with respect to $\prec^h$, its dehomonization $\{h^{\mathrm{deh}} \mid h \in H\}$ is also a Gröbner basis of $\langle F \rangle$ with respect to $\prec$.

3.1.2 Signature and F_5 algorithm

Here we briefly outline the F_5 algorithm, which is an improvement of Buchberger's algorithm. For details, see a survey [17]. Let $F = \{f_1, \dots, f_m\} \subset R = K[X]$ be a given generating set. For each polynomial h constructed during Gröbner basis computation of $\langle F \rangle$, the F_5 algorithm attaches a *special label called a signature* as follows: Since h belongs to $\langle F \rangle$, it can be written as

$$h = a_1 f_1 + a_2 f_2 + \cdots + a_m f_m, \tag{23}$$

for some $a_1, \dots, a_m \in R$. Then, we assign h to $a_1 \mathbf{e}_1 + \cdots + a_m \mathbf{e}_m \in R^m$ and we call its leading monomial $t\mathbf{e}_i$ with respect to a monomial (module) ordering in R^m *the signature* of h. As the expression (23) is not unique, in order to determine the signature, we construct the expression procedurally or use the uniquely determined residue in $R^m/\mathrm{syz}(f_1, \dots, f_m)$ by a module Gröbner basis of $\mathrm{syz}(f_1, \dots, f_m)$. (For the latter case, we call it the *minimal* signature.) Here we denote the signature of h by $\mathrm{sig}(h)$. Anyway, in the F_5 algorithm, we can meet the both by carefully choosing S-polynomials and by applying restricted reduction steps (called Σ-reductions) for S-polynomials without any change of the signature. (So, we need not to compute a module Gröbner basis of $\mathrm{syz}(f_1, \dots, f_m)$.) We note that for the S-polynomial $S(h_1, h_2) = c_1 t_1 h_1 - c_2 t_2 h_2$ with $c_1, c_2 \in K$ and $t_1, t_2 \in \mathcal{T}$, the signature $\mathrm{sig}(S(h_1, h_2))$ is determined as the largest one between $\mathrm{sig}(c_1 t_1 h_1)$ and $\mathrm{sig}(c_2 t_2 h_2)$. Then, we have the following criteria which are very useful to avoid the computation of unnecessary S-polynomials. (The latter one is called the *syzygy criterion*.)

Proposition 4 (cf. [11], [17]) *In the F_5 algorithm, we need not compute an S-polynomial if some S-polynomial of the same signature was already proceeded,* *since both are reduced to the same polynomial. Moreover, we need not compute an S-polynomial of signature s if there is a signature s' such that s' divides s and some S-polynomial with the signature s' is reduced to* 0*.*

3.2 Complexity of the Gröbner basis computation

In general, determining the complexity of computing a Gröbner basis is very hard; in the worst-case, the complexity is doubly exponential in the number of variables, see e.g., [9], [26], [31] for surveys. It is well-known that a Gröbner basis with respect to a degree-compatible monomial ordering (in particular, DRL ordering) can be computed quite more efficiently than ones with respect to other orderings in general. Moreover, in the case where the input polynomials generate a zero-dimensional ideal, once a Gröbner basis with respect to an efficient monomial ordering is computed, one with respect to any other ordering can be computed easily by the FGLM basis conversion algorithm [20]. From this, we focus on the case where the monomial ordering is degree-compatible, and if necessary we also assume that the ideal generated by the input polynomials is zero-dimensional.

A typical way to estimate the complexity is to count the number of S-polynomials that are reduced during the Gröbner basis computation. In the case where the chosen monomial ordering is degree-compatible, the most efficient strategy to compute Gröbner bases is the so-called normal strategy, which proceed *degree by degree*, namely increase the degree of critical pairs defining S-polynomials, as in the F_4 and F_5 algorithms. For an algorithm adopting this strategy, several S-polynomials are dealt with consecutively at the same degree, which is called the *step degree*. The lowest step degree at which an intermediate ideal basis contains a minimal Gröbner basis is called the *solving degree* of the algorithm. Determining (or finding a tight bound for) the solving degree is difficult without computing any Gröbner basis. On the other hand, once it is specified, we can estimate the complexity of the algorithm in a straightforward manner. For example, for a linear algebra-based algorithm, such as an F_4-family including the (matrix-)F_5 algorithm and the XL family (cf. [10]), that follows Lazard's atrategy [24] to reduces S-polynomials by the Gaussian elimination on Macaulay matrices, its complexity is estimated to be $O(N^\omega)$ with $N = \binom{n+D}{n}$, where D is the solving degree, and where ω is the *matrix multiplication exponent* with $2 \le \omega < 3$.

Caminata et al. (cf. [5], [7], [8]) defined the solving degree of an algorithm based on Lazard's strategy in a mathematically rigorous way, see their papers for the definition. For a given polynomial sequence $F = (f_1, \ldots, f_m) \in R^m$ and a degree-compatible monomial ordering $\prec$, it is denoted by $\mathrm{sd}_\prec(F)$. They provided a mathematical formulation for the relation between the solving degree $\mathrm{sd}_\prec(F)$ and algebraic invariants coming from F, such as the maximal Gröbner basis degree, the degree of regularity, the Castelnuovo–Mumford regularity, the first and last degrees, and so on. Here, the *maximal Gröbner basis degree* of the ideal $\langle F \rangle_R$ is the maximal degree of elements in the reduced Gröbner basis of $\langle F\rangle_R$ with respect to a fixed monomial ordering $\prec$, and is denoted by $\mathrm{max.GB.deg}_{\prec}(F)$.

In the following, we recall some of Caminata et al.'s results. We set $\prec$ as the DRL ordering on R with $x_n \prec \cdots \prec x_1$, and fix it throughout the rest of this subsection. Let y be an extra variable for homogenization as in the previous subsection, and $\prec^h$ the homogenization of $\prec$, so that $y \prec x_i$ for any i with $1 \le i \le n$. Then, we have

$$\mathrm{max.GB.deg}_{\prec}(F) \le \mathrm{sd}_{\prec}(F) = \mathrm{sd}_{\prec^h}(F^h) = \mathrm{max.GB.deg}_{\prec^h}(F^h),$$

see [7] for a proof. Here, we also recall Lazard's bound for the maximal Gröbner basis degree of $\langle F^h\rangle_{R'}$ with $R' = R[y]$:

Theorem 4 (Lazard; [24, Theorem 2]) *With notation as above, we assume that the number of projective zeros of F^h is finite (and therefore $m \ge n$), and that $f_1^h = \cdots = f_m^h = 0$ has no non-trivial solution over the algebraic closure $\overline{K}$ with $y = 0$, i.e., F^{top} has no solution in $\overline{K}^n$ other than $(0,\dots,0)$. Then, supposing also that $d_1 \ge \cdots \ge d_m$, we have*

$$\mathrm{max.GB.deg}_{\prec^h}(F^h) \le d_1 + \cdots + d_\ell - \ell + 1 \tag{24}$$

with $\ell := \min\{m, n+1\}$.

Lazard's bound given in (24) is also referred to as the *Macaulay bound*, and it provides an upper-bound for the solving degree of F with respect to a DRL ordering.

As for the maximal Gröbner basis degree of $\langle F\rangle$, if $\langle F^{\mathrm{top}}\rangle$ is Aritinian, we have

$$\mathrm{max.GB.deg}_{\prec'}(F) \le d_{\mathrm{reg}}(\langle F^{\mathrm{top}}\rangle) \tag{25}$$

for any degree-compatible ordering $\prec'$ on R, see [7, Remark 15] or Lemma 4 below for a proof. Both $d_{\mathrm{reg}}(\langle F^{\mathrm{top}}\rangle)$ and $\mathrm{sd}_{\prec}(F)$ are greater than or equal to $\mathrm{max.GB.deg}_{\prec}(F)$, whereas the degree of regularity (or the first fall degree) used in the cryptographic literature as a proxy (or a heuristic upper-bound) for the solving degree. However, it is pointed out in [5], [7], and [8] by explicit examples that *any* of the degree of regularity and the first fall degree does *not* produce an estimate for the solving degree in general, even when F is an affine (cryptographic) semi-regular sequence. Caminata-Gorla proved in [8] that the solving degree is nothing but the *last fall degree* if it is greater than the maximal Gröbner basis degree:

Theorem 5 ([8, Theorem 3.1]) *With notation as above, for any degree-compatible monomial ordering $\prec'$ on R, we have the following inequality:*

$$\mathrm{sd}_{\prec'}(F) = \max\{d_F, \mathrm{max.GB.deg}_{\prec'}(F)\},$$

where d_F denotes the last fall degree of F defined in [8, Definition 1.5].

By this theorem, if $d_{\mathrm{reg}}(\langle F^{\mathrm{top}}\rangle) < d_F$, the degree of regularity is no longer an upper-bound on the solving degree.

On the other hand, Semaev and Tenti claimed that the solving degree is linear in the degree of regularity, if K is a (large) finite field, and if the input system contains polynomials related to the *field equations*, say $x_i^q - x_i$ for $1 \leq i \leq n$:

Theorem 6 ([33, Theorem 2.1], [34, Corollary 3.67]) *With notation as above, assume that $K = \mathbb{F}_q$, and that F contains $x_i^q - x_i$ for $1 \leq i \leq n$. Put $D = d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle)$. If $D \geq \max\{\deg(f) : f \in F\}$ and $D \geq q$, then we have*

$$\mathrm{sd}_{<}(F) \leq 2D - 2. \tag{26}$$

In Subsection 5.2 below, we will prove a similar inequality (26) for the case where F not necessarily contains a field equation but is cryptographic semi-regular.

4 Hilbert-Poincaré series of affine semi-regular sequence

As in the previous section, let K be a field, and $R = K[X] = K[x_1, \ldots, x_n]$ denote the polynomial ring of n variables over K. We denote by R_d the homogeneous part of degree d, that is, the set of homogeneous polynomials of degree d and 0. Recall Definition 7 for the definition of cryptographic semi-regular sequences.

The Hilbert-Poincaré series of (homogeneous) cryptographic semi-regular (resp. semi-regular) sequences is given by (22). On the other hand, the Hilbert-Poincaré series of the homogenizaton F^h cannot be computed without knowing its Gröbner basis in general, but we shall prove that it can be computed till the degree $d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle)$ if F is affine cryptographic semi-regular, namely F^{top} is cryptographic semi-regular.

Theorem 7 *Let $R = K[x_1, \ldots, x_n]$ and $R' = R[y]$, and let $F = (f_1, \ldots, f_m)$ be a sequence of not necessarily homogeneous polynomials in R. Assume that F is affine cryptographic semi-regular. Then, for each d with $d < D := d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle)$, we have*

$$\mathrm{HF}_{R'/\langle F^h \rangle}(d) = \mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(d) + \mathrm{HF}_{R'/\langle F^h \rangle}(d-1), \tag{27}$$

and hence

$$\mathrm{HF}_{R'/\langle F^h \rangle}(d) = \mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(d) + \cdots + \mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(0), \tag{28}$$

whence we can compute the value $\mathrm{HF}_{R'/\langle F^h \rangle}(d)$ from the formula (22).

Proof. Let $K_\bullet = K_\bullet(f_1^h, \ldots, f_m^h)$ by the Koszul complex associated to $(f_1^h, \ldots, f_m^h)$, which is given by (3). By tensoring $K_\bullet$ with $R'/\langle y \rangle_{R'} \cong K[x_1, \ldots, x_n] = R$ over R', we obtain the following exact sequence of chain complexes:

$$0 \longrightarrow K_\bullet \xrightarrow{\times y} K_\bullet \xrightarrow{\pi_\bullet} K_\bullet \otimes_{R'} R \longrightarrow 0,$$

where $\times y$ is a graded homomorphism of degree 1 multiplying each entry of a vector with y. Note that there is an isomorphism

$$K_i \otimes_{R'} R \cong \bigoplus_{1 \le j_1 < \cdots < j_i \le m} R(-d_{j_1 \cdots j_i}) \mathbf{e}_{j_1 \cdots j_i},$$

by which $\pi_i : K_i \to K_i \otimes_{R'} R$ reduces each entry of a vector in K_i modulo y. In particular, we have

$$\begin{aligned} K_0 \otimes_{R'} R &= R'/\langle f_1^h, \ldots, f_m^h \rangle_{R'} \otimes_{R'} R'/\langle y \rangle_{R'} \\ &\cong R'/\langle f_1^h, \ldots, f_m^h, y \rangle_{R'} \\ &\cong R/\langle f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}} \rangle_R \end{aligned}$$

for $i = 0$. This means that the chain complex $K_\bullet \otimes_{R'} R$ gives rise to the Kosuzul complex associated with $(f_1^{\mathrm{top}}, \ldots, f_m^{\mathrm{top}})$. We induce a long exact sequence of homology groups. In particular, for each degree d, we have the following long exact sequence:

$$\begin{tikzcd} H_{i+1}(K_\bullet)_{d-1} \arrow[r, "\times y"] & H_{i+1}(K_\bullet)_d \arrow[r, "\pi_{i+1}"] & H_{i+1}(K_\bullet \otimes_{R'} R)_d \arrow[dll, "\delta_{i+1}"'] \\ H_i(K_\bullet)_{d-1} \arrow[r, "\times y"'] & H_i(K_\bullet)_d \arrow[r, "\pi_i"'] & H_i(K_\bullet \otimes_{R'} R)_d \end{tikzcd}$$

where δ_{i+1} is a connecting homomorphism produced by the Snake lemma. For $i = 0$, we have the following exact sequence:

$$H_1(K_\bullet \otimes_{R'} R)_d \longrightarrow H_0(K_\bullet)_{d-1} \xrightarrow{\times y} H_0(K_\bullet)_d \longrightarrow H_0(K_\bullet \otimes_{R'} R)_d \longrightarrow 0.$$

From our assumption that F^{top} is cryptographic semi-regular, it follows from Propostion 3 that $H_1(K_\bullet \otimes_{R'} R)_{\le D-1} = 0$ for $D := d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle)$. Therefore, if $d \le D - 1$, we have an exact sequence

$$0 \longrightarrow H_0(K_\bullet)_{d-1} \xrightarrow{\times y} H_0(K_\bullet)_d \longrightarrow H_0(K_\bullet \otimes_{R'} R)_d \longrightarrow 0$$

of K-linear spaces, so that

$$\dim_K H_0(K_\bullet)_d = \dim_K H_0(K_\bullet \otimes_{R'} R)_d + \dim_K H_0(K_\bullet)_{d-1}$$

by the dimension theorem. Since $H_0(K_\bullet) = R'/\langle F^h \rangle$ and $H_0(K_\bullet \otimes_{R'} R) = R/\langle F^{\mathrm{top}} \rangle$, we have the equality (27), as desired. □

Remark 6 In the proof of Theorem 7, the multiplication map $H_0(K_\bullet)_{d-1} \to H_0(K_\bullet)_d$ by y is injective for all $d < D$, whence $\mathrm{HF}_{R'/\langle F^h \rangle}(d)$ is monotonically increasing for $d < D - 1$. On the other hand, since $H^0(K_\bullet \otimes_{R'} R)_d = (R/\langle F^{\mathrm{top}} \rangle)_d = 0$ for all $d \ge D$ by the definition of the degree of regularity, the multiplication map $H_0(K_\bullet)_{d-1} \to H_0(K_\bullet)_d$ by y is surjective for all $d \ge D$, whence $\mathrm{HF}_{R'/\langle F^h \rangle}(d)$ is monotonically decreasing for $d \ge D - 1$. By this together with [9, Theorem 3.3.4], the homogeneous ideal $\langle F^h \rangle$ is zero-dimensional or trivial, i.e., there are at most a finite number of projective zeros of F^h (and thus there are at most a finite number of affine zeros of F).

By Theorem 4, it can be proved that the Hilbert-Poincaré series of $R'/\langle F^h \rangle$ satisfies the following which may correspond to [3, Proposition 6]:

Corollary 1 *Let $D = d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle)$. Then*

$$\mathrm{HS}_{R'/\langle F^h \rangle}(z) \equiv \frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^{n+1}} \pmod{z^D}. \tag{29}$$

Therefore, by Theorem 3 ([13, Theorem 1]), F^h is D-regular. Here, we note that $D = \deg(\mathrm{HS}_{R/\langle F^{\mathrm{top}} \rangle}) + 1 = \deg\left(\left[\frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^n}\right]\right) + 1$.

Proof. Let $\mathrm{HS}'(z) = \frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^{n+1}} \bmod z^D$ and let $\mathrm{HF}'(d)$ denote the coefficient of $\mathrm{HS}'(z)$ of degree d for $d < D$. First we remark that, as F^{top} is a cryptographic semi-regular sequence, the Hilbert-Poincaré series of $R/\langle F^{\mathrm{top}} \rangle$ satisfies the following:

$$\mathrm{HS}_{R/\langle F^{\mathrm{top}} \rangle}(d) = \left[\frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^n}\right] = \frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^n} \bmod z^D,$$

since $\mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(d) = 0$ for $d \geq D$. Then we have

$$\begin{aligned}
\mathrm{HS}'(z) \bmod z^D &= \frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^{n+1}} \bmod z^D \\
&= \frac{\prod_{i=1}^{m}(1-z^{d_i})}{(1-z)^n} \times (1 + z + \cdots + z^{D-1}) \bmod z^D \\
&= \mathrm{HS}_{R/\langle F^{\mathrm{top}} \rangle}(z) \cdot (1 + z + \cdots + z^{D-1}) \bmod z^D.
\end{aligned}$$

Therefore, for $d < D$, the equation (28) gives

$$\mathrm{HF}'(d) = \mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(d) + \cdots + \mathrm{HF}_{R/\langle F^{\mathrm{top}} \rangle}(0) = \mathrm{HF}_{R'/\langle F^h \rangle}(d),$$

which implies the desired equality (29). □

To prove the following corollary, we use a fact that, for a homogeneous ideal I in R, the equality $\sum_{i=0}^{d} \dim_K I_i = \dim_K (IR')_d$ holds for each $d \geq 0$.

Corollary 2 *With notation as above, assume that $F = (f_1, \ldots, f_m) \in R^m$ is affine cryptographic semi-regular. Put $\overline{I} := \langle F^{\mathrm{top}} \rangle_R$ and $\tilde{I} := \langle F^h \rangle_{R'}$. Then, we have $(\langle \mathrm{LM}(\tilde{I}) \rangle_{R'})_d = (\langle \mathrm{LM}(\overline{I}) \rangle_{R'})_d$ for each d with $d < D := d_{\mathrm{reg}}(\overline{I})$.*

Proof. We prove $(\langle \mathrm{LM}(\tilde{I}) \rangle_{R'})_d \subset (\langle \mathrm{LM}(\overline{I}) \rangle_{R'})_d$ by the induction on d. The case where $d = 0$ is clear from Theorem 7, and so we assume $d > 0$. Any element in $(\langle \mathrm{LM}(\tilde{I}) \rangle_{R'})_d$ is represented as a finite sum of elements in R' of the form $g \cdot \mathrm{LM}(h)$ with $g \in R'$, $h \in \tilde{I}$, and $\deg(gh) = d$. Hence, we can also write each $g \cdot \mathrm{LM}(h)$

as a K-linear combination of elements of the form $\mathrm{LM}(th)$ for a monomial t in R' of dgeree $\deg(g)$, where th is an element in $\tilde{I}$ of degree d. Therefore, it suffices for showing "$\subset$" to prove that $\mathrm{LM}(f) \in (\langle \mathrm{LM}(\overline{I})\rangle_{R'})_d$ for any $f \in \tilde{I}$ with $\deg(f) = d$. We may assume that f is homogeneous. It is straightforward that $f|_{y=0} \in \overline{I}_{\leq d}$. If $\mathrm{LM}(f) \in R = K[x_1, \ldots, x_n]$, then we have $\mathrm{LM}(f) = \mathrm{LM}(f|_{y=0}) \in \mathrm{LM}(\overline{I})$. Thus, we may also assume that $y \mid \mathrm{LM}(f)$. In this case, it follows from the definition of the DRL ordering that any other term in f is also divisible by y, so that $f \in \langle y\rangle_{R'}$. Thus, we can write $f = yh$ for some $h \in R'$, where h is homogeneous of degree $d-1$. As in the proof of Theorem 7, the multiplication map

$$(R'/\tilde{I})_{d'-1} \longrightarrow (R'/\tilde{I})_{d'} \ ; \ h' \bmod \tilde{I} \mapsto yh' \bmod \tilde{I}$$

is injective for any $d' < d_{\mathrm{reg}}(\overline{I})$, since F is cryptographic semi-regular. Therefore, it follows from $f \in \tilde{I}_d$ that $h \in \tilde{I}_{d-1}$, whence $f = yh \in y\tilde{I}_{d-1}$. By the induction hypothesis, there exists $g \in \overline{I}$ such that $\mathrm{LM}(g) \mid \mathrm{LM}(h)$, whence $\mathrm{LM}(f) \in (\langle \mathrm{LM}(\overline{I})\rangle_{R'})_d$.

Here, it follows from Theorem 7 that

$$\begin{aligned}\dim_K(R')_d - \dim_K\tilde{I}_d &= \sum_{i=0}^{d}\left(\dim_K R_i - \dim_K \overline{I}_i\right) = \sum_{i=0}^{d}\dim_K R_i - \sum_{i=0}^{d}\dim_K \overline{I}_i\\ &= \dim_K(R')_d - \dim_K(\overline{I}R')_d,\end{aligned}$$

and thus $\dim_K\tilde{I}_d = \dim_K(\overline{I}R')_d$. Hence, it follows from $\langle \mathrm{LM}(\overline{I})\rangle_{R'} = \langle \mathrm{LM}(\overline{I}R')\rangle_{R'}$ that

$$\dim_K(\langle \mathrm{LM}(\tilde{I})\rangle_{R'})_d = \dim_K(\langle \mathrm{LM}(\overline{I})\rangle_{R'})_d,$$

whence $(\langle \mathrm{LM}(\tilde{I})\rangle_{R'})_d = (\langle \mathrm{LM}(\overline{I})\rangle_{R'})_d$, as desired. □

Example 1 We give a simple example. Let $p = 73$, $K = \mathbb{F}_p$, and

$$\begin{aligned}f_1 &= x_1^2 + (3x_2 - 2x_3 - 1)x_1 + x_2^2 + (-2x_3 - 2)x_2 + x_3^2 + x_3,\\ f_2 &= 4x_1^2 + (3x_2 + 4x_3 - 2)x_1 - x_2 + x_3^2 + 2x_3,\\ f_3 &= 3x_1^2 - x_1 + 9x_2^2 + (-6x_3 + 1)x_2 + x_3^2 - x_3,\\ f_4 &= x_1^2 + (-6x_2 + 2x_3 - 2)x_1 + 9x_2^2 + (-6x_3 + 1)x_2 + 2x_3^2.\end{aligned}$$

Then, $d_1 = d_2 = d_3 = d_4 = 2$. As their top parts (maximal total degree parts) are

$$\begin{aligned}f_1^{\mathrm{top}} &= x_1^2 + (3x_2 - 2x_3)x_1 + x_2^2 - 2x_3x_2 + x_3^2,\\ f_2^{\mathrm{top}} &= 4x_1^2 + (3x_2 + 4x_3)x_1 + x_3^2,\\ f_3^{\mathrm{top}} &= 3x_1^2 + 9x_2^2 - 6x_3x_2 + x_3^2,\\ f_4^{\mathrm{top}} &= x_1^2 + (-6x_2 + 2x_3)x_1 + 9x_2^2 - 6x_3x_2 + 2x_3^2,\end{aligned}$$

one can verify that F^{top} is a cryptographic semi-regular sequence. Moreover, its degree of regularity is equal to 3. Indeed, the reduced Gröbner basis G_{top} of the ideal $\langle F^{\mathrm{top}}\rangle$ with respect to the DRL ordering $x_1 \succ x_2 \succ x_3$ is

$\{x_3^2x_2, x_3^3, x_1^2{+}68x_3x_2{+}55x_3^2, x_2x_1{+}27x_3x_2{+}29x_3^2, x_2^2{+}x_3x_2{+}71x_3^2, x_3x_1{+}3x_3x_2{+}33x_3^2\}$.

Then its leading monomials are $x_3^3, x_3^2x_2, x_1^2, x_1x_2, x_2^2, x_3x_1$ and its Hilbert-Poincaré series satisfies

$$\mathrm{HS}_{R/\langle F^{\mathrm{top}}\rangle}(z) = 2z^2 + 3z + 1 = \left(\frac{(1-z^2)^4}{(1-z)^3} \bmod z^3\right),$$

whence the degree of regularity of $\langle F^{\mathrm{top}}\rangle$ is 3.

On the other hand, the reduced Gröbner basis G_{hom} of the ideal $\langle F^h\rangle$ with respect to the DRL ordering $x_1 \succ x_2 \succ x_3 \succ y$ is

$$\begin{aligned}
&\{y^3x_1, y^3x_2, y^3x_3, 60y^2x_1 + (x_3^2 + 22y^2)x_2 + 39y^2x_3,\\
&72y^2x_1 + 14y^2x_2 + x_3^3 + 56y^2x_3, 16y^2x_1 + (yx_3 + 55y^2)x_2 + 38y^2x_3,\\
&72y^2x_1 + 66y^2x_2 + yx_3^2 + 70y^2x_3, x_1^2 + 72yx_1 + (68x_3 + 40y)x_2 + 55x_3^2 + 14yx_3,\\
&(x_2 + 20y)x_1 + (27x_3 + 37y)x_2 + 29x_3^2 + 12yx_3,\\
&57yx_1 + x_2^2 + (x_3 + 3y)x_2 + 71x_3^2 + 52yx_3,\\
&(x_3 + 22y)x_1 + (3x_3 + 5y)x_2 + 33x_3^2 + 14yx_3\}
\end{aligned}$$

and its leading monomials are $y^3x_1, y^3x_2, y^3x_3, x_3^2x_2, x_3^3, yx_2x_3, yx_3^2, x_1^2, x_1x_2, x_2^2, x_1x_3$. Then the Hilbert-Poincaré series of $R'/\langle F^h\rangle$ satisfies

$$\left(\mathrm{HS}_{R'/\langle F^h\rangle}(z) \bmod z^3\right) = \left(6z^2 + 4z + 1 \bmod z^3\right) = \left(\frac{(1-z^2)^4}{(1-z)^4} \bmod z^3\right).$$

We note that $\mathrm{HF}_{R'/\langle F^h\rangle}(3) = 4$ and $\mathrm{HF}_{R'/\langle F^h\rangle}(4) = 1$. We can also examine $\mathrm{LM}(G_{\mathrm{hom}})_{d<D} = \mathrm{LM}(G_{\mathrm{top}})_{d<D}$ and, for $g \in G_{\mathrm{hom}}$, if $\mathrm{LM}(g)$ is divided by y, then $\deg(g) \geq D = 3$. Thus, at the degree 3, there occurs a *degree-fall*. See the detail in Subsection 3.2.1. Also, the reduced Gröbner basis of $\langle F\rangle$ with respect to $\prec$ is $\{x, y, z\}$ and we can examine that the dehomonization of G_{hom} is also a Gröbner basis of $\langle F\rangle$.

5 Application to Gröbner bases computation

We use the same notation as in the previous section, and assume that F is cryptographic semi-regular. Here we apply results in the previous section to the computation of Gröbner bases of ideals $\langle F\rangle$ and $\langle F^h\rangle$. Let G, G_{hom}, and G_{top} be reduced Gröbner bases of $\langle F\rangle$, $\langle F^h\rangle$, and $\langle F^{\mathrm{top}}\rangle$, respectively, where their monomial ordering are DRL $\prec$ or its extension $\prec^h$. Also we let $D = d_{\mathrm{reg}}(\langle F^{\mathrm{top}}\rangle)$.

As to the computation of G, in special settings on F such as F containing field equations or F appearing in a multivariate polynomial cryptosystem, methods using the value D or those of the Hilbert function for degrees less than D were proposed.

(See [33, 32].) Our results in the section can be considered as a *certain extension* and to *give exact mathematical proofs* for the correctness of the methods.

Here, we extend the notion of *top part* to a homogeneous polynomial h in $R' = R[y]$. We call $h|_{y=0}$ the *top part* of h and denote it by h^{top}. Thus, if h^{top} is not zero, it coincides with the top part $(h|_{y=1})^{\mathrm{top}}$ of the dehomogenization $h|_{y=1}$ of h. We remark that $g^{\mathrm{top}} = (g^h)^{\mathrm{top}}$ for a polynomial g in R.

5.1 Gröbner basis elements of degree less than D

Here we show relations between $(G_{\mathrm{hom}})_{<D}$ and $(G_{\mathrm{top}})_{<D}$ with proofs which are useful for the computations of G_{hom} and G.

Since F^{top} is cryptographic semi-regular and F^h is D-regular by Corollary 1, $H_1(K_\bullet(F^{\mathrm{top}}))_{<D} = H_1(K_\bullet(F^h))_{<D} = 0$. As $H_1(K_\bullet(F^h)) = \mathrm{syz}(F^h)/\mathrm{tsyz}(F^h)$ and $H_1(K_\bullet(F^{\mathrm{top}})) = \mathrm{syz}(F^{\mathrm{top}})/\mathrm{tsyz}(F^h)$ (see (8)), we have the following corollary, where $\mathrm{tsyz}(F^h)$ denotes the module of trivial syzygies (see Definition 1).

Corollary 3 ([13, Theorem 1]) *It follows that* $\mathrm{syz}(F^{\mathrm{top}})_{<D} = \mathrm{tsyz}(F^{\mathrm{top}})_{<D}$ *and* $\mathrm{syz}(F^h)_{<D} = \mathrm{tsyz}(F^h)_{<D}$.

This implies that, in the Gröbner basis computation G_{hom} with respect to a degree-compatible ordering $\prec^h$, if an S-polynomial $S(g_1, g_2) = t_1 g_1 - t_2 g_2$ of degree less than D is reduced to 0, it comes from some trivial syzyzy, that is, $\sum_{i=1}^m (t_1 a_i^{(1)} - t_2 a_i^{(2)} - b_i)\mathbf{e}_i$ belongs to $\mathrm{tsyz}(F^h)_{<D}$, where $g_1 = \sum_{i=1}^m a_i^{(1)} f_i^h$, $g_2 = \sum_{i=1}^m a_i^{(2)} f_i^h$, $S(g_1, g_2) = \sum_{i=1}^m b_i f_i^h$ and $\mathrm{sig}(S(g_1, g_2)) = \mathrm{LM}(\sum_{i=1}^m (t_1 a_i^{(1)} - t_2 a_i^{(2)} - b_i)\mathbf{e}_i)$ Thus, since the F_5 algorithm (or its variant such as the matrix-F_5 algorithm) with the *Schreyer ordering* automatically discards an S-polynomial whose signature is the LM of some trivial syzygy, we can avoid unnecessary S-polynomials. See Subsection 3.1.2 for a brief outline of the F_5 algorithm and the syzygy criterion (Proposition 4).

In addition to the facts above, as mentioned (somehow implicitly) in [1, Section 3.5] and [3], when we compute a Gröbner basis of $\langle F^h \rangle$ for the degree less than D by the F_5 algorithm with respect to $\prec^h$, for each computed non-zero polynomial g from an S-polynomial, say $S(g_1, g_2)$, of degree less than D, its signature does not come from any trivial syzygy and so the reductions of $S(g_1, g_2)$ are done only at its top part. This implies that the Gröbner basis computation process of $\langle F^h \rangle$ corresponds exactly to that of $\langle F \rangle$ for each degree less than D. Especially, the following lemma holds. Here we give a *concrete and easy* proof using Corollary 2. We also note that the following can be considered as a corrected version of [32, Theorem 4].

Lemma 2 *For each degree $d < D$,*

$$\mathrm{LM}(G_{\mathrm{hom}})_d = \mathrm{LM}(G_{\mathrm{top}})_d. \tag{30}$$

Proof. We can prove the equality (30) by the induction on d. Assume that the equality (30) holds for $d < D - 1$.

Consider any $t \in \mathrm{LM}(G_{\mathrm{hom}})_{d+1}$. Then, there is a polynomial $g \in G_{\mathrm{hom}}$ such that $\mathrm{LM}(g) = t$. By Corollary 2, for $d + 1 < D$, we have

$$(\langle \mathrm{LM}(\langle F^h \rangle) \rangle_{R'})_{d+1} = (\langle \mathrm{LM}(\langle F^{\mathrm{top}} \rangle_R) \rangle_{R'})_{d+1}$$

and $\mathrm{LM}(g)$ is divided by $\mathrm{LM}(g')$ for some $g' \in G_{\mathrm{top}}$. Since G_{hom} is reduced, $\mathrm{LM}(g)$ is not divisible by any monomial in $\mathrm{LM}(G_{\mathrm{hom}})_{\le d} = \mathrm{LM}(G_{\mathrm{top}})_{\le d}$, and hence $\deg(g') = d + 1$. In this case, we have $\mathrm{LM}(g) = \mathrm{LM}(g')$. Thus, $\mathrm{LM}(G_{\mathrm{hom}})_{d+1} \subset \mathrm{LM}(G_{\mathrm{top}})_{d+1}$ holds.

By the same argument, $\mathrm{LM}(G_{\mathrm{hom}})_{d+1} \supset \mathrm{LM}(G_{\mathrm{top}})_{d+1}$ can be shown. We note that for each $t \in \mathrm{LM}(G_{\mathrm{top}})_{d+1}$, there is a polynomial $g \in (G_{\mathrm{top}})_{d+1} \subset \langle F^{\mathrm{top}} \rangle_{d+1}$ such that $t = \mathrm{LM}(g)$. In this case, there are homogeneous polynomials $a_1, \ldots, a_m$ such that $g = \sum_{i=1}^m a_i f^{\mathrm{top}}$. Then $g' = \sum_{i=1}^m a_i f^h$ in $\langle F^h \rangle_{d+1}$ has t as its LM. ◻

Next we consider $(G_{\mathrm{hom}})_D$.

Lemma 3 *For each monomial M in X of degree D, there is an element g in $(G_{\mathrm{hom}})_{\le D}$ such that* $\mathrm{LM}(g)$ *divides M. Therefore,*

$$\langle \mathrm{LM}((G_{\mathrm{hom}})_{\le D}) \rangle \cap R_D = R_D. \tag{31}$$

Moreover, for each element g in $(G_{\mathrm{hom}})_D$ with $g^{\mathrm{top}} \ne 0$, the top-part g^{top} consists of one term, that is, $g^{\mathrm{top}} = \mathrm{LT}(g)$, where LT *denotes the leading term of g. (We recall* $\mathrm{LT}(g) = \mathrm{LC}(g)\mathrm{LM}(g)$.*)*

Proof. Since $\langle F^{\mathrm{top}} \rangle_D = R_D$, for each monomial M in X of degree D, there are homogeneous polynomials $a_1, \ldots, a_m$ such that $M = \sum_{i=1}^m a_i f_i^{\mathrm{top}}$. Now consider $h = \sum_{i=1}^m a_i f_i^h$ which belongs to $\langle F^h \rangle$. Then, as $f_i^h = f_i^{\mathrm{top}} + y h_i$ for some h_i in R', we have

$$h = \sum_{i=1}^m a_i (f_i^{\mathrm{top}} + y h_i) = \sum_{i=1}^m a_i f_i^{\mathrm{top}} + y \sum_{i=1}^m a_i h_i = M + y \sum_{i=1}^m a_i h_i$$

and $\mathrm{LM}(h) = M$. As G_{hom} is the reduced Gröbner basis of $\langle F^h \rangle$, there is some g in $(G_{\mathrm{hom}})_{\le D}$ whose LM divides M.

Next we prove the second assertion. Let $g_1, \ldots, g_k$ be all elements of $(G_{\mathrm{hom}})_D$ which have non-zero top parts, and set $\mathrm{LM}(g_1) \prec \cdots \prec \mathrm{LM}(g_k)$. We show that $g_i^{\mathrm{top}} = \mathrm{LT}(g_i)$ for all i. Suppose, to the contrary, that our claim does not hold for some g_i. Then, g_i^{top} can be written as $= \mathrm{LT}(g_i) + T_2 + \cdots + T_s$ for some terms $T_2, \ldots, T_s$ in R_D. Since $\mathrm{LM}(T_j) \prec \mathrm{LM}(g_i)$ for $2 \le j \le s$, it follows from equality (31) that each $\mathrm{LM}(T_j)$ is equal to $\mathrm{LM}(g_\ell)$ for some $\ell < i$ or is divisible by $\mathrm{LM}(g')$ for some $g' \in (G_{\mathrm{hom}})_{<D}$. This contradicts to the fact that G_{hom} is reduced. ◻

Remark 7 If we apply a signature-based algorithm such as the F_5 algorithm or its variant to compute the Gröbner basis of $\langle F^h \rangle$, its Σ-Gröbner basis is a Gröbner basis, but is not always *reduced* in the sense of ordinary Gröbner basis, in general.

In this case, we have to compute so called *inter-reduction* among elements of the Σ-Gröbner basis for obtaining the reduced Gröbner basis.

5.2 Gröbner basis elements of degree not less than D

In this subsection, we shall extend the upper bound on solving degree given in [33, Theorem 2.1] to our case.

Remark 8 In [33], polynomial ideals over $\mathbb{F}_q[X]$ are considered. Under the condition where the generating set F contains the field equations $x_i^q - x_i$ for $1 \leq i \leq n$, the solving degree of $\langle F \rangle$ is bounded by $2D-2$, where $D = d_{\mathrm{deg}}(\langle F^{\mathrm{top}} \rangle)$. For estimating the upper bound, the property $\langle F^{\mathrm{top}} \rangle_D = R_D$ was essentially used. As the property also holds in our case, we may apply their arguments. Also in [5, Section 3.2], the case where F^h is cryptographic semi-regular is considered. The results on the solving degree and the maximal degree of the Gröbner basis are heavily related to our result in this subsection. See Example 2.

When we use the normal selection strategy on the choice of S-polynomials, the Gröbner basis computation of $\langle F \rangle$ proceeds along with the graded structure of R in its early stage, By Lemma 2 it simulates faithfully that of $\langle F^{\mathrm{top}} \rangle$ until the degree of computed polynomial becomes $D-1$, that is, it produces $\{g|_{y=1} : g \in (G_{\mathrm{hom}})_{<D}\}$. Also, by Lemma 3, it also produces $\{g|_{y=1} : g \in (G_{\mathrm{hom}})_D, g^{\mathrm{top}} \neq 0\}$. We note that, as $g^{\mathrm{top}} \neq 0$, $g|_{y=1}$ is also produced during the computation of the Gröbner basis G. We also note that the F_5 algorithm actually uses the normal strategy.

Lemma 4 *If $D \geq \max\{\deg(f) : f \in F\}$, then the maximal Gröbner basis degree and the solving degree are bounded as follows:*

$$\mathrm{max.GB.deg}_{<}(F) \leq D \textit{ and } \mathrm{sd}_{<}(F) \leq 2D-2.$$

Proof. We consider $H = \{g|_{y=1} : g \in (G_{\mathrm{hom}})_{\leq D}, g^{\mathrm{top}} \neq 0\}$. Then H is a subset obtained at earlier stage of Gröbner basis computation of $\langle F \rangle$. By Lemma 3, $\langle \mathrm{LM}(H) \rangle$ contains all monomials in X of degree D.

Then we continue the Gröbner basis computation from H. In this *latter process*, all polynomials generated from S-polynomials are reduced by elements of H. Therefore, their LM's are reduced with respect to any monomial (in X) of degree D and thus, their degrees are not more than $D-1$. Thus, the maximal Gröbner basis degree is bounded by D, and the degree of S-polynomials dealt in the whole computation is bounded by $2D$.

Next we show that we can avoid any S-polynomial of degree $2D$ or $2D-1$.
(i) If an S-polynomial $S(g_1, g_2)$ has its degree $2D$, then $\deg(g_1) = \deg(g_2) = D$ and $\gcd(\mathrm{LM}(g_1), \mathrm{LM}(g_2)) = 1$. Then, Buchberger's criterion predicts that $S(g_1, g_2)$ is always reduced to 0.

(ii) If an S-polynomial $S(g_1, g_2)$ has its degree $2D-1$, then $\deg(g_1) = \deg(g_2) = D$, $\deg(g_1) = D, \deg(g_2) = D-1$ or $\deg(g_1) = D-1, \deg(g_2) = D$.

For the case where $\deg(g_1) = D, \deg(g_2) = D-1$ or $\deg(g_1) = D-1, \deg(g_2) = D$, we have $\gcd(\mathrm{LM}(g_1), \mathrm{LM}(g_2)) = 1$, and hence $S(g_1, g_2)$ is always reduced to 0 by Buchberger's criterion.

Finally, we consider the remaining case where $\deg(g_1) = \deg(g_2) = D$. In this case, g_1 and g_2 should belong to H and recall from Lemma 3 that both of $(g_1)^{\mathrm{top}}$ and $(g_2)^{\mathrm{top}}$ are single terms. Then $S(g_1, g_2)$ cancels the top parts of $t_1 g_1$ and $t_2 g_2$, where $S(g_1, g_2) = t_1 g_1 - t_2 g_2$ for some terms t_1 and t_2. Thus, the degree of $S(g_1, g_2)$ is less than $2D-1$. □

We refer to [7, Remark 15] for another proof of $\mathrm{max.GB.deg}_{<}(F) \leq D$.

As to the computation of G_{hom}, we have a result similar to Lemma 4. Since $\langle \mathrm{LM}(G_{\mathrm{hom}})_{\leq D} \rangle$ contains all monomials in X of degree D, for any polynomial g generated in the middle of the computation of G_{hom} the degree of X part of $\mathrm{LM}(g)$ is less than D. Because g is reduced by $(G_{\mathrm{hom}})_{\leq D}$. Thus, letting $\mathcal{U}$ be the set of all polynomials generated during the computation of G_{hom}, we have

$$\{\text{The } X\text{-part of } \mathrm{LM}(g) : g \in \mathcal{U}\} \subset \{x_1^{e_1} \cdots x_n^{e_n} : e_1 + \cdots + e_n \leq D\}.$$

As different $g, g' \in \mathcal{U}$ can not have the same X part in their leading terms, the $\#\mathcal{U}$ is bounded by the number of monomials in X of degree not greater than D, that is $\binom{n+D}{n}$. By using the F_5 algorithm or its efficient variant, under an assumption that every unnecessary S-polynomial can be avoided, the number of computed S-polynomials during the computation of G_{hom} coincides with the number $\#\mathcal{U}$ and is bounded by $\binom{n+D}{n}$.

Example 2 When $m = n+1$ and $d_1 = \cdots = d_m = 2$, the Hilbert-Poincáre series of $R/\langle F^{\mathrm{top}} \rangle$ is $\left[\frac{(1-z^2)^{n+1}}{(1-z)^n} \right]$. Since $\frac{(1-z^2)^n}{(1-z)^n} = (1+z)^n = \sum_{i=0}^n \binom{n}{i} z^i$, we have

$$\frac{(1-z^2)^{n+1}}{(1-z)^n} = (1+z)^n(1-z^2) = 1 + nz + \sum_{i=2}^{n} \left(\binom{n}{i} - \binom{n}{i-2} \right) z^i - nz^{n+1} - z^{n+2},$$

so that $D = d_{\mathrm{reg}}(\langle F^{\mathrm{top}} \rangle) = \min \left\{ i : \binom{n}{i} - \binom{n}{i-2} \leq 0 \right\} = \lfloor (n+1)/2 \rfloor + 1$. In this case, as

$$2D - 2 = 2(\lfloor (n+1))/2 \rfloor + 1) - 2 \leq n+1,$$

we have $\mathrm{sd}_{<}(F) \leq n+1$ in Lemma 4; see [5, Theorem 4.2, Therorem 4.7] for the bound in the case where F^h is a generic sequence.

We note that, in the homogeneous case, we can apply the bound on the maximal Gröbner basis degree to the solving degree. However, in our case where F is inhomogeneous, the last fall degree of F coincides with the solving degree. (See Theorem 5.) On the other hand, the solving degree of F can be bounded by the maximal Gröbner basis degree of F^h and we can apply Theorem 4, as our case satisfies its conditions. Then, for the case where $m = n+1$ and $d_1 = \cdots = d_{n+1} = 2$, it gives the bound $n+2$ for $\mathrm{max.GB.deg}_{<^h}(F^h) = \mathrm{sd}_{<^h}(F^h) \geq \mathrm{sd}_{<}(F)$.

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